## CS 5154

## Graph Coverage Criteria

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The following are modified versions of the publicly-available slides for Chapter 7 in the Ammann and Offutt Book, "Introduction to Software Testing"
(http://www.cs.gmu.edu/~offutt/softwaretest)

## Graph Coverage

Four Structures for Modeling Software


Use cases


## Covering Graphs

- Graphs are the most commonly used structure for testing
- Graphs can come from many sources
- Control flow graphs
- Design structure
- FSMs and statecharts
- Use cases
- Tests usually are intended to "cover" the graph somehow


## Why Graph Coverage?

- Some of the most widely-used coverage criteria
- The "R" in the RIPR model
- Graph coverage criteria help create tests that reach different parts of software


## The next two classes

- Today:
- Review of graph concepts
- Coverage criteria defined over generic graphs
- Next class (depending on progress today):
- Apply concepts learned in today's class to source code
- Not in CS 5154
- Applying graph coverage criteria to design, specs, and use cases


## Definition of a Graph

- A set $N$ of nodes, $N$ is not empty
- A set $N_{0}$ of initial nodes, $N_{0}$ is not empty
- A set $N_{f}$ of final nodes, $N_{f}$ is not empty
- A set $E$ of edges, each edge from one node to another
- ( $\left.n_{i}, n_{j}\right), i$ is predecessor, $j$ is successor


$$
\begin{aligned}
N_{0} & =\{1\} \\
N_{f} & =\{1\} \\
E & =\{ \}
\end{aligned}
$$

## Example Graphs



## Paths in Graphs

- Path : A sequence, $p$, of nodes $\left[n_{1}, n_{2}, \ldots, n_{M}\right]$ s.t. there is an edge between each pair of nodes in $p$
- Length of a path : The number of edges in $p$
- A single node is a path of length 0
- Subpath : A subsequence of nodes in $p$ is a subpath of $p$



## Test Paths and SESE graphs

- Test Path : A path that starts at an initial node and ends at a final node
- Test paths represent execution of test cases
- Some test paths can be executed by many tests
- Some test paths cannot be executed by any tests
- SESE graphs : All test paths start at a single node and end at another node
- Single-entry, single-exit
- $\mathrm{N}_{0}$ and $\mathrm{N}_{\mathrm{f}}$ have exactly one node


Double-diamond graph Four test paths

$$
\begin{aligned}
& {[1,2,4,5,7]} \\
& {[1,2,4,6,7]} \\
& {[1,3,4,5,7]} \\
& {[1,3,4,6,7]}
\end{aligned}
$$

## Visiting and Touring

- Visit : A test path $p$ visits node $n$ if $n$ is in $p$

A test path $p$ visits edge $e$ if $e$ is in $p$

- Tour : A test path $p$ tours subpath $q$ if $q$ is a subpath of $p$

Test path [ 1, 2, 4, 5, 7]
Visits nodes ? 1, 2, 4, 5, 7
Visits edges ? $(1,2),(2,4),(4,5),(5,7)$
Tours subpaths? $[1,2,4],[2,4,5],[4,5,7],[1,2,4,5]$,

$$
[2,4,5,7],[1,2,4,5,7]
$$

(Also, each edge is technically a subpath)

## Tests and Test Paths

- path $(t)$ : The test path executed by test $t$
- path $(T)$ : The set of test paths executed by set of tests $T$
- Each test executes one and only one test path
- Complete execution from a start node to a final node

$$
\begin{aligned}
& \text { I.s the last bonlet true? } \\
& \text {-recursion }
\end{aligned}
$$

## Tests and Test Paths (2)

- A location in a graph (node or edge) can be reached from another location if there is a sequence of edges from the first location to the second
- Syntactic reach : A subpath exists in the graph
- Semantic reach : A test exists that can execute that subpath
- This distinction (semantic vs syntactic) is important when applied to source code


## Tests and Test Paths (3) <br> 

Test Path

## test 3

Deterministic software: a test always execute same test path


Non-deterministic software: a test can execute >I test paths

## Testing and Covering Graphs

- We use graphs in testing as follows :
- Develop a model of the software as a graph
- Require tests to visit/tour sets of nodes, edges or sub-paths


