CS4450

Computer Networks: Architecture and Protocols

Lecture 2
Sharing Networks:
"Circuits" and "Packets"

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Announcements

- The webpage is up!
 - https://www.cs.cornell.edu/courses/cs4450/2022fa/
 - Please read everything on the webpage carefully
 - Especially, Admin page
 - All slides, problem sets, readings, etc. will be on the webpage
 - Solutions etc. will be on Ed Discussions
- You should all be now on Ed Discussions
 - For those of you, who enrolled recently: give it a day or two
 - If you are unable to access by next week, send us an email: cs4450-staff@cornell.edu
- Tentative office hours on the webpage
 - Office hours start tomorrow
 - Please fill out the poll (we sent over email this afternoon)
 - We will announce the final office hours soon
- I do not expect you to read notes/slides before lecture

Announcements

- Communication with staff in 4450
 - All enrollment-related questions: courses@cis.cornell.edu
 - Everything: First check the webpage
 - Everything that is not answered on the webpage: Ed discussions
 - Time-sensitive: cs4450-staff@cornell.edu
 - Sensitive: <u>cs445-prof@cornell.edu</u>
 - Do not expect answers if you do not follow the above protocol
- Please inform us about any exam conflict before 09/07
- If you have sent us an email about exam conflict:
 - Please wait until 09/07
 - Once we know everyone's conflicts, we will find solutions.

Goal of Today's Lecture

- Learn about:
 - Two important performance metrics:
 - Bandwidth
 - Delay, or latency
 - Why are these important?
 - Two ways of sharing networks:
 - Circuit switching
 - Packet switching
 - Why do current computer networks use packet switching?

But first, Recap from last lecture

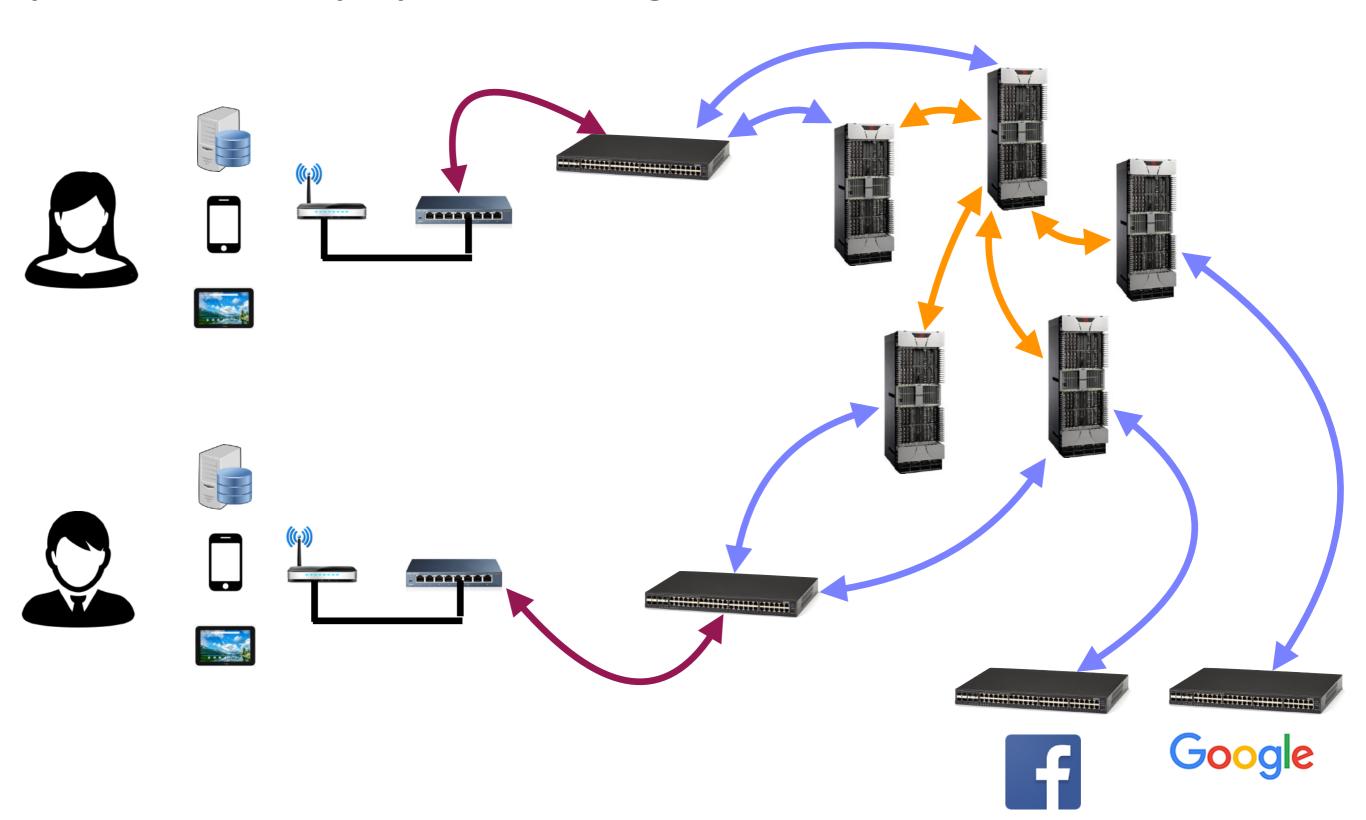
Recap: what is a computer network?

