

CS 5154: Software Testing

Active Clause Coverage

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Recall the four software models in this course

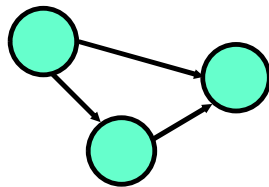


Input
Domains

```
A: {0, 1, >1}
B: {600, 700, 800}
C: {cs, ece, is, sds}
```



Graphs



Logic
Expressions

```
(!x | !y) & a & b
```

Syntax

```
if (x > y)
  z = x - y;
else
  z = 2 * x;
```

We need criteria that are not as costly as CoC

- The general idea is quite simple:

Test each clause independently from the other clauses

- But, getting the details right is hard
 - e.g., what exactly does “independently” mean ?
- The book presents this idea as “*making clauses active*” ...

Active Clauses

- A **weakness of Clause Coverage**: values do not always make a difference
- *Values* $((5 < 10) \vee \text{true}) \wedge (1 \geq 1 * 1)$ for $((a < b) \vee D) \wedge (m \geq n * o)$
 - Only the last clause counts!
- To really test the results of a clause, the clause should be the **determining factor** in what the predicate evaluates to

Determination

A clause c_i in predicate p , called the *major clause*, determines p if and only if the values of the remaining *minor clauses* c_j are such that changing c_i changes the value of p

- Making c_i determine p is said to **make the clause active**
- Condition under which c_i determines p

$\forall c_j \in C_p \setminus c_i \exists$ assignment (c_j) s.t. $p(c_i = \text{true}) \neq p(c_i = \text{false})$
where a assignment (c_j) is $c_j = \text{true}$ or $c_j = \text{false}$

Active clause criteria

The essence of testing with determination

1. Pick one clause in predicate p to be the major clause c_i
2. Find conditions under which c_i determines p
3. Find a test that makes c_i true and a test that makes c_i false
4. Repeat steps 1 to 3 for all other clauses in p
5. Eliminate redundant tests

Examples: determining predicates

$$\underline{P = A \vee B}$$

if $B = \text{true}$, p is always true.
so if $B = \text{false}$, A determines p .
if $A = \text{false}$, B determines p .

$$\underline{P = A \wedge B}$$

if $B = \text{false}$, p is always false.
so if $B = \text{true}$, A determines p .
if $A = \text{true}$, B determines p .

	a	b	$a \vee b$
1	T	T	T
2	T	F	T
3	F	T	T
4	F	F	F

	a	b	$a \wedge b$
1	T	T	T
2	T	F	F
3	F	T	F
4	F	F	F

More examples: determining predicates

$$\underline{P = A \oplus B}$$

if $B = \text{true}$, A determines p .

if $B = \text{false}$, A determines p .

so, A determines p for any B .

$$\underline{P = A \leftrightarrow B}$$

if $B = \text{true}$, A determines p .

if $B = \text{false}$, A determines p .

so, A determines p for any B .

	a	b	$a \oplus b$
1	T	T	F
2	T	F	T
3	F	T	T
4	F	F	F

	a	b	$a \leftrightarrow b$
1	T	T	T
2	T	F	F
3	F	T	F
4	F	F	T

Testing with determination 😊

- **Goal** : Find tests for each clause when that clause determines the value of the predicate
- This goal is formalized in a **family of criteria** that have subtle, but very important, differences