## Testing and Covering Graphs (2)

- Test Requirements (TR) : Describe properties of test paths
- Test Criterion : Rules that define test requirements
- Satisfaction : Given a set TR of test requirements for a criterion $C$, a set of tests $T$ satisfies $C$ on a graph if and only if for every test requirement tr in $T R$, there is a test path in path $(T)$ that meets the test requirement tr


## Two kinds of graph coverage criteria

I. Structural Coverage Criteria : Defined on a graph just in terms of nodes and edges
2. Data Flow Coverage Criteria : Requires a graph to be annotated with references to variables

## Node Coverage

- The first (and simplest) two criteria require that each node and edge in a graph be executed

Node Coverage (NC) : Test set T satisfies node coverage on graph G iff for every syntactically reachable node $\boldsymbol{n}$ in $N$, there is some path $p$ in path( $T$ ) such that $p$ visits $n$.

- This statement is a bit cumbersome, so we abbreviate it in terms of the set of test requirements

> Node Coverage (NC): TR contains each reachable node in G.

## Edge Coverage

- Edge coverage is slightly stronger than node coverage

Edge Coverage (EC) : TR contains each reachable path of ength up to Dinclusive, in $\mathbf{G}$.

- The phrase "length up to $l$ " allows for graphs with one node and no edges


## Node and Edge Coverage

- NC and EC are only different when there is an edge and another subpath between a pair of nodes (as in an "ifelse" statement)


Node Coverage : ? TR = \{ 1, 2, 3 \}
Test Path = [1, 2, 3]
Edge Coverage : ? TR = \{(1, 2), (1, 3), (2, 3) \} Test Paths = [1, 2, 3] $V$ $[1,3]$

## Paths of Length 1 and 0

- A graph with only one node will not have any edges
- It may seem trivial, but formally, Edge Coverage needs to require Node Coverage on this graph
- Else, Edge Coverage will not subsume Node Coverage
- So, we define "length up to 1 " instead of simply "length 1 "
- We have the same issue with graphs that only have one edge - for Edge-Pair Coverage ...



## Covering Multiple Edges

- Edge-pair coverage requires pairs of edges, or subpaths of length 2
Edge-Pair Coverage (EPC) : TR contains each reachable path of length up to 2 , inclusive, in $\mathbf{G}$.
- The phrase "length up to 2 " is used to include graphs that have less than 2 edges


$$
\begin{aligned}
& \text { Edge-Pair Coverage : ? } \\
& \text { TR }=\{[1,4,5],[1,4,6],[2,4,5] \text {, } \\
& [2,4,6],[3,4,5],[3,4,6]\}
\end{aligned}
$$

- A logical extension is to require covering all paths ...


## Covering Multiple Edges

## Complete Path Coverage (CPC): TR contains all paths in G.

Unfortunately, this is impossible if the graph has a loop, so a weak compromise makes the tester decide which paths:

Specified Path Coverage (SPC) : TR contains a set S of test paths, where S is supplied as a parameter.

## Structural Coverage Example



## Node Coverage

$$
T R=\{1,2,3,4,5,6,7\}
$$

$$
\text { Test Paths: }[1,2,3,4,7][1,2,3,5,6,5,7]
$$

## Edge Coverage

$\operatorname{TR}=\{(1,2),(1,3),(2,3),(3,4),(3,5),(4,7),(5,6),(5,7)$, $(6,5)\}$
Test Paths: [ 1, 2, 3, 4, 7 ] [1, 3, 5, 6, 5, 7 ]

## Edge-Pair Coverage

$\operatorname{TR}=\{[1,2,3],[1,3,4],[1,3,5],[2,3,4],[2,3,5],[3,4,7]$,

$$
[3,5,6],[3,5,7],[5,6,5],[6,5,6],[6,5,7]\}
$$

Test Paths: $[1,2,3,4,7][1,2,3,5,7][1,3,4,7]$ $[1,3,5,6,5,6,5,7][1,3,5,7]$

## Complete Path Coverage

Test Paths: [ 1, 2, 3, 4, 7 ] [ 1, 2, 3, 5, 7 ] [ 1, 2, 3, 5, 6, 5, 7 ] [ 1, 2, 3, 5, 6, 5, 6, 5, 7 ] [ 1, 2, 3, 5, 6, 5, 6, 5, 6, 5, 7 ] ...

