15 Efficient meshes

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Basics of efficiency for meshes

Use triangle or quad meshes

- general polygon meshes lead to too much complexity
- quad meshes are great for some applications but more constrained

Use shared-vertex triangle meshes for GPU applications

- major memory/bandwidth savings over separate triangles
- · if you get separate triangles, merge them in a pre-process

Store most data at vertices

- there are ~half as many vertices as faces
- vertex data may be interpolated across faces
- · in typical GPU mesh representation, vertices must be duplicated to create discontinuities

More sophistication in mesh storage

Optimizing vertex order

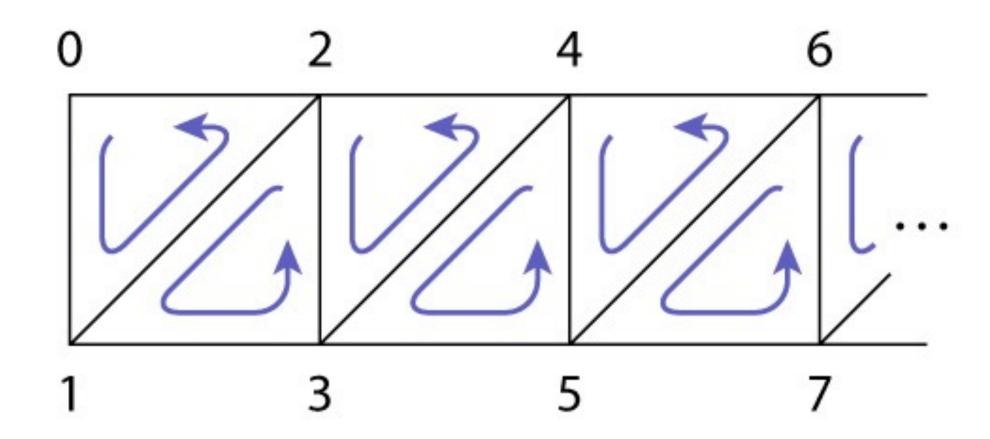
- strips and fans as classic examples (when per-frame bandwidth was the concern)
- modern systems don't use these but optimize for hit rate in vertex cache

Reducing the number of triangles

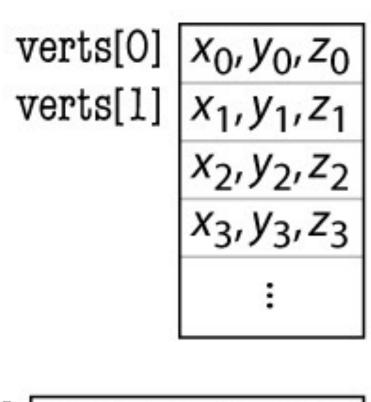
- ultimately this is needed to save more time and space
- many levels of detail are useful
 - simpler meshes for faraway objects
 - simpler meshes for lower-resolution screens
 - simpler meshes for lower-performance hardware or networks

Take advantage of the mesh property

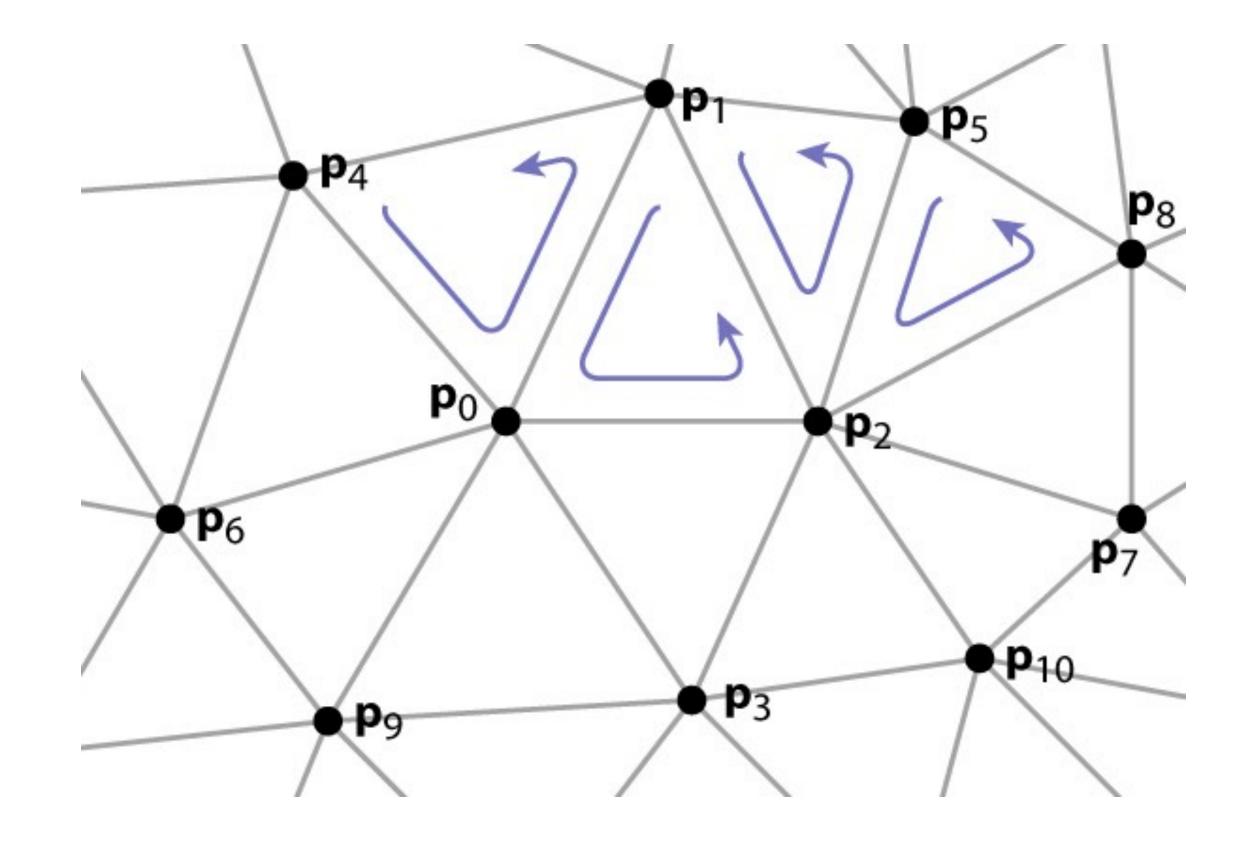
each triangle is usually adjacent to the previous