A set of network elements connected together, that implement a set of protocols for the purpose of sharing resources at the end hosts

- Three important components:
 - Core infrastructure:
 - A set of network elements connected together
 - Protocols:
 - Needed to use the network
 - Purpose:
 - Sharing resources at the end hosts (computing devices)

What is a computer network?

A set of network elements connected together, that implement a set of protocols for the purpose of sharing resources at the end hosts



Recap: what do computer networks do?

A computer network delivers data between the end points

- One and only one task: Delivering the data
- This delivery is done by:
 - Chopping the data into packets
 - Sending individual packets across the network
 - Reconstructing the data at the end points

Recap: what do computer networks look like?

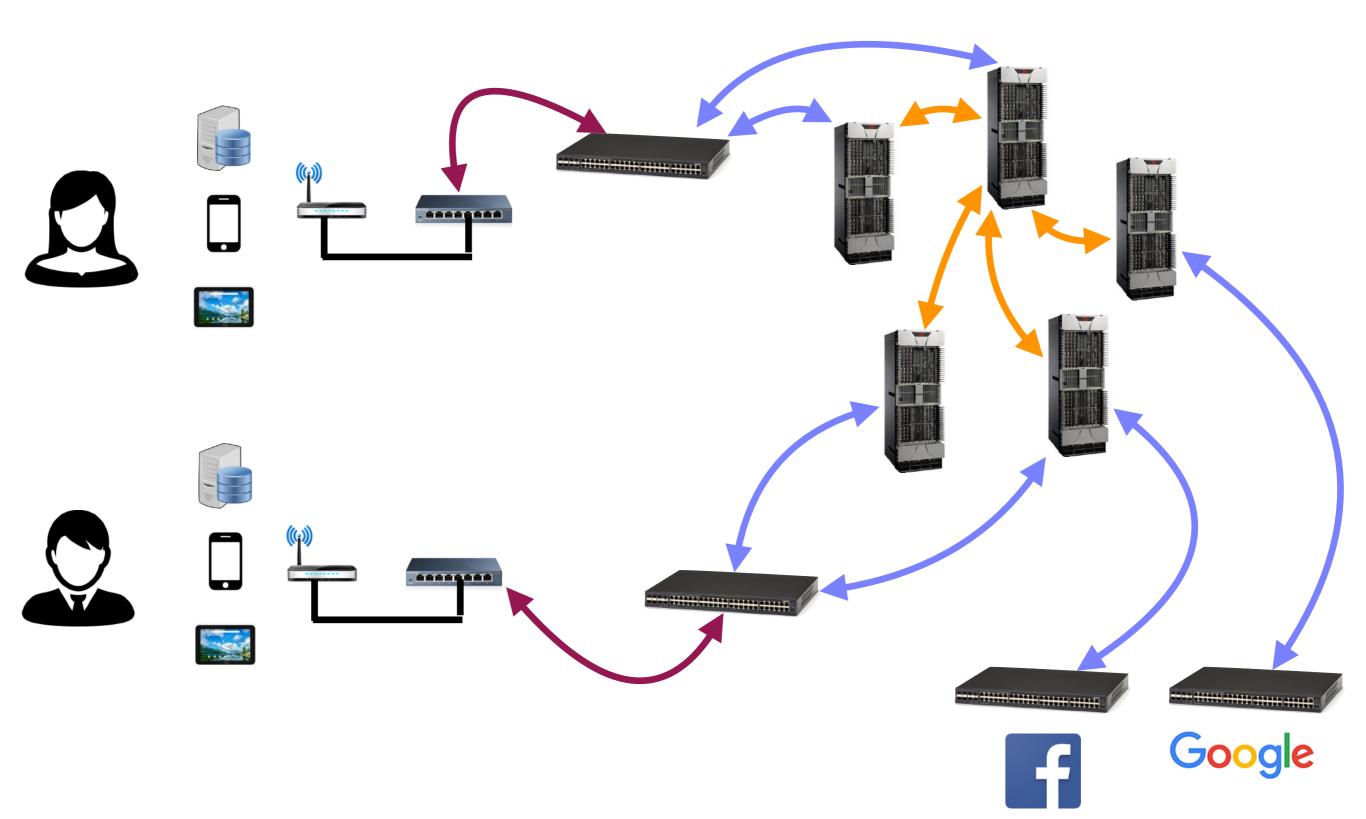
Three Basic components

- End hosts: they send/receive packets
- Switches/Routers: they forward packets
- Links: connect end hosts to switches, and switches to each other