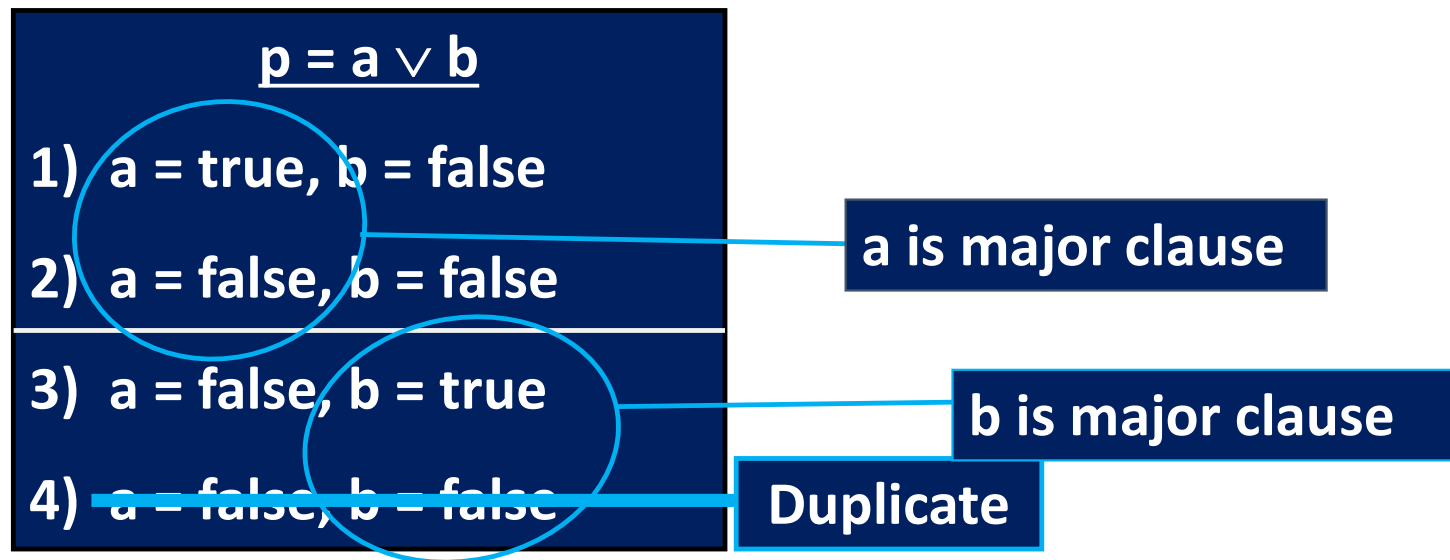
Active Clause Coverage

- Step 1: For each p in P and each major clause c_i in C_p , choose minor clauses $c_j, j \neq i$, so that c_i determines p .

Active Clause Coverage (ACC) : TR has two requirements for each c_i : c_i evaluates to true and c_i evaluates to false.

- ACC is a form of Multiple Condition Decision Coverage (MCDC)
- MCDC is required by the FAA for safety-critical software

Example on Active Clause Coverage



A formulaic way of determining predicates?

- Finding values for minor clauses c_j is easy for simple predicates
- How to find values for more complicated predicates?
- We need some “formula” that is easy to apply

A definitional way: when does c determine p ?

- Let $p_{c=true}$ be predicate p with every occurrence of c replaced by *true*
- Let $p_{c=false}$ be predicate p with every occurrence of c replaced by *false*
- To find values for the minor clauses, connect $p_{c=true}$ and $p_{c=false}$ with XOR

$$p_c = p_{c=true} \oplus p_{c=false}$$

- After solving, p_c describes exactly the values needed for c to determine p

An example using the definitional way

- Let $p = a \vee (b \wedge c)$. What values of b and c will cause a to determine p ?

$$\begin{aligned} & \underline{p = a \vee (b \wedge c)} \\ p_a &= p_{a=\text{true}} \oplus p_{a=\text{false}} \\ &= (\text{true} \vee (b \wedge c)) \oplus (\text{false} \vee (b \wedge c)) \\ &= \text{true} \oplus (b \wedge c) \\ &= !(b \wedge c) \\ &= !b \vee !c \end{aligned}$$

- “ $!b \vee !c$ ” means a determines p when either b or c is false

Exercise 1: using the definitional way

- Let $p = a \vee b$. What values of b will cause a to determine p ?

$$\begin{aligned} p &= a \vee b \\ p_a &= p_{a=\text{true}} \oplus p_{a=\text{false}} \\ &= (\text{true} \vee b) \oplus (\text{false} \vee b) \\ &= \text{true} \oplus b \\ &= !b \end{aligned}$$

	a	b	$a \vee b$
1	T	T	T
2	T	F	T
3	F	T	T
4	F	F	F

- “ $!b$ ” means a determines p when b is false
- We obtained the same result from reasoning about the truth table

