# Machine Learning for Data Science (CS4786) Lecture 6 

Random Projections

Feb 09, 2015

Course Webpage :
http://www.cs.cornell.edu/Courses/cs4786/2015sp/

CCA Demo Redo?




## Back to Single View: Recap




## The Tall, the Fat and the Ugly



## The Tall, the Fat and the Ugly



K

## The Tall, the Fat and the Ugly



## THE TALL, the Fat AND THE UGLY

## $n$ <br>  <br> $d$

 $\operatorname{SVD}(X)$$n$

$>$

## $W^{\top}$

## THE TALL, the Fat AND THE UGLY

## $n$ <br> 

$n$

$n$


## THE TALL, THE FAT AND the Ugly



- $d$ and $n$ so large we can't even store in memory
- Only have time to be linear in $\operatorname{size}(X)=n \times d$

I there any hope?

## PICK A Random W

$$
Y=X \times\left[\begin{array}{ccc}
+1 & \ldots & -1 \\
-1 & \ldots & +1 \\
+1 & \ldots & -1 \\
& \cdot & \\
& \cdot & \\
+1 & \ldots & -1
\end{array}\right] d / \sqrt{K}
$$

## Random Projection

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Distances between all pairs of data-points in low dim. projection is roughly the same as their distances in the high dim. space.

That is, when $K$ is "large enough", with "high probability", for all pairs of data points $i, j \in\{1, \ldots, n\}$,

$$
(1-\epsilon)\left\|\mathbf{y}_{i}-\mathbf{y}_{j}\right\|_{2} \leq\left\|\mathbf{x}_{i}-\mathbf{x}_{j}\right\|_{2} \leq(1+\epsilon)\left\|\mathbf{y}_{i}-\mathbf{y}_{j}\right\|_{2}
$$

## WHY ShOULD RANDOM PROJECTIONS EVEN WORK?!

Consider any vector $\tilde{\mathbf{x}} \in \mathbb{R}^{d}$ and let $\tilde{\mathbf{y}}=W^{\top} \tilde{\mathbf{x}}$. Note that

$$
\tilde{\mathbf{y}}[j]^{2}=\left(\sum_{i=1}^{d} W[i, j] \cdot \tilde{\mathbf{x}}[i]\right)^{2}
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& =\sum_{i, i^{\prime}}\left(W[i, j] \cdot W\left[i^{\prime}, j\right]\right) \cdot\left(\tilde{\mathbf{x}}[i] \cdot \tilde{\mathbf{x}}\left[i^{\prime}\right]\right)
\end{aligned}
$$

## Why should Random Projections even work?!

Hence,
$\mathbb{E}\left[\tilde{\mathbf{y}}[j]^{2}\right]=\sum_{i, i^{\prime}=1}^{d} \mathbb{E}\left[\left(W[i, j] \cdot W\left[i^{\prime}, j\right]\right)\right] \cdot\left(\tilde{\mathbf{x}}[i] \cdot \tilde{\mathbf{x}}\left[i^{\prime}\right]\right)$

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if $i \neq i^{\prime}, W[i, j]$ and $W\left[i^{\prime}, j\right]$ are independent and so

$$
=\sum_{i=1}^{d} \mathbb{E}\left[\left(W[i, j]^{2}\right)\right] \tilde{\mathbf{x}}[i]^{2}+\sum_{i \neq i^{\prime}}\left(\mathbb{E}[W[i, j]] \cdot \mathbb{E}\left[W\left[i^{\prime}, j\right]\right]\right) \cdot\left(\tilde{\mathbf{x}}[i] \cdot \tilde{\mathbf{x}}\left[i^{\prime}\right]\right)
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& =\sum_{i=1}^{d} \tilde{\mathbf{x}}[i]^{2} / \sqrt{K}^{2}=\|\tilde{\mathbf{x}}\|_{2}^{2} / K
\end{aligned}
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Lets try this in Matlab ...

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This is called the Johnson-Lindenstrauss lemma or JL lemma for short.

WHY is THis so Ridiculously Magical?
$\mathrm{n}=$ 1000

$$
d=1000
$$

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If we take $K=69.1 / \epsilon^{2}$, with probability 0.99 distances are preserved to accuracy $\epsilon$

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n= 1000

$$
d=1000000
$$

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