CSCI-GA.3033.003 Scripting Languages

12/02/2013 OCaml

Acknowledgement

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About OCaml

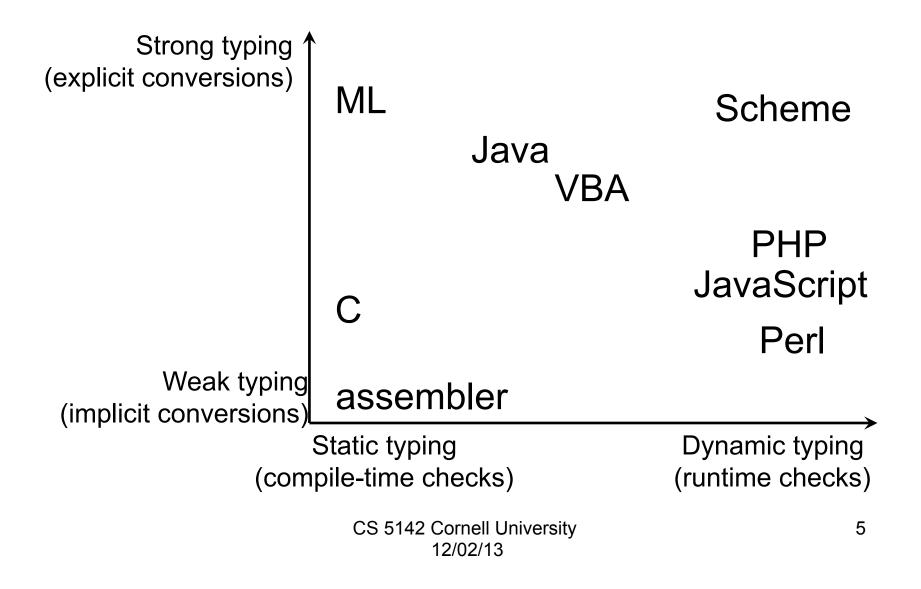
- A functional programming language
 - All computation is done with functions
 - No state (mutable variables), simple and clean
 - Functions are first class: you can pass them, return them, etc.
 - Support for object-oriented programming (the O in OCaml)
- Statically typed, type-safe language
 - You can't apply the wrong operations to the wrong data (e.g., can't try to divide two strings)
 - Avoids many "silly errors" and provides security (e.g., no buffer overflows)

Ocaml Features

- Garbage collection: automatic memory management
- Type inference: You don't have to write the type information down. The compiler will figure it out.
- Parametric polymorphism: Write functions and data structures that can be used with any type. Note that Java provides subtype polymorphism.
- Algebraic data types: Can build data structures easily.
 Pattern matching makes working with them convenient.
- Advanced modules: Encapsulate implementations behind interfaces. Functors (functions that manipulate modules) add functionality beyond most languages' module systems.

Concepts

Weak/Strong, Static/Dynamic Typing



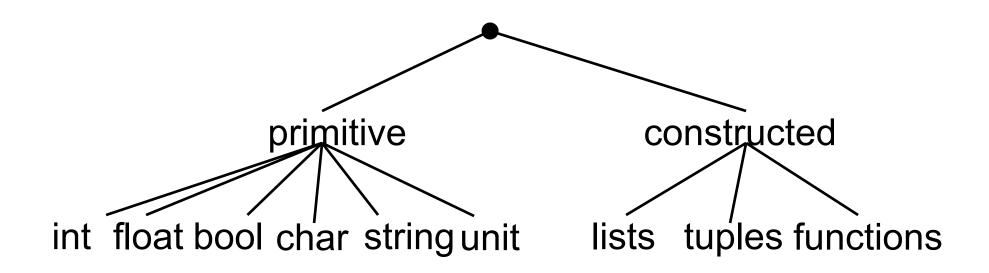
History

- Robin Milner and others at the University of Edinburgh were working on theorem provers
- Problem: the provers would sometimes put incorrect "proofs" together, and claim they were valid
- Designed ML (Meta Language) as a language to let you construct valid proofs
- In the early '80s, there was a schism in the ML community
 - French developed Caml and later Objective Caml (OCaml)
 - British and Americans developed Standard ML (SML)

How to Write + Run Code

- ocaml
 - REPL
- ocaml file.ml
 - Run the interpreter
- ocamlc file.ml; ./a.out
 - Run the compiler

Types



Type Inference

The compiler deduces the type of an expression

$$y = x + 1$$

 Compiler knows that + takes two integers and returns an integer, so x and y should be integers.

$$z = x + 1.0$$

- This should produce an error, since x is being used as both an int and a float
- Compiler infers types by aggregating type information, and reducing expressions to implicitly typed values

Type Inference

- Statically determining whether a program will have a type-error is impossible
- All statically typed languages are conservative, and may reject some programs that are perfectly okay
- Milner formulated the type-inference system for ML, and proved its soundness
- Note that Milner also worked on concurrent programming languages (CCS, CSP, pi-Calculus), and won the Touring award – in large part to his work on ML

Syntax (subset)