Exercise 2: using the definitional way

- Let $p = a \leftrightarrow b$. What values of b will cause a to determine p ?

$$\begin{aligned} & \underline{p = a \leftrightarrow b} \\ p_a &= p_{a=\text{true}} \oplus p_{a=\text{false}} \\ &= (\text{true} \leftrightarrow b) \oplus (\text{false} \leftrightarrow b) \\ &= b \oplus !b \\ &= \text{true} \end{aligned}$$

	a	b	$a \leftrightarrow b$
1	T	T	T
2	T	F	F
3	F	T	F
4	F	F	T

- “*true*” means that a always determines p
- We obtained the same result from reasoning about the truth table

Is there a problem with Active Clause Coverage?

- Step 1: For each p in P and each major clause c_i in C_p , choose minor clauses $c_j, j \neq i$, so that c_i determines p .

Active Clause Coverage (ACC) : TR has two requirements for each c_i : c_i evaluates to true and c_i evaluates to false.

- **Ambiguity** : Must minor clauses have the **same values** when the major clause is true and when the major clause is false?

Illustrating the ambiguity in ACC

- Recall: a determines p when “ $!b \vee !c$ ”, i.e., when either b or c is false

$$p = a \vee (b \wedge c)$$

Major clause : a

a = true, b = false, c = true

a = false, b = false, c = false

Is this allowed ?

Three options for resolving ACC ambiguity

- Minor clauses **do not** need to be the same
- Minor clauses **must** be the same
- Minor clauses **allow the predicate** to become both true and false

Option 1: minor clauses **don't** need to be the same

- Step 1: For each p in P and each major clause c_i in C_p , choose minor clauses $c_j, j \neq i$, so that c_i determines p .
- Step 2 (ACC): TR has two requirements for each c_i : c_i evaluates to true and c_i evaluates to false.

General Active Clause Coverage (GACC) : The values chosen for the minor clauses c_j do not need to be the same when c_i is true as when c_i is false, that is, $c_j(c_i = true) = c_j(c_i = false)$ for all c_j OR $c_j(c_i = true) \neq c_j(c_i = false)$ for all c_j .

Problem: GACC doesn't subsume Predicate Coverage

Major clause : a

	a	b	$a \leftrightarrow b$
1	T	T	T
2	T	F	F
3	F	T	F
4	F	F	T

Option 2: minor clauses **do** need to be the same

- Step 1: For each p in P and each major clause c_i in C_p , choose minor clauses $c_j, j \neq i$, so that c_i determines p .
- Step 2 (ACC): TR has two requirements for each c_i : c_i evaluates to true and c_i evaluates to false.

Restricted Active Clause Coverage (RACC) : The values chosen for the minor clauses c_j must be the same when c_i is true as when c_i is false, that is, it is required that $c_j(c_i = true) = c_j(c_i = false)$ for all c_j .

Exercise 3: using the definitional way

- Let $p = a \wedge (b \vee c)$. What values of b and c will cause a to determine p ?

$$\begin{aligned} & \underline{p = a \wedge (b \vee c)} \\ p_a &= p_{a=\text{true}} \oplus p_{a=\text{false}} \\ &= (\text{true} \wedge (b \vee c)) \oplus (\text{false} \wedge (b \vee c)) \\ &= (b \vee c) \oplus \text{false} \\ &= (b \vee c) \\ &= b \vee c \end{aligned}$$

- “ $b \vee c$ ” means a determines p when either b or c is true

Example on Restricted Active Clause Coverage

Major clause : a , $P_a = \mathbf{b} \vee \mathbf{c}$

	a	b	c	$a \wedge (b \vee c)$
1	T	T	T	T
2	T	T	F	T
3	T	F	T	T
4	T	F	F	F
5	F	T	T	F
6	F	T	F	F
7	F	F	T	F
8	F	F	F	F

RACC ($c_i = a$) can only be satisfied by row pairs (1, 5), (2, 6), or (3, 7)
Only three pairs can be used

Notes on RACC

- Does RACC subsume predicate and clause coverage?
- RACC was a **common interpretation** by developers for FAA
- Problem: RACC often leads to **infeasible test requirements**

Option 3: minor clauses **allow** predicate to be true and false

- Step 1: For each p in P and each major clause c_i in C_p , choose minor clauses $c_j, j \neq i$, so that c_i determines p .
- Step 2 (ACC): TR has two requirements for each c_i : c_i evaluates to true and c_i evaluates to false.