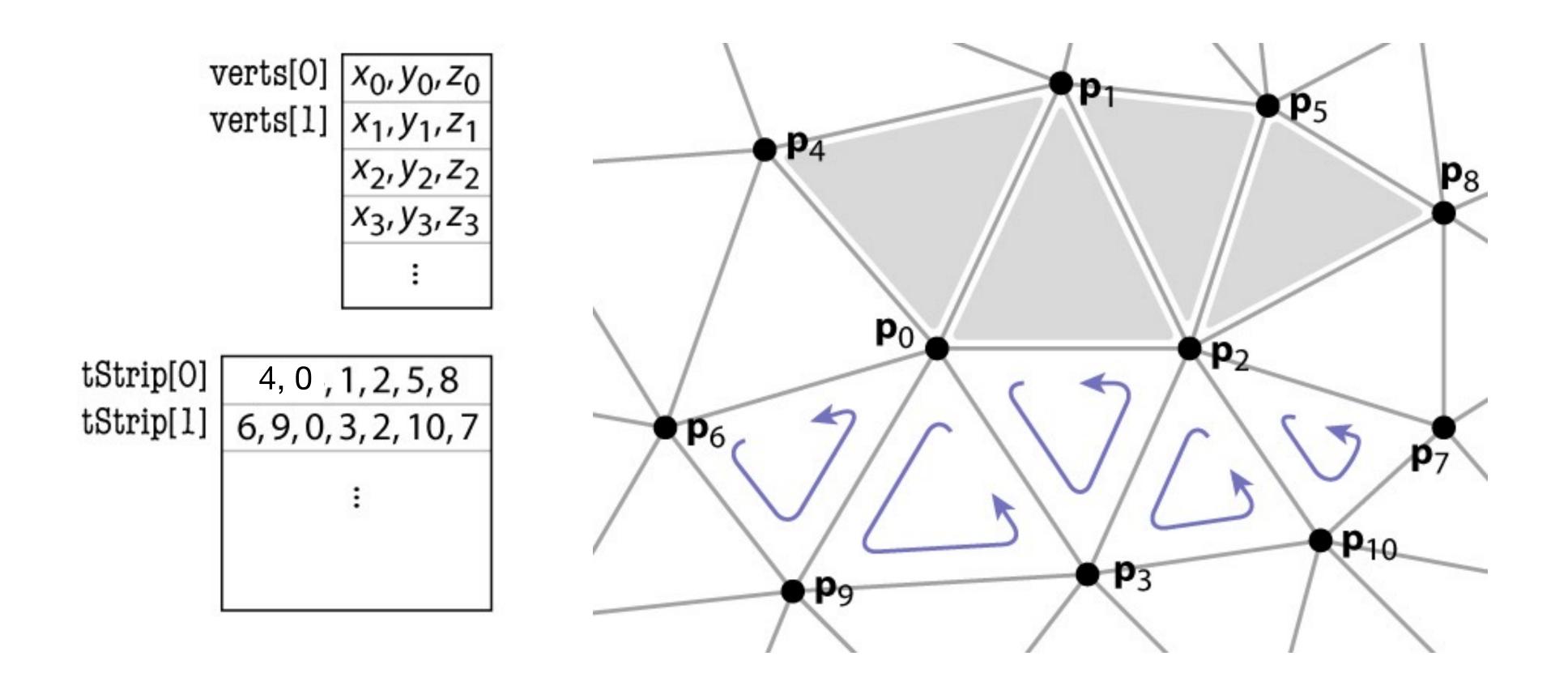


- let every vertex create a triangle by reusing the second and third vertices of the previous triangle
- every sequence of three vertices produces a triangle (but not in the same order)
- e. g., 0, 1, 2, 3, 4, 5, 6, 7, ... leads to
 (0 1 2), (2 1 3), (2 3 4), (4 3 5), (4 5 6), (6 5 7), ...
- for long strips, this requires about one index per triangle