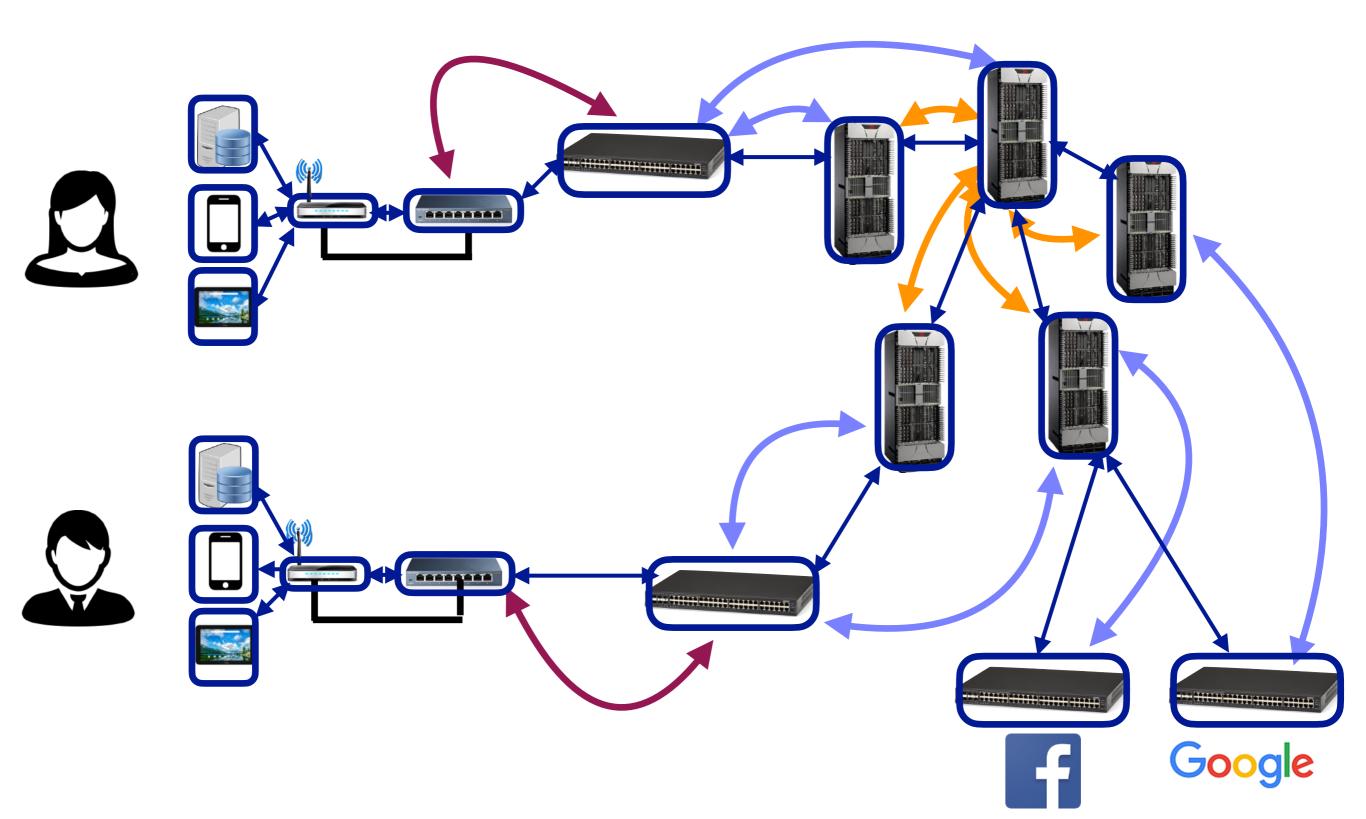
What is a computer network?

A set of network elements connected together, that implement a set of protocols for the purpose of sharing resources at the end hosts