Correlated Active Clause Coverage (CACC) : The values chosen for the minor clauses c_j must cause p to be true for one value of the major clause c_i and false for the other, that is, it is required that $p(c_i = true) \neq p(c_i = false)$.

Example on CACC

	a	b	c	$a \wedge (b \vee c)$
1	T	T	T	T
2	T	T	F	T
3	T	F	T	T
4	T	F	F	F
5	F	T	T	F
6	F	T	F	F
7	F	F	T	F
8	F	F	F	F

a determines P when (b=true or c = true)

CACC ($c_i = a$) can be satisfied by choosing any of rows 1, 2, 3 AND any of rows 5, 6, 7 – a total of nine pairs

Notes on CACC

- CACC **implicitly** allows minor clauses to have different values
- CACC explicitly **subsumes** predicate coverage
- Does CACC subsume clause coverage?

Does CACC subsume clause coverage?

	a	b	c	$a \wedge (b \vee c)$
1	T	T	T	T
2	T	T	F	T
3	T	F	T	T
4	T	F	F	F
5	F	T	T	F
6	F	T	F	F
7	F	F	T	F
8	F	F	F	F

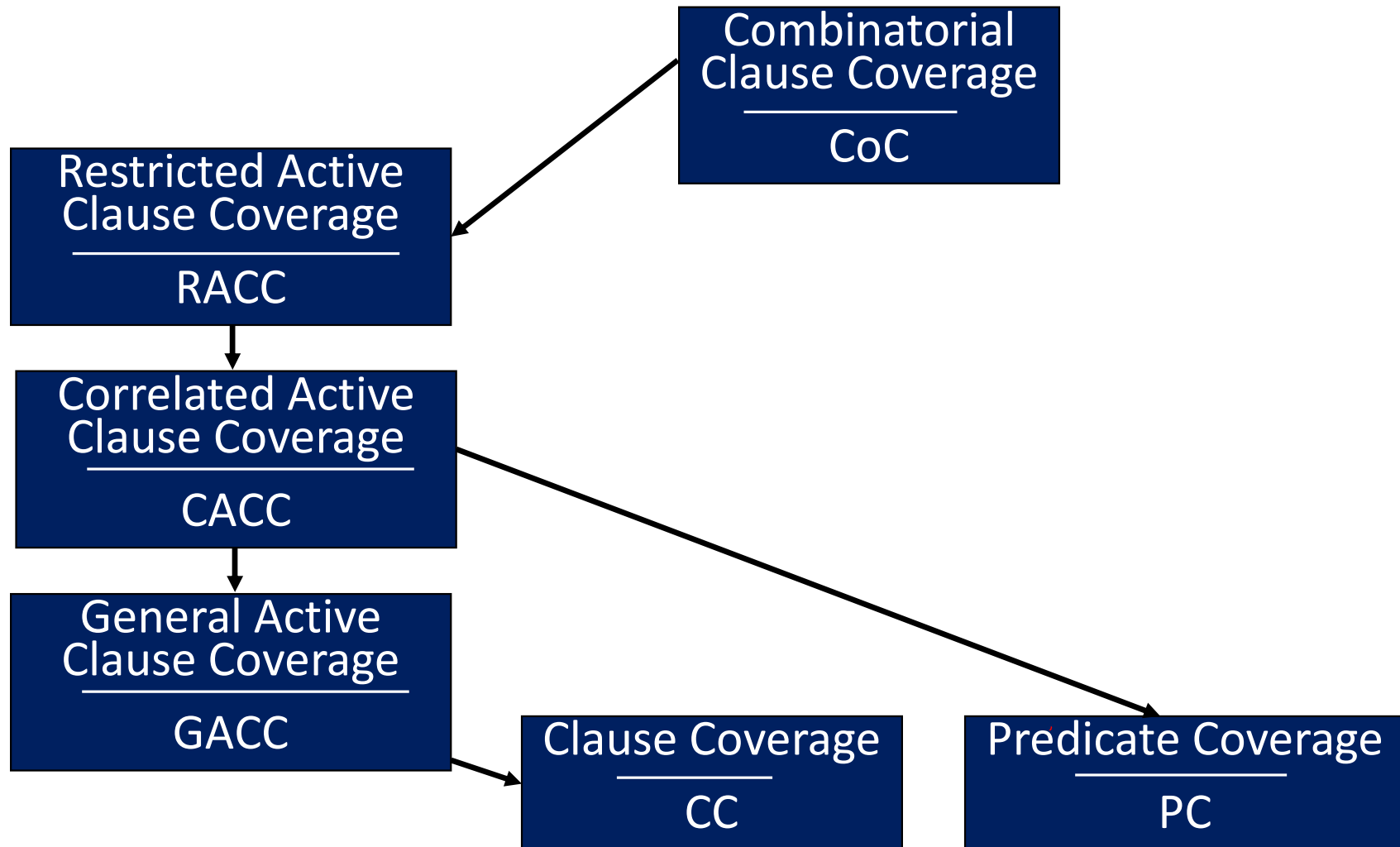
a determines P when (b=true or c = true)