tStrip[0] 4, 0, 1, 2, 5, 8 tStrip[1] 6, 9, 0, 3, 2, 10, 7





array of vertex positions

- float[n_V][3]: 12 bytes per vertex
 - (3 coordinates x 4 bytes) per vertex

array of index lists

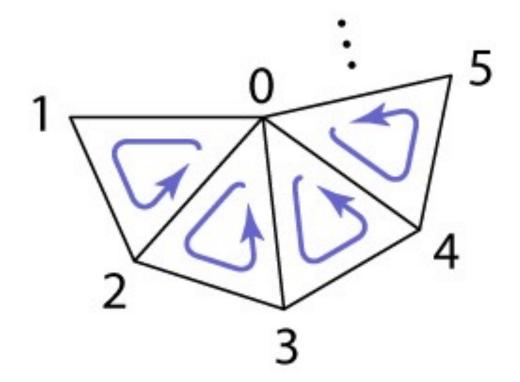
- $\inf[n_{S}][variable]: 2 + n indices per strip$
- on average, $(I + \varepsilon)$ indices per triangle (assuming long strips)
 - 2 triangles per vertex (on average)
 - about 4 bytes per triangle (on average)

total is 20 bytes per vertex (limiting best case)

- factor of 3.6 over separate triangles; 1.8 over indexed mesh

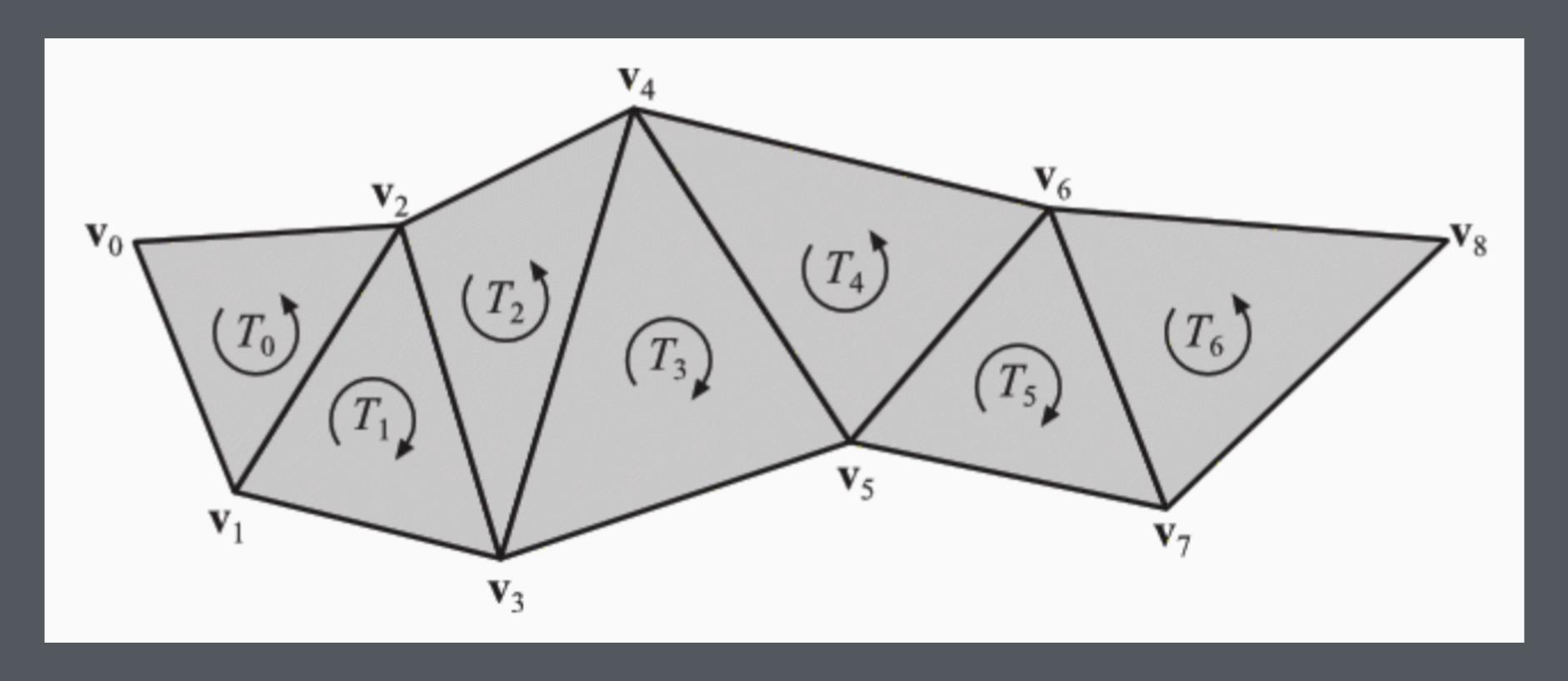
Triangle fans