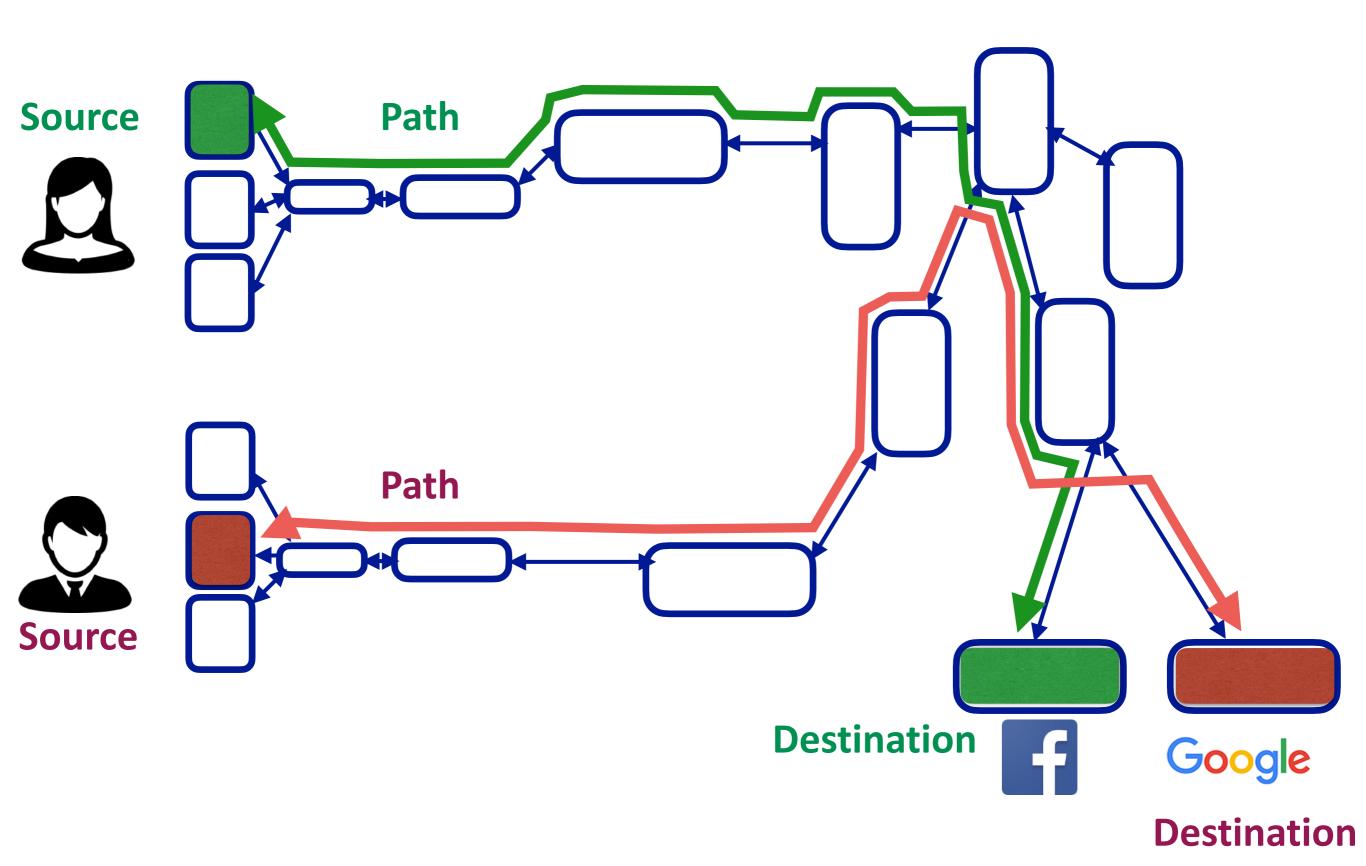


What is a computer network?

A set of network elements connected together, that implement a set of protocols for the purpose of sharing resources at the end hosts



