Adversarial Search

## Game Playing

## An AI Favorite

- structured task, often a symbol of "intelligence"
- clear definition of success and failure
- does not require large amounts of knowledge (at first glance)
- focus on games of perfect information
- multiplayer, chance


## Game Playing

Initial State is the initial board/position
Successor Function defines the set of legal moves from any position
Terminal Test determines when the game is over
Utility Function gives a numeric outcome for the game

For chess, only win, lose, draw. Backgammon: +192 to -192.

## Partial Search Tree for Tic-Tac-Toe



## Game Playing as Search



Two Ply

## Simplified Minimax Algorithm

1. Expand the entire tree below the root.
2. Evaluate the terminal nodes as wins for the minimizer or maximizer (i.e. utility).
3. Select an unlabeled node, $n$, all of whose children have been assigned values. If there is no such node, we're done --- return the value assigned to the root.
4. If $n$ is a minimizer move, assign it a value that is the minimum of the values of its children. If $n$ is a maximizer move, assign it a value that is the maximum of the values of its children. Return to Step 3.

## Another Example



According to minimax, which action to take? $A=A 1 \quad B=A 2 \quad C=A 3$

## Another Example



## Minimax

function MINIMAX-DECISION(game) returns an operator for each op in OPERATORS[game]do VALUE[op]<MINIMAX-VALUE(APPLY(op,game),game)
end
return the op with the highest VALUE[op]
function MINIMAX-VALUE(state,game) returns a utility value
if TERMINAL-TEST[game](state) then
return UTILITY[game](state)
else if MAX is to move in state then
return the highest MINIMAX-VALUE of SUCCESSORS(state)
else
return the lowest MINIMAX-VALUE of SUCCESSORS(state)

# Improving Minimax: $\alpha-\boldsymbol{\beta}$ Pruning 

Idea: Avoid generating the whole search tree
Approach: Analyze which subtrees have no influence on the solution

## $\boldsymbol{\alpha}-\boldsymbol{\beta}$ Search

$\alpha=$ best choice (highest) found so far for max, initially $-\infty$
$\beta=$ best choice (lowest) found so far for min , initially $+\infty$

We'll call $\alpha-\beta$ procedure recursively with a narrowing range between $\alpha$ and $\boldsymbol{\beta}$.

Maximizing levels may reset $\alpha$ to a higher value; Minimizing levels may reset $\boldsymbol{\beta}$ to a lower value.

## Features of Evolution

Player

Opponent

Player

Opponent


If $m$ is better than $n$ for Player, never get to $n$ in play.

## $\boldsymbol{\alpha}-\boldsymbol{\beta}$ Search Algorithm

1. If terminal state, compute $\mathrm{e}(\mathrm{n})$ and return the result.
2. Otherwise, if the level is a minimizing level,

- Until no more children or $\beta \leq \alpha$
- $v_{i} \leftarrow \alpha-\beta$ search on a child
- If $v_{i}<\beta, \beta \leftarrow v_{i}$.
- Return $\min \left(\boldsymbol{v}_{i}\right)$

3. Otherwise, the level is a maximizing level:

- Until no more children or $\alpha \geq \beta$,
$-v_{i} \leftarrow \alpha-\beta$ search on a child.
- If $v_{i}>\alpha$, set $\alpha \leftarrow v_{i}$
- Return $\max \left(v_{i}\right)$


## Search Space Size Reductions

Worst Case: In an ordering where worst options evaluated first, all nodes must be examined.

Best Case: If nodes ordered so that the best options are evaluated first, then what?


## The Need for Imperfect Decisions

Problem: Minimax assumes the program has time to search to the terminal nodes.

Solution: Cut off search earlier and apply a heuristic evaluation function to the leaves.

## Static Evaluation Functions

Minimax depends on the translation of board quality into single, summarizing number. Difficult. Expensive.

- Add up values of pieces each player has (weighted by importance of piece).
- Isolated pawns are bad.
- How well protected is your king?
- How much maneuverability to you have?
- Do you control the center of the board?
- Strategies change as the game proceeds.


## Design Issues for Heuristic Minimax

Evaluation Function:

Need to be carefully crafted and depends on game! What criteria should an evaluation function fulfill?

## Linear Evaluation Functions

$$
\boldsymbol{w}_{1} f_{1}+\boldsymbol{w}_{2} \boldsymbol{f}_{2}+\ldots+\boldsymbol{w}_{\boldsymbol{n}} \boldsymbol{f}_{\boldsymbol{n}}
$$

- This is what most game playing programs use
- Steps in designing an evaluation function:

1. Pick informative features.
2. Find the weights that make the program play well

## Design Issues for Heuristics Minimax

Search: search to a constant depth
What are problems with constant search depth?

## Backgammon Board

0

$242322212019 \quad 181716151413$

## Backgammon - Rules

- Goal: move all of your pieces off the board before your opponent does.
- Black moves counterclockwise toward 0.
- White moves clockwise toward 25.
- A piece can move to any position except one where there are two or more of the opponent's pieces.
- If it moves to a position with one opponent piece, that piece is captured and has to start it's journey from the beginning.