- Same idea as triangle strips, but keep oldest rather than newest
 - every sequence of three vertices produces a triangle
 - e. g., 0, 1, 2, 3, 4, 5, ... leads to(0 1 2), (0 2 3), (0 3 4), (0 4 5), ...
 - for long fans, this requires
 about one index per triangle
- Memory considerations exactly the same as triangle strip



Vertex cache and mesh ordering

Triangle strips gain efficiency by caching the most recent two vertices



- we are essentially using a FIFO cache policy with a size of 2
- cache miss rate approaches 1 miss / triangle

Optimizing for larger caches

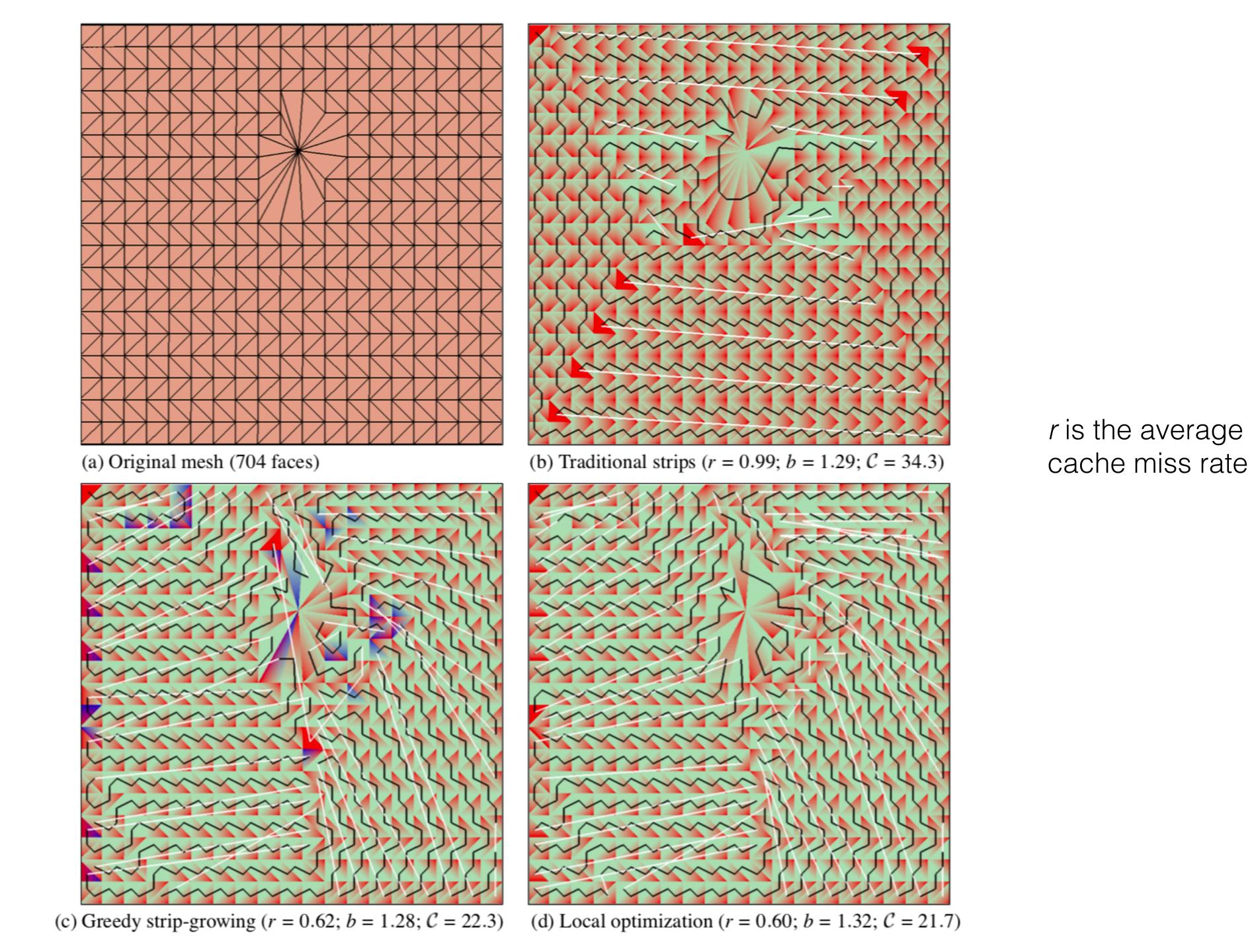
With indexed meshes, saving indices is less important