A computer network can be abstractly represented as a graph

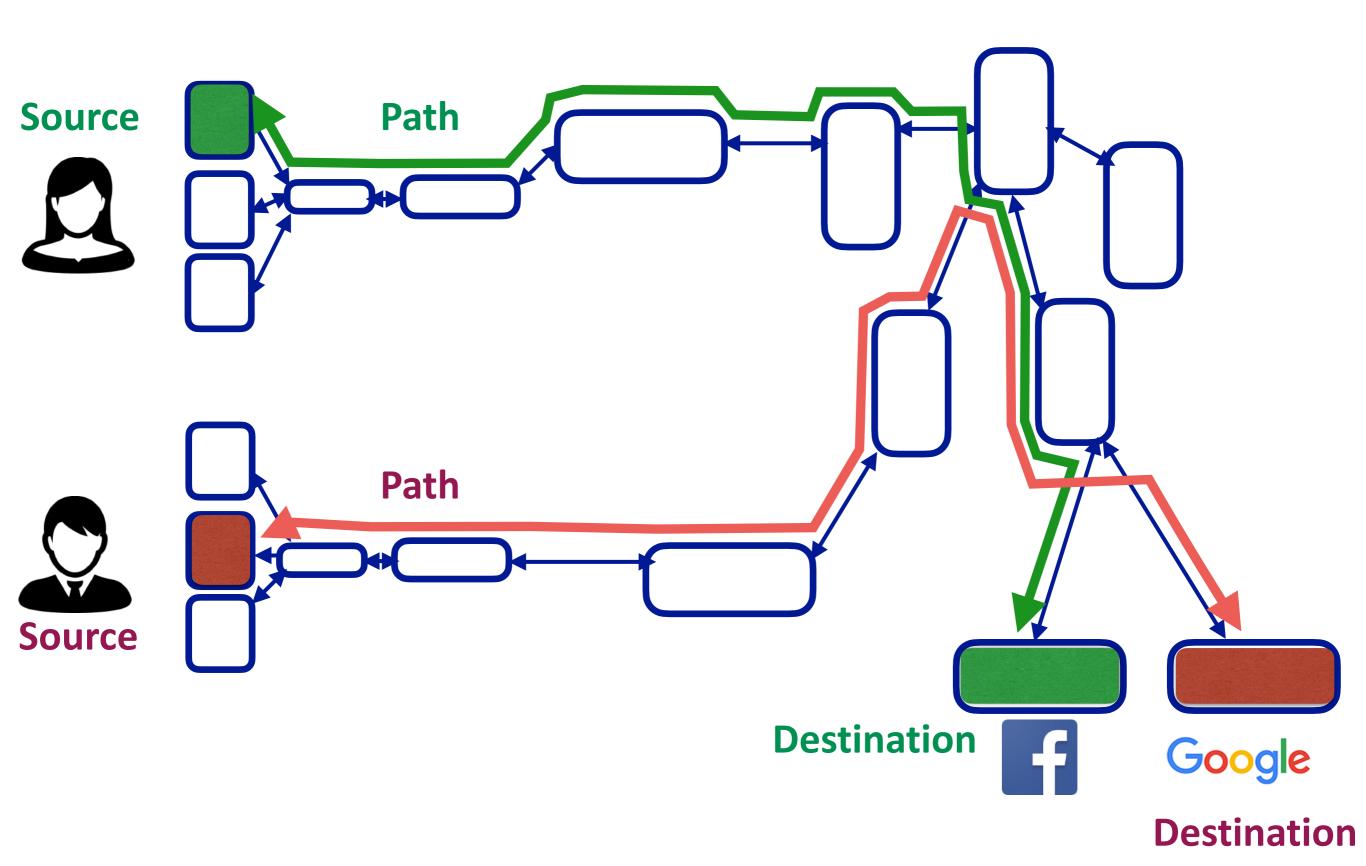


Many mechanisms underneath!

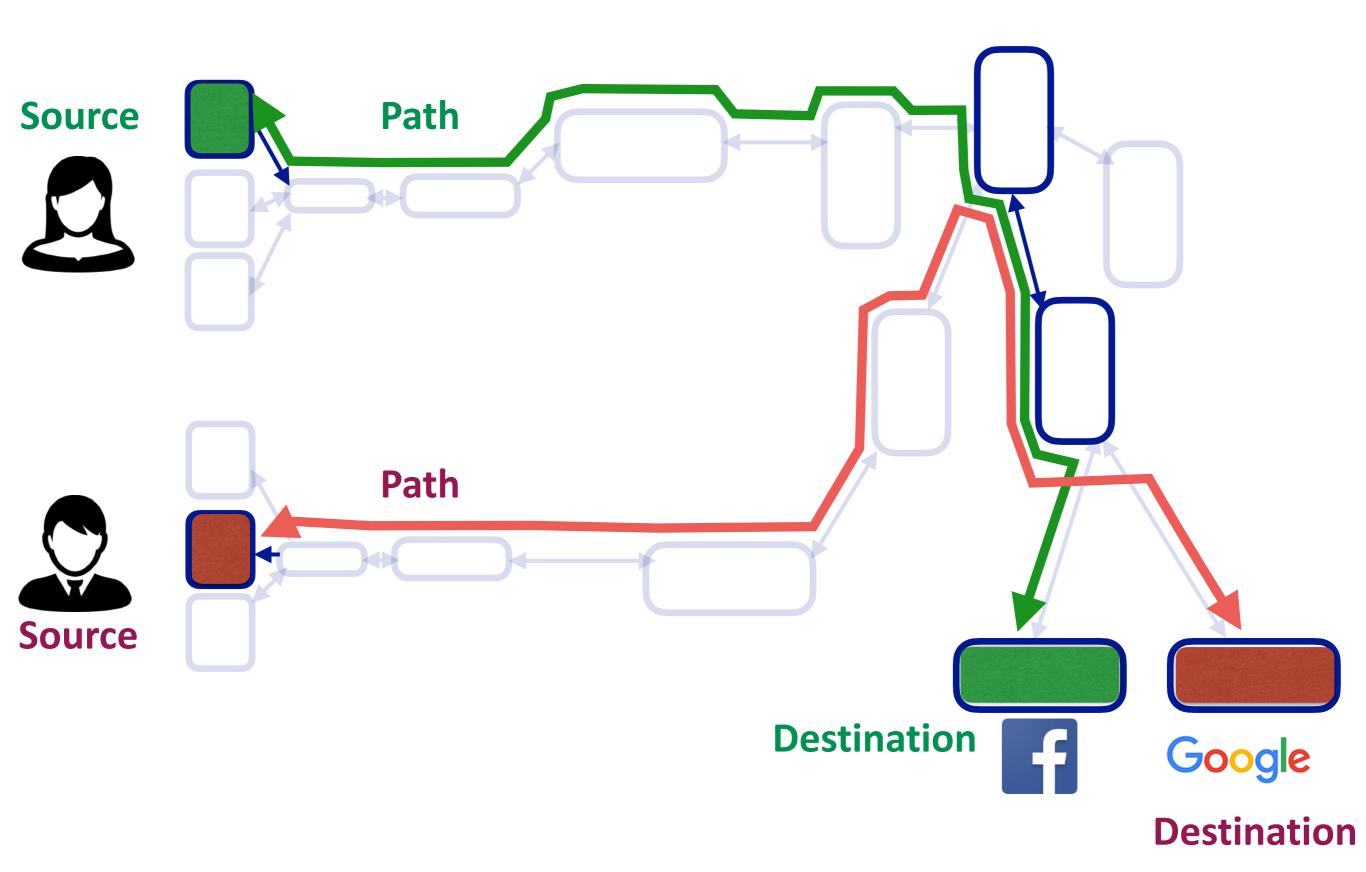
- Locating the destination: Naming, addressing
- Finding a path to the destination: Routing
- Sending data to the destination: Forwarding
- Failures, reliability, etc.: Distributed routing and congestion control