Syntactic class	Grammar rule	Example
identifiers	x,f	a, x,y,x_y, z100
constants	С	1, 01, 1.0, true, "hello" 'A'
unary operator	и	-, not
binary operator	b	+,*,-, <,>, ^, !=,
terms	e::= x c u e e b e if e then e else e let d and d in e e (e, ,e)	not b, 2 + 2, foo
declarations	$t := x = e \mid f(x,,x) : t = e$	one = 1

Program Errors

- An expression can have a legal syntax, but may be wrong. The expression must be well-typed.
- Ways that a program can be wrong:
 - Syntax errors: let 0 x =
 - Type errors: "x" + 3
 - Semantic errors: 1/0
 - More general errors: computes the wrong output

Example

```
let abs (x : int) : int =
if x < 0 then -x else x
```

```
let abs : int -> int =
function x -> if x < 0 then -x else x
```

```
let abs = fun x \rightarrow fun x < 0 then -x else x
```

```
val abs : int -> int = <fun>
```

- Every expression and declaration has both a type and a value.
- Here, we have bound the name abs to a function whose type is int -> int.

Type Checking

```
$ ocaml
# let abs = fun x -> if x < 0 then -x else x ;;
val abs : int -> int = <fun>

# abs 3;;
- : int = 3

# abs "foo";;
Error: This expression has type string but an expression was expected of type
    int
```

Scope

Variable declarations bind variables within a scope

let x = e1 in e2

- The scope of x is the expression e2.
- Functions also bind variables

let f x = e1 in e2

The scope of the formal parameter x is the expression in e1. The scope of the variable f (which is bound to a function value) is the body of the let, e2.

Scope

A let expression can introduce multiple variables

```
let x = 2 and y = 3 in x + y
```

• To declare a recursive function, the function must be in scope in its own body. In Ocaml, you use let rec instead of let for this.

```
let rec even x = x = 0 || odd (x-1)
and odd x = not (x = 0 || not (even (x-1)))
in
odd 3110
```

Curried functions

 A function with multiple parameters is really just syntactic sugar for a function passed a tuple as an argument

```
let plus (x, y) = x + y;
val plus : int * int -> int = <fun>
plus (2, 3)
```

 Ocaml has another way to declare functions with multiple arguments. In curried form:

let plus
$$x y = x + y$$

Notice there are no commas between the parameters

Curried functions

There are also no commas when calling the function:

plus 23

 Functions really only have one argument. The plus function is being passed one argument, 2, and it returns a function that takes one argument. (plus 2) (3)

```
plus 2 3
= ((function (x : int) -> function (y : int) -> x + y) 2) 3
= (function (y : int) -> 2 + y) 3
= 2 + 3
= 5
```

Lists

- Lists are **immutable**. You cannot change the elements of a list.
- Lists are homogenous. The can only contain one type.
- Examples: [1;2;3], 1::[2;3], [1;2]@[3;4]

Pattern Matching

Use match expressions to extract elements from a list

match 1st with

 Returns 0 if the list is empty, 1 if the list has 1 element, and 2 if it has 2 or more elements

Pattern Matching

 Use recursive functions to do something to every element in a list

```
let rec length (lst : string list) : int =
  match lst with
  [] -> 0
  | h :: t -> 1 + length t
```

This function computes the length of the list lst.

Variant Types

- Allows you to have variables that contain more than one kind of value
- Similar to an enum in java, or a tagged union in C

type answer = Yes | No | Maybe;;

- Type keyword declared name for the type.
- Declared with a set of constructors

```
type eitherPoint = TwoD of float * float
| ThreeD of float * float * float
```

Pattern Matching

- Ocaml patterns are very rich
- Can match on variant types, tuples, records, and can pull the values apart

```
let answer_to_string (a : answer) : string =
  match a with
  | Yes -> "yes"
  | No -> "no"
```

Polymorphism

 Suppose we want to write a function that swaps positions of a pair:

```
let swapInt ((x : int), (y : int)) : int * int = (y, x)
and swapReal ((x : float), (y : float)) : float * float = (y, x)
and swapString ((x : string), (y : string)) : string * string = (y, x)
```

Better way:

```
let swap ((x : 'a), (y : 'b)) : 'b * 'a = (y, x);;
```

Or without type declarations:

```
# let swap (x, y) = (y, x);;
val swap : 'a * 'b -> 'b * 'a = <fun>
```

Mapping

 Apply a function to every element in a list, and return a new list:

List.map f[a; b; c] = [f a; f b; f c]

Copy a list (using an anonymous function)

let copy = map (fun $x \rightarrow x$)

Folding

 Apply a function to every element in a list, accumulates a result:

```
fold_right f [a; b; c] r = f a (f b (f c r))
fold_left f r [a; b; c] = f (f (f r a) b) c
```

Sum all the elements in a list:

```
let sum_right_to_left il = fold_right (+) il 0
let sum_left_to_right = fold_left (+) 0
```

Folding

- Fold is very powerful!
- Can write many other list functions in terms of fold

```
type intlist = Nil | Cons of (int * intlist)

let mapp f lst = fold_right (fun x y -> Cons (f x, y)) lst Nil
```

Note that fold_left would give the result in reverse order

```
let length = fold_left (fun x = -> 1 + x) 0
```

Select a subset of the list

```
let filter f lst = fold_right (fun x y -> if f x then Cons (x, y) else y) lst Nil
```



Last Slide

Next lecture: Review!