- we store lots of data at vertices; ~6 indices is the least of our worries
- just putting meshes in triangle-strip order gives you the same vertex caching behavior ("transparent" vertex caching)

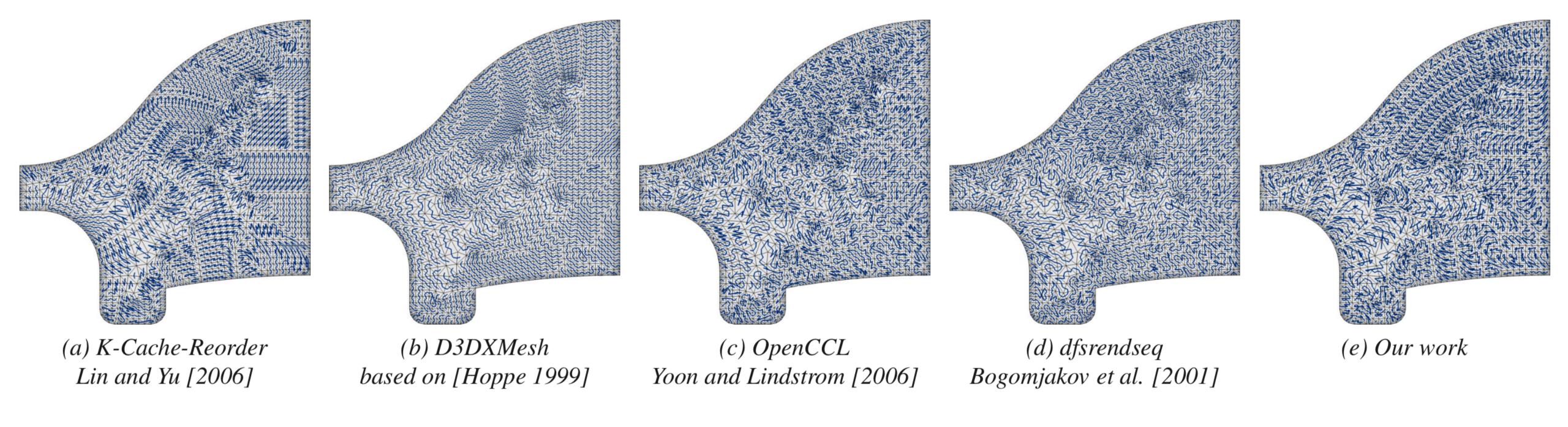
GPU pipelines are built with post-transform vertex caches

- cache the results of the vertex processing stage
- cache hits can save substantial computation
- (for parallelism newer systems process primitives in batches, but the effect is similar)

As with other applications of caches, now order of data access matters



[Hoppe 1999]