Will take the entire course to learn these

A computer network can be abstractly represented as a graph



Today's focus: sharing the network (graph)



Today's lecture: sharing computer networks

- 1. What does network sharing mean?
- 2. What are the performance metrics?
- 3. What are the various mechanisms for sharing networks?
- 4. Why "packets" and "flows"?
- 5. Understanding bandwidth and latency for packets

What does network sharing mean?

The problem of sharing networks

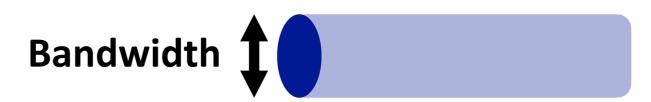
- Must support many "users" and "applications" at the same time
- Each user/application wants to use the network (send and receive data)
- Limited resources
 - We will learn, over the semester, that network has different resources.
- Fundamental question:
 - How does the network decide which resource to allocate to which user/application at any given point of time?

Resources relate to performance.

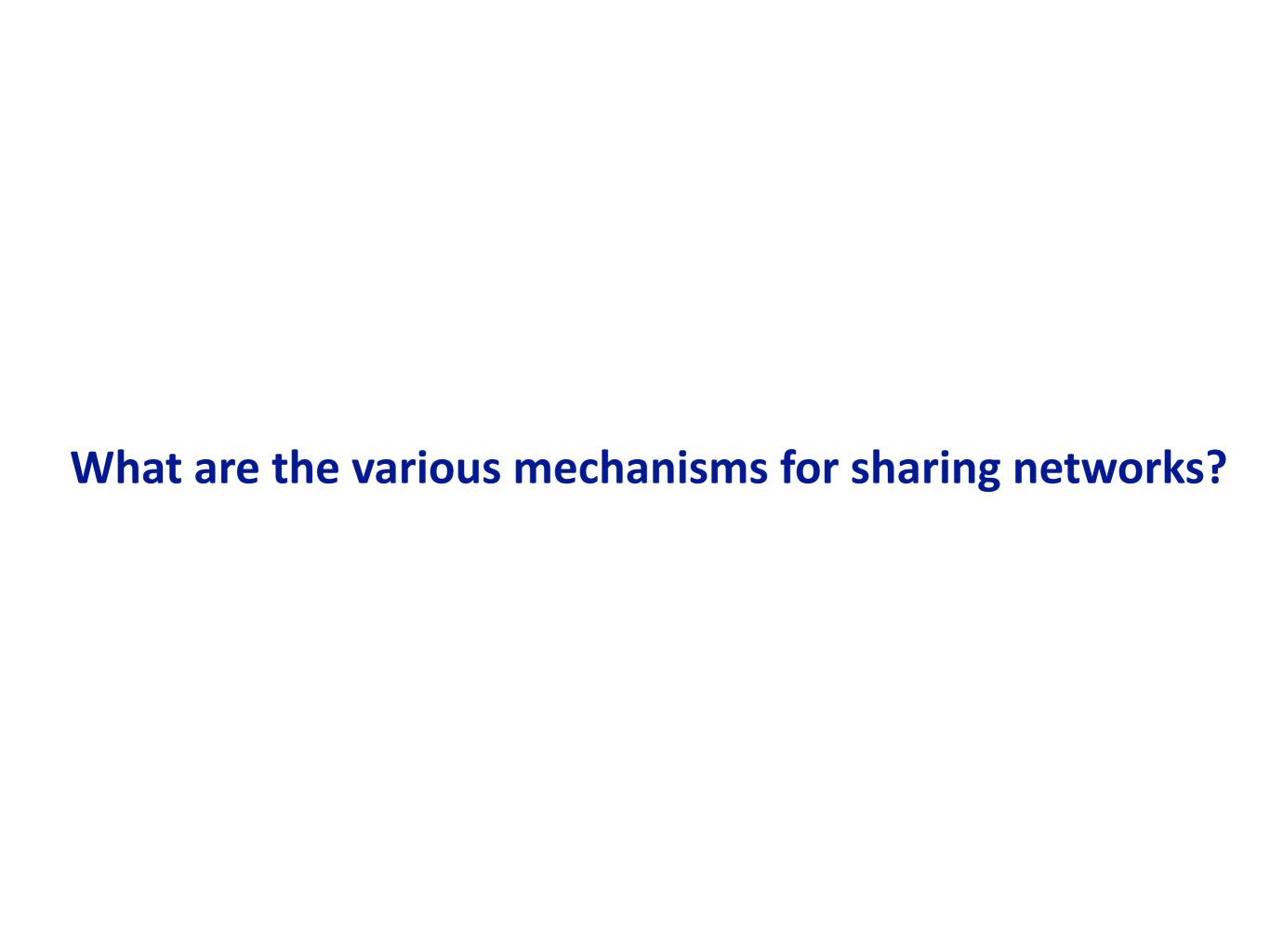
What are the performance metrics?