CACC ($c_i = a$) can be satisfied by choosing any of rows 1, 2, 3 AND any of rows 5, 6, 7 – a total of nine pairs

Infeasibility

- Consider the predicate: $(a > b \wedge b > c) \vee c > a$
- *Infeasible*: $(a > b) = \text{true}, (b > c) = \text{true}, (c > a) = \text{true}$ is **infeasible**
- As with other criteria, infeasible test requirements must be **recognized** and **dealt with**
- Recognizing infeasible test requirements is hard, and in general, **undecidable**

Subsumption among Logic coverage criteria



An end-to-end example with RACC

b & c are false, cause a to determine the value of p
 Likewise, for clause c , only one pair, TFF and TFF, cause c to determine the value of p

	a	b	c	$a \wedge (b \vee c)$	p_a	p_b	p_c
1	T	T	T	T	★		
2	T	T	F	T	★	★	
3	T	F	T	T			★
4	T	F	F	F		★	★
5	F	T	T	F	★		
6	F	T	F	F	★		
7	F	F	T	F			★
8	F	F	F	F			

In sum, three separate pairs of rows can cause a to determine the value of p , and only one pair each for b and c

How many tests does RACC yield, compared to Combinatorial Clause Coverage?

A more subtle exercise on determination

$$\underline{p = (a \wedge b) \vee (a \wedge !b)}$$

$$\begin{aligned} p_a &= p_{a=\text{true}} \oplus p_{a=\text{false}} \\ &= ((\text{true} \wedge b) \vee (\text{true} \wedge !b)) \oplus ((\text{false} \wedge b) \vee (\text{false} \wedge !b)) \\ &= (b \vee !b) \oplus \text{false} \\ &= \text{true} \oplus \text{false} \\ &= \text{true} \end{aligned}$$

$$\underline{p = (a \wedge b) \vee (a \wedge \neg b)}$$

$$\begin{aligned} p_b &= p_{b=\text{true}} \oplus p_{b=\text{false}} \\ &= ((a \wedge \text{true}) \vee (a \wedge \neg \text{true})) \oplus ((a \wedge \text{false}) \vee (a \wedge \neg \text{false})) \\ &= (a \vee \text{false}) \oplus (\text{false} \vee a) \\ &= a \oplus a \\ &= \text{false} \end{aligned}$$

A more subtle exercise on determination (2)

$$p = (a \wedge b) \vee (a \wedge \neg b)$$

- a always determines the value of this predicate
- b never determines the value – b is irrelevant !
- So, why would anyone write a predicate like this?

Logic Coverage Summary

- Predicates are often *very simple*—in practice, most have <3 clauses
 - In fact, most predicates only have one clause !
- Only clause? PC is enough
- 2 or 3 clauses? CoC is practical
- Advantages of ACC criteria can be significant for large (no. of) predicates
 - CoC is impractical for predicates with many clauses

Next

- Applying Logic Coverage to source code