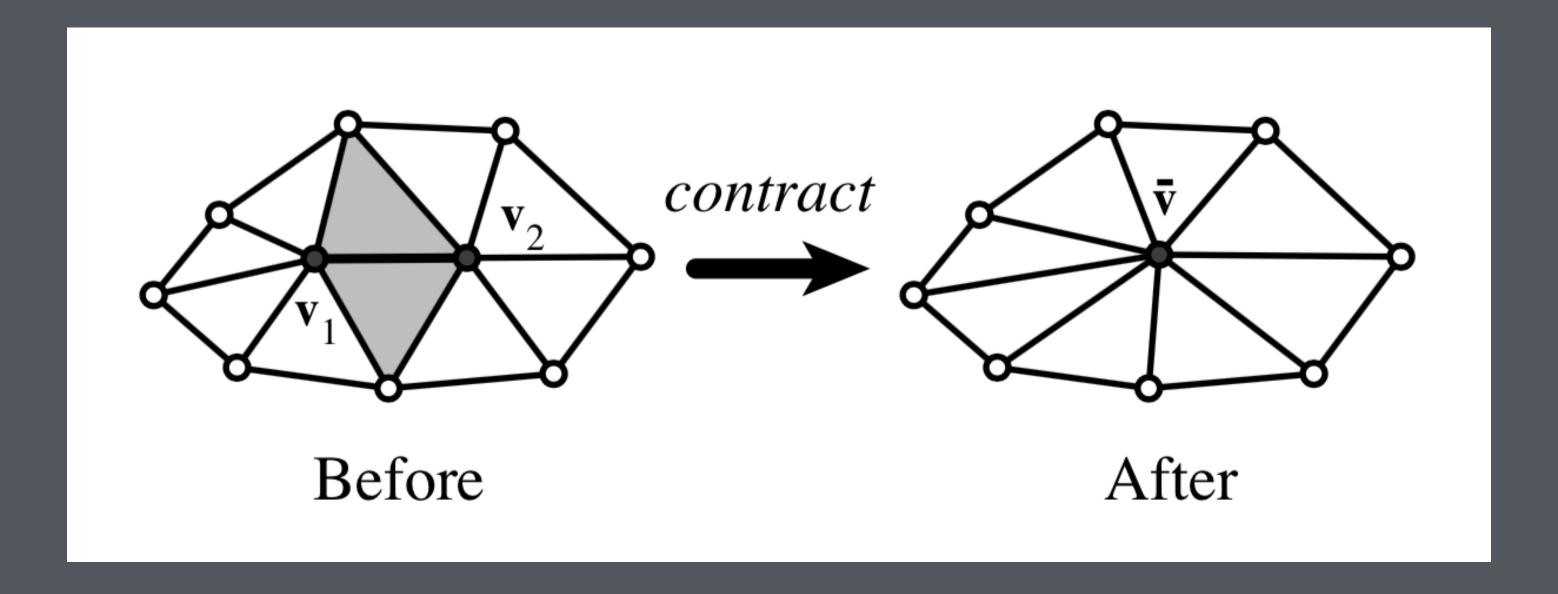


Mesh simplification

Many ways to simplify meshes

- remove chunks, retriangulate hole
- quantize vertices to centers of voxels

Particularly simple and effective is edge collapses, or edge contractions:



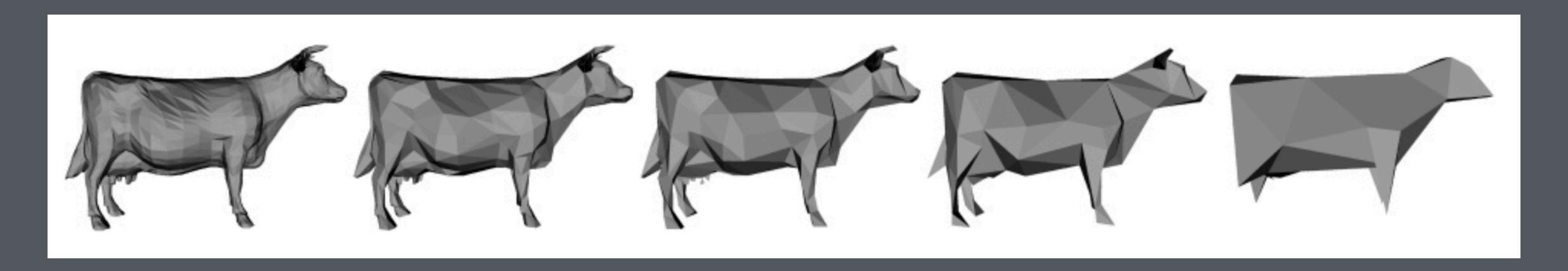
Quadric Error Metric

Edge-collapse simplification produces a sequence of meshes

- each mesh has one fewer face
- · each is derived from the previous by a single edge collapse

Key question: where to put the vertex after the collapse?

- at first vertex? at second? at midpoint?
- can choose location as the solution to an optimization



Where to put the new vertex?

It depends on the mesh geometry:

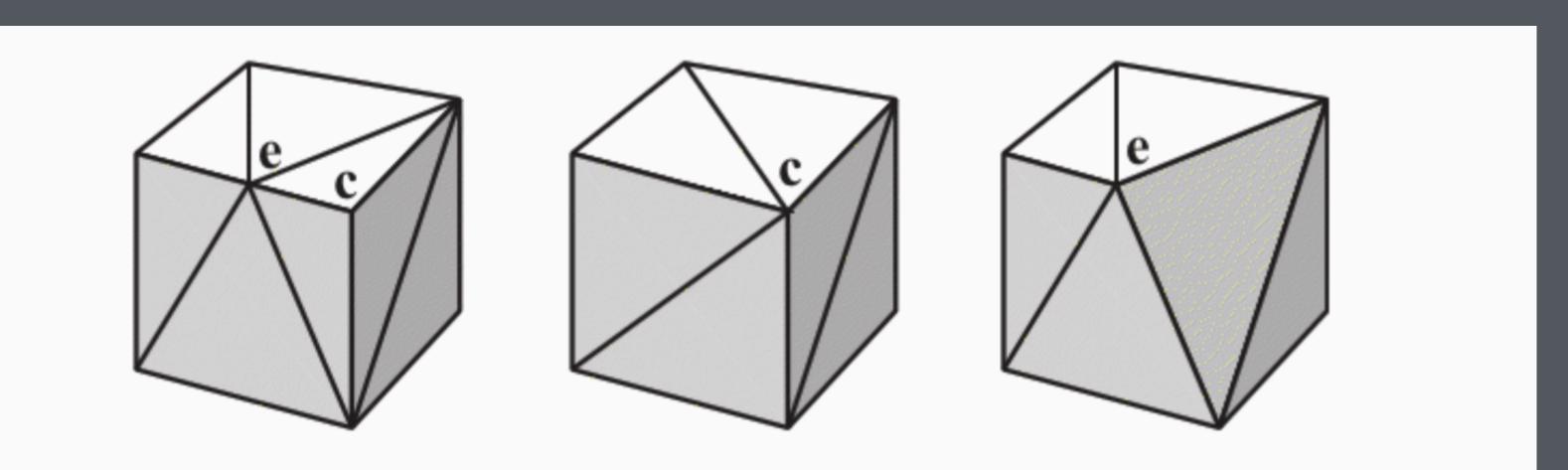


Figure 16.19. The left figure shows a cube with an extra point along one edge. The middle figure shows what happens if this point \mathbf{e} is collapsed to corner \mathbf{c} . The right figure shows \mathbf{c} collapsed to \mathbf{e} .