Performance metrics in computer networks!

- Bandwidth: Number of bits sent per second (bits per second, or bps)
 - Depends on
 - Hardware



- Network traffic conditions
-
- Delay: Time for <u>all bits</u> to go from source to destination (seconds)
 - Depends on
 - Hardware
 - Distance
 - Traffic from other sources
 - ...
- Many other performance metrics (reliability, etc.)
 - We will come back to other metrics later ...



Group Exercise 1:

How would you design a sharing mechanism?

Hint:

Think about sharing any resource (say, a computer)

Two approaches to sharing networks

- Reservations
- On demand

Two approaches to sharing networks

- First: Reservations
 - Reserve bandwidth needed in advance
 - Set up <u>circuits</u> and send data over that circuit
 - Must reserve for peak bandwidth
- How much bandwidth to reserve?
 - Applications may generate data at rate varying over time
 - 100MB in first second
 - 10MB in second second ...
 - Reservations must be made for "peak"

Circuit switching: Implementing reservations since ...

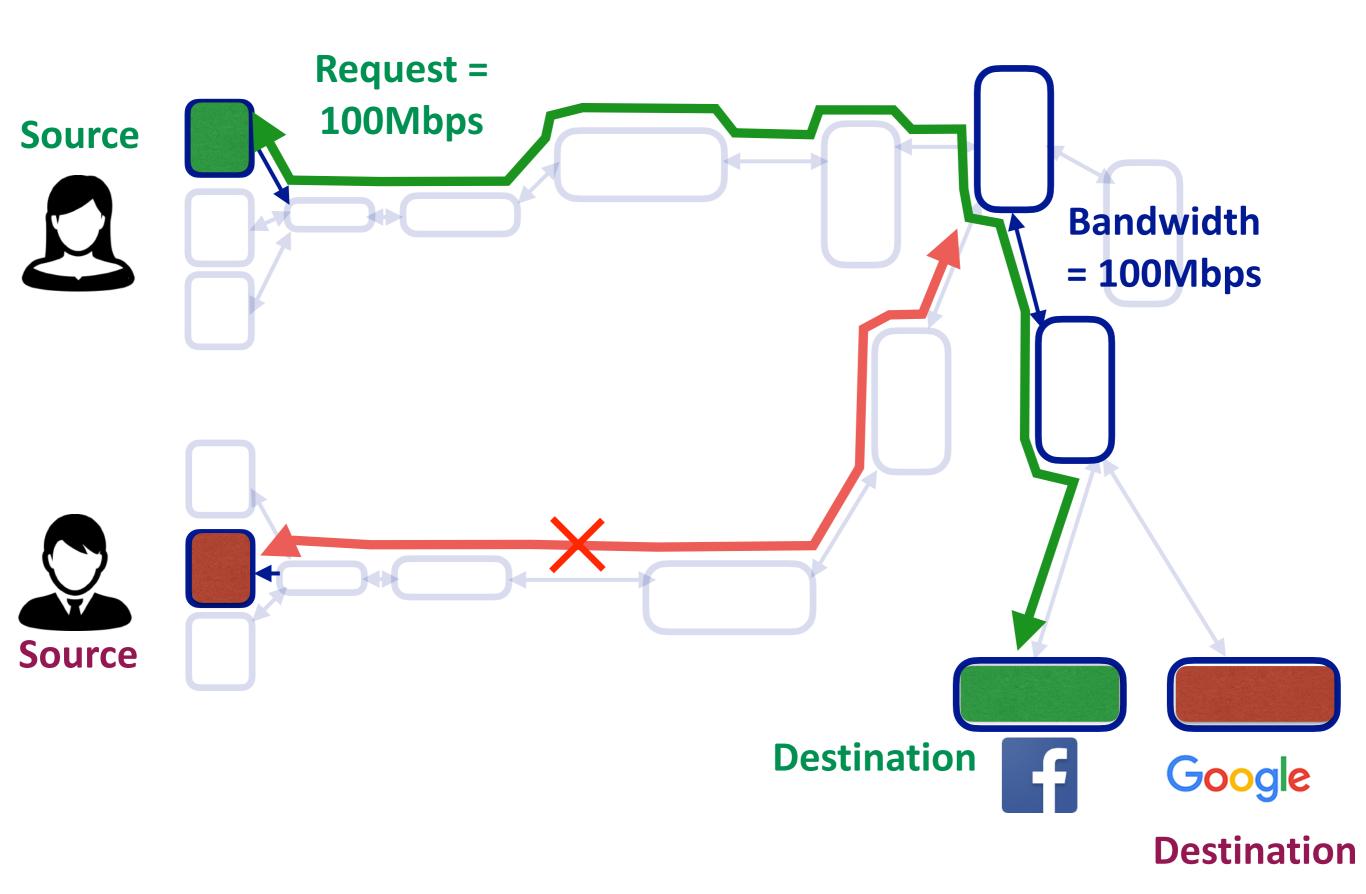
Telephone networks

One of the many approaches to implementing reservations

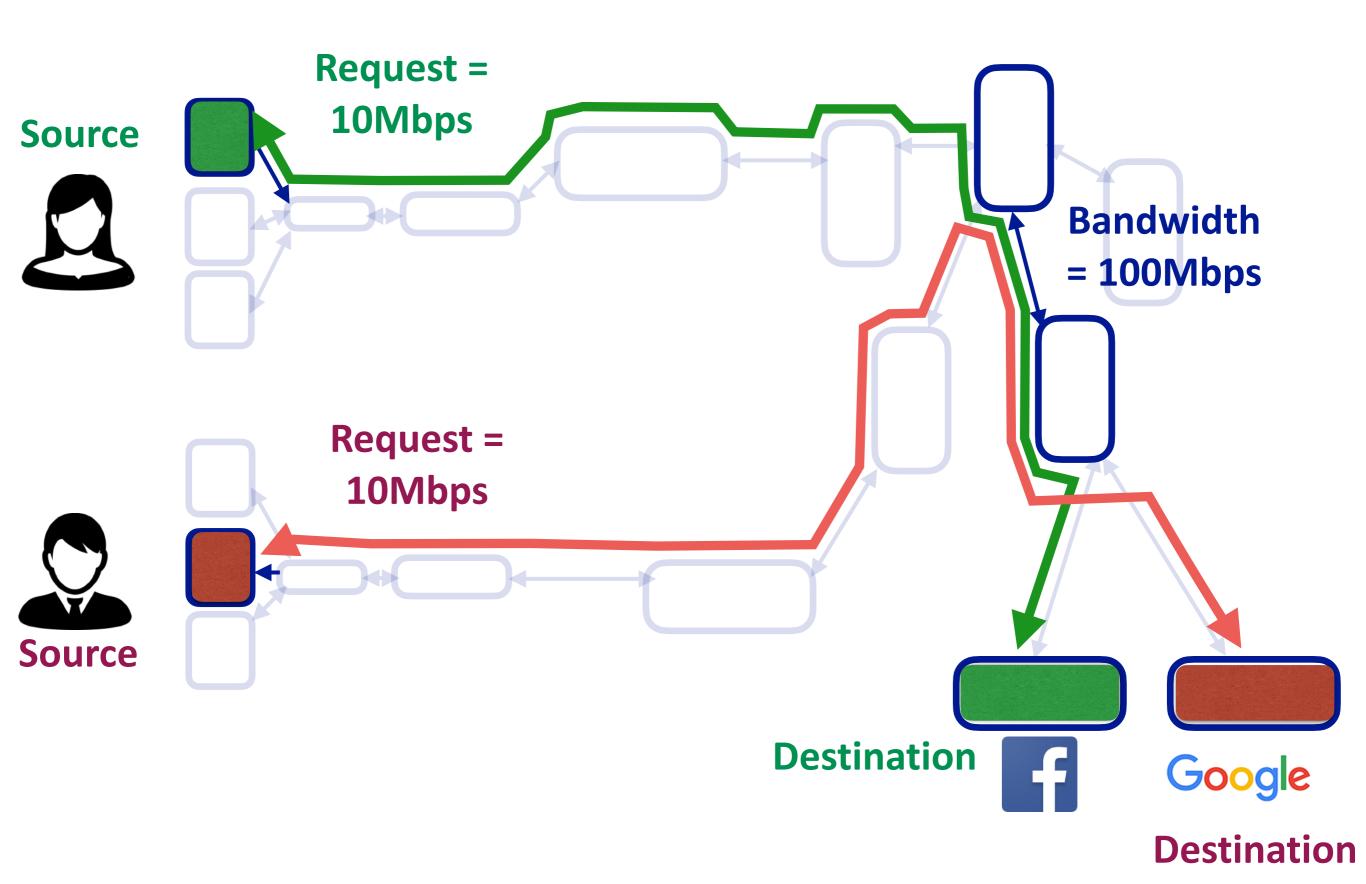
Mechanism:

- Source sends a reservation request for peak demand to destination
- Switches/routers establish a "circuit"
- Source sends data
- Source sends a "teardown circuit" message

Circuit switching: an example (red request fails)



Circuit switching: another example (red request succeeds)



Circuit switching and failures

- Circuit is established
- Link fails along path (!!!!!!)
 - First time we have seen failures making our life complicated.
 - Remember this moment.
 - Its gonna happen, over and over again.
- Must establish new circuit

Circuit switching doesn't route around failures!!

Circuit switching summary

Goods:

- Predictable performance
- Reliable delivery (assuming no failures)
- Simple forwarding mechanism