## Handling Loops in Graphs

- If a graph contains a loop, it has an infinite number of paths
- Thus, Complete Path Coverage is not feasible
- SPC is not satisfactory because the results are subjective and vary with the tester
- Attempts to "deal with" loops:
- 1970s : Execute cycles once ([5, 6, 5] in previous example, informal)
- 1980s : Execute each loop, exactly once (formalized)
- 1990s: Execute loops 0 times, once, more than once (informal description)
- 2000s : Prime paths (touring, sidetrips, and detours)


## Simple Paths and Prime Paths

- Simple Path : A path from node $n_{i}$ to $n_{j}$ is simple if no node appears more than once, except possibly the first and last nodes are the same
- No internal loops
- A loop is a simple path

$$
\begin{aligned}
& {[2,5,6,5,7]^{x}} \\
& {[5,6,5]}
\end{aligned}
$$

- Prime Path : A simple path that does not appear as a proper subpath of any other simple path


Simple Paths : $[1,2,4,1],[1,3,4,1],[2,4,1,2],[2,4,1,3]$, $[3,4,1,2],[3,4,1,3],[4,1,2,4],[4,1,3,4],[1,2,4],[1,3,4]$, $[2,4,1],[3,4,1],[4,1,2],[4,1,3],[1,2],[1,3],[2,4],[3,4]$, [4,1], [1], [2], [3], [4]

Prime Paths : $[2,4,1,2],[2,4,1,3],[1,3,4,1],[1,2,4,1]$, [3,4,1,2], [4,1,3,4], [4,1,2,4], [3,4,1,3]

## Prime Path Coverage

- A simple, elegant and finite criterion that requires loops to be executed as well as skipped


## Prime Path Coverage (PPC) : TR contains each prime path in G.

- Will tour all paths of length $0,1, \ldots$
- That is, it subsumes node and edge coverage
- PPC almost, but not quite, subsumes EPC ...

Why does PPC not subsume EPC?

## PPC Does Not Subsume EPC

- If a node $n$ has an edge to itself (self edge), EPC requires $[n, n, m]$ and $[m, n, n]$
- Neither $[n, n, m]$ nor $[m, n, n]$ are simple paths (not prime)


> PPC Requirements : ?
> TR = \{ $[1,2,3],[2,2]\}$

## Prime Path Example

- The previous example has 38 simple paths
- Only nine prime paths



## Touring, Sidetrips, and Detours

- Prime paths have no internal loops ... test paths might
- Tour : A test path $p$ tours subpath $q$ if $q$ is $a$ subpath of $p$
- Tour With Sidetrips : A test path $p$ tours subpath $q$ with sidetrips iff every edge in $q$ is also in $p$ in the same order
- Tour can have a sidetrip if it comes back to the same node
- Tour With Detours : A test path $p$ tours subpath $q$ with detours iff every node in $q$ is also in $p$ in the same order
- Tour can have a detour from node $n_{i j}$ if it returns to the prime path at a successor of $n_{i}$


## Sidetrips and Detours Example



## Dealing with Infeasible TRs

- An infeasible test requirement cannot be satisfied
- Unreachable statement (dead code)
- Subpath that can only be toured if a contradiction holds, e.g., if $(x>0$ and $(x<0)$
- Most test criteria have some infeasible test requirements
- It is usually undecidable whether all test requirements are feasible


## Infeasible TRs and Sidetrips

- When sidetrips are not allowed, many structural criteria have more infeasible test requirements
- However, always allowing sidetrips weakens the test criteria


## Practical recommendation-Best Effort Touring

 - First, satisfy as many test requirements as possible without sidetrips- Then, allow sidetrips to try to satisfy remaining test requirements


## Simple path \& prime path example

"!" Means "cannot be extended to a simple path"


## Round Trips

- Round-Trip Path : A prime path that starts and ends at the same node

Simple Round Trip Coverage (SRTC) : TR contains at least one round-trip path for each reachable node in G that begins and ends a round-trip path.