 one way to formalize: the new vertex should be close to the planes of the triangles around it before the edge collapse

Garland & Heckbert QEM

A particularly convenient error metric: sum of squared distances to planes

- each plane has an equation, can be represented as a 4-vector (a, b, c, d) with (a, b, c) components normalized
- distance of a vertex v from the plane p is then the inner product p^Tv
- squared distance from plane is in the form v^TMv for a 4x4 M (a quadric)

$$\Delta(\mathbf{v}) = \sum_{\mathbf{p} \in \text{planes}(\mathbf{v})} (\mathbf{v}^{\mathsf{T}} \mathbf{p}) (\mathbf{p}^{\mathsf{T}} \mathbf{v})$$

$$= \sum_{\mathbf{p} \in \text{planes}(\mathbf{v})} \mathbf{v}^{\mathsf{T}} (\mathbf{p} \mathbf{p}^{\mathsf{T}}) \mathbf{v}$$

$$= \mathbf{v}^{\mathsf{T}} \left(\sum_{\mathbf{p} \in \text{planes}(\mathbf{v})} \mathbf{K}_{\mathbf{p}} \right) \mathbf{v}$$

• and better yet, the sum-squared distance from several planes is still in the form $\mathbf{v}^{\mathsf{T}}\mathbf{Q}\mathbf{v}$

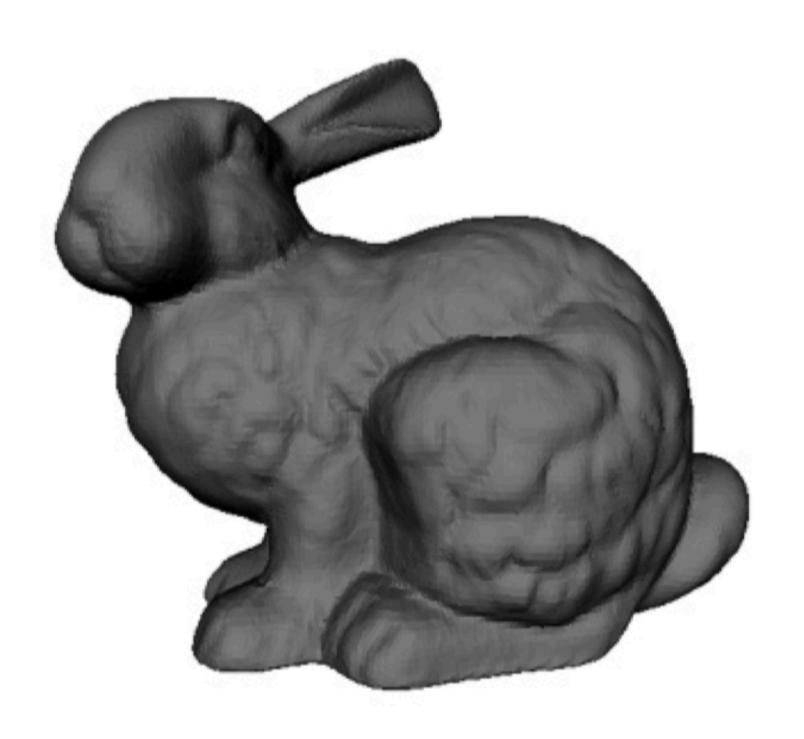
QEM simplification

With the error in the form of a quadric per vertex:

- the matrix is easy to compute from the surrounding triangles
- the error is easy to optimize. Given \mathbf{Q}_1 and \mathbf{Q}_2 belonging to a pair of vertices \mathbf{v}_1 and \mathbf{v}_2 , we simply sum the errors of the two vertices:

$$egin{aligned} \Delta(\mathbf{v}) &= \Delta_1(\mathbf{v}) + \Delta_2(\mathbf{v}) \ &= \mathbf{v}^T \mathbf{Q}_1 \mathbf{v} + \mathbf{v}^T \mathbf{Q}_2 \mathbf{v} \ &= \mathbf{v}^T (\mathbf{Q}_1 + \mathbf{Q}_2) \mathbf{v} \end{aligned}$$