## Backgammon - Rules

- If you roll doubles you take 4 moves (example: roll 5,5, make moves 5,5,5,5).
- Moves can be made by one or two pieces (in the case of doubles by 1, 2, 3 or 4 pieces)
- And a few other rules that concern bearing off and forced moves.


White has rolled 6-5 and has 4 legal moves: (5-10,5-11), (5-11,19-24), (5-10,10-16) and (5-11,11-16).

## Game Tree for Backgammon

MAX

DICE

MIN


TERMINAL

## Expectiminimax

$\operatorname{Expectiminimax}(\mathrm{n})=$
Utility(n)
for $n$, a terminal state
$\max _{s \in S u c c(n)}$ expectiminimax $(s)$ for $\mathbf{n}$, a Max node $\min _{s \in \operatorname{Succ}(n)} \operatorname{expectiminimax}(s)$ for $\mathbf{n}$, a Min node $\Sigma_{s \in \operatorname{Succ}(\boldsymbol{n})} \boldsymbol{P}(\boldsymbol{s}) * \operatorname{expectiminimax}(\boldsymbol{s})$ for n , a chance node

## Evaluation function



## State of the Art in Backgammon

- 1980: BKG using two-ply (depth 2) search and lots of luck defeated the human world champion.
- 1992: Tesauro combines Samuel's learning method with neural networks to develop a new evaluation function (search depth 2-3), resulting in a program ranked among the top 3 players in the world.


## State ofthe Artincheckers

- 1952: Samuel developed a checkers program that learned its own evaluation function through self play.
- 1990:

Chinook (J. Schaeffer) wins the U.S. Open. At the world championship, Marion Tinsley beat Chinook.

- 2005: Schaeffer et al. solved checkers for "White Doctor" opening (draw) (about 50 other openings).


## State of the Art in Go

Large branching factor makes regular search methods inappropriate.

## Best computer Go programs ranked only "weak amateur'.

Employ pattern recognition techniques and limited search.
$\$ 2,000,000$ prize available for first computer program to defeat a top level player.

## History of Chess in AI

| 500 | Legal chess |
| :--- | :--- |
| 1200 | Occasional player |
| 2000 | World-ranked |
| 2900 | Gary Kasparov |

Early 1950's Shannon and Turing both had programs that (barely) played legal chess (500 rank).

1950's Alex Bernstein's system, $(500+\varepsilon)$
1957 Herb Simon claims that a computer chess program would be world chess champion in 10 years...yeah, right.

1966 McCarthy arranges computer chess match, Stanford vs. Russia. Long, drawn-out match. Russia wins.

1967 Richard Greenblatt, MIT. First of the modern chess programs, MacHack (1100 rating).

1968 McCarthy, Michie, Papert bet Levy (rated 2325) that a computer program would beat him within 10 years.

1970 ACM started running chess tournaments. Chess 3.0-6 (rated 1400).

1973 By 1973...Slate: "It had become too painful even to look at Chess 3.6 any more, let alone work on it."

1973 Chess 4.0: smart plausible-move generator rather than speeding up the search. Improved rapidly when put on faster machines.

1976 Chess 4.5: ranking of 2070.

1977 Chess 4.5 vs. $\sim$ Levy. Levy wins.

1980's Programs depend on search speed rather than knowledge (2300 range).
1993 DEEP THOUGHT: Sophisticated special-purpose computer; $\boldsymbol{\alpha}-\boldsymbol{\beta}$ search; searches 10-ply; singular extensions; rated about 2600.

1995 DEEP BLUE: searches 14-ply; iterative deepening $\boldsymbol{\alpha}-\boldsymbol{\beta}$ search; considers 100-200 billion positions per move; regularly reaches depth 14; evaluation function has $8000+$ features; singular extensions to 40-ply; opening book of 4000 positions; end-game database for 5-6 pieces.

1997 DEEP BLUE: first match won against world-champion (Kasparov). 2002 IBM declines re-match. FRITZ played world champion Vladimir Kramnik. 8 games. Ended in a draw.

## Concludes "Search"

Uninformed search: DFS / BFS / Uniform cost search
time / space complexity
size search space: up to approx. $10^{11}$ nodes
special case: Constraint Satisfaction / CSPs generic framework: variables \& constraints backtrack search (DFS); propagation (forward-checking / arc-consistency, variable / value ordering

Informed Search: use heuristic function guide to goal
Greedy best-first search
A* search / provably optimal
Search space up to approximately $10^{25}$

## Local search

Greedy / Hillclimbing
Simulated annealing
Tabu search
Genetic Algorithms / Genetic Programming
search space $10^{100}$ to $10^{1000}$

## Aversarial Search / Game Playing

minimax Up to $\sim 10^{10}$ nodes, $6-7$ ply in chess. alpha-beta pruning Up to ${ }^{\sim} 10^{20}$ nodes, 14 ply in
chess. provably optimal

## Search and AI

## Why such a central role?

- Basically, because lots of tasks in Al are intractable. Search is "only" way to handle them.
- Many applications of search, in e.g., Learning / Reasoning / Planning / NLU / Vision
- Good thing: much recent progress ( $10^{30}$ quite feasible; sometimes up to $10^{1000}$ ).

Qualitative difference from only a few years ago!