Complete Round Trip Coverage (CRTC) : TR contains all round-trip paths for each reachable node in $\mathbf{G}$.

- The criteria omit nodes \& edges that are not in round trips
- They do not subsume edge-pair, edge, or node coverage


## Data Flow Criteria

## Goal : Ensure that values are computed and used correctly

- Definition (def) : A location where a value for a variable is stored into memory
- Use : A location where a variable's value is accessed


$$
\begin{aligned}
& \text { Defs: } \operatorname{def}(1)=\{\mathbf{X}\} \\
& \text { def }(5)=\{\mathbb{Z}\} \begin{array}{l}
\text { Fill in } \\
\text { these }
\end{array} \\
& \operatorname{def}(6)=\{\mathbb{Z}\} \text { sets } \\
& \text { Uses: use (5) }=\{\mathbf{X}\} \\
& \text { use (6) }=\{\mathbf{X}\}
\end{aligned}
$$

The values given in defs should reach at least one, some, or all possible uses

## DU Pairs and DU Paths

- def $(\mathrm{n})$ or def $(\mathrm{e})$ :The set of variables that are defined by node n or edge e
- use ( n ) or use (e) : The set of variables that are used by node n or edge e
- DU pair: A pair of locations $\left(l_{i}, l_{j}\right)$ such that a variable $v$ is defined at $l_{i}$ and used at $l_{j}$
- Def-clear : A path from $I_{i}$ to $l_{j}$ is def-clear with respect to variable $v$ if $v$ is not given another value on any of the nodes or edges in the path
- Reach : If there is a def-clear path from $I_{i}$ to $I_{j}$ with respect to $v$, the def of $v$ at $I_{i}$ reaches the use at $l_{j}$
- du-path :A simple subpath that is def-clear with respect to $v$ from a def of $v$ to a use of $v$
- du $\left(n_{i}, n_{j}, v\right)$ - the set of du-paths from $n_{i}$ to $n_{j}$
- du $\left(n_{i}, v\right)$ - the set of du-paths that start at $n_{i}$


## Touring DU-Paths

- A test path $p$ du-tours subpath $d$ with respect to $v$ if $p$ tours $d$ and $d$ is def-clear with respect to $v$
- Sidetrips can be used, just as with previous touring
- Three criteria
- Use every def
- Get to every use
- Follow all du-paths


## Data Flow Test Criteria

- First, we make sure every def reaches a use

All-defs coverage (ADC) : For each set of du-paths $S=d u$ $(n, v)$, TR contains at least one path $d$ in $S$.

- Then we make sure that every def reaches all possible uses

All-uses coverage (AUC) : For each set of du-paths to uses $S=d u\left(n_{i}, n_{j}, v\right)$,TR contains at least one path $d$ in $S$.

- Finally, we cover all the paths between defs and uses

All-du-paths coverage (ADUPC) : For each set $S=d u$ (ni, $n j, v)$, TR contains every path $d$ in $S$.

## Data Flow Testing Example



All-defs for $X$ [1, 2, 4, 5 ]

All-uses for $X$
[ $1,2,4,5$ ]
[1, 2, 4, 6 ]

All-du-paths for $X$
[1, 2, 4, 5 ]
$[1,3,4,5]$
[ 1, 2, 4, 6 ]
[1, 3, 4, 6 ]

## Graph-based criteria subsumption



## Summary

- Graphs are a very powerful abstraction for designing tests
- The various criteria allow lots of cost / benefit tradeoffs
- These two sections are entirely at the "design abstraction level" from chapter 2
- Graphs appear in many situations in software
- Next: we will apply these criteria to source code
- Design, specs, and use cases are not covered in CS 5I54