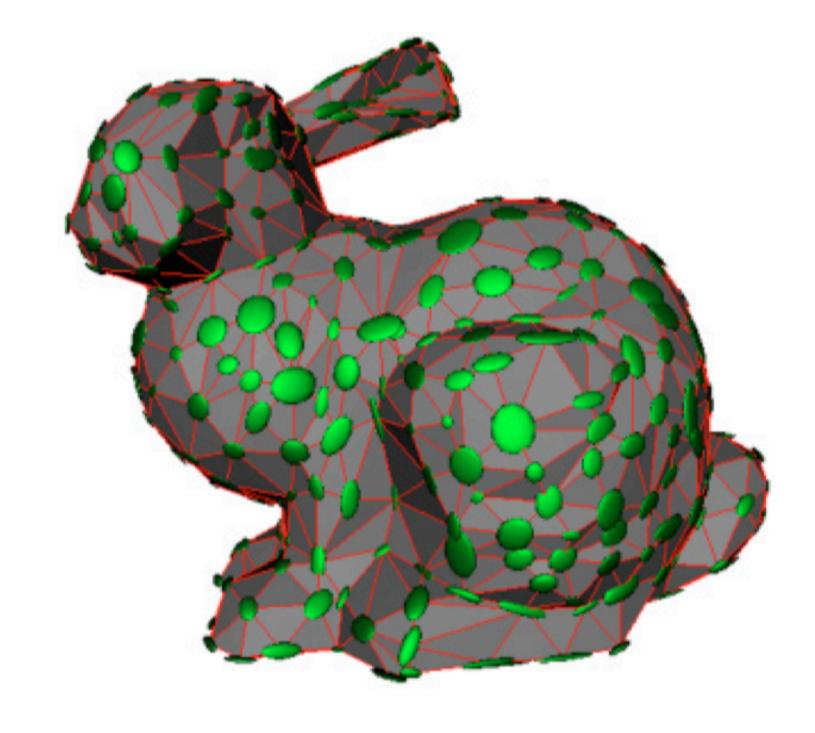
- minimizing this error is a 4x4 linear system—very fast
- algorithm
 - 0. compute Qs for all vertices, compute errors for all potential edge collapses.
 - 1. use priority queue to find smallest-error edge. Collapse it; update the neighboring Qs.
 - 2. repeat until mesh is small enough!







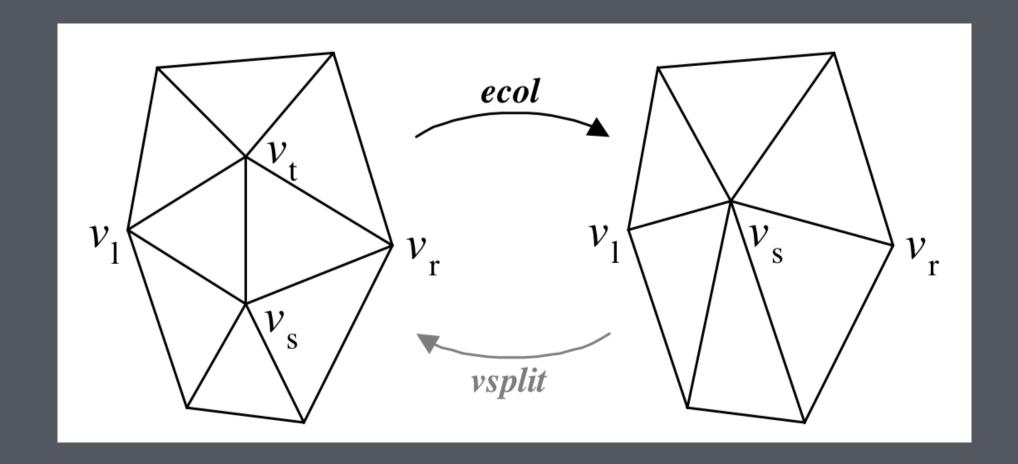
1k faces



surfaces of constant cost for reducing to 999 faces

Continuous level-of-detail: Progressive Meshes

Key observation: edge collapse is invertible



just need to store (offsets to) the locations of the two new vertices

Thus a sequence of edge collapses, reversed, is a representation for a mesh

$$(\hat{M}=M^n) \stackrel{ecol_{n-1}}{\longrightarrow} \dots \stackrel{ecol_1}{\longrightarrow} M^1 \stackrel{ecol_0}{\longrightarrow} M^0$$
.

$$M^0 \xrightarrow{vsplit_0} M^1 \xrightarrow{vsplit_1} \dots \xrightarrow{vsplit_{n-1}} (M^n = \hat{M})$$

Progressive Meshes

Store full representation, load various levels of detail

- just load or transmit a prefix of the list of edge splits
- can change level of detail smoothly depending on size/distance/salience/etc.

Can interpolate ("geomorphs")

- sudden edge splits/collapses are jarring
- interpolate new vertices from merged position to new positions
- leads to truly continuous LoD

Extra details (of QEM and PM)

- boundaries, creases—want to preserve them
- · merging of small pieces—otherwise can't simplify enough
- maintenance of additional attributes—throw them in the metric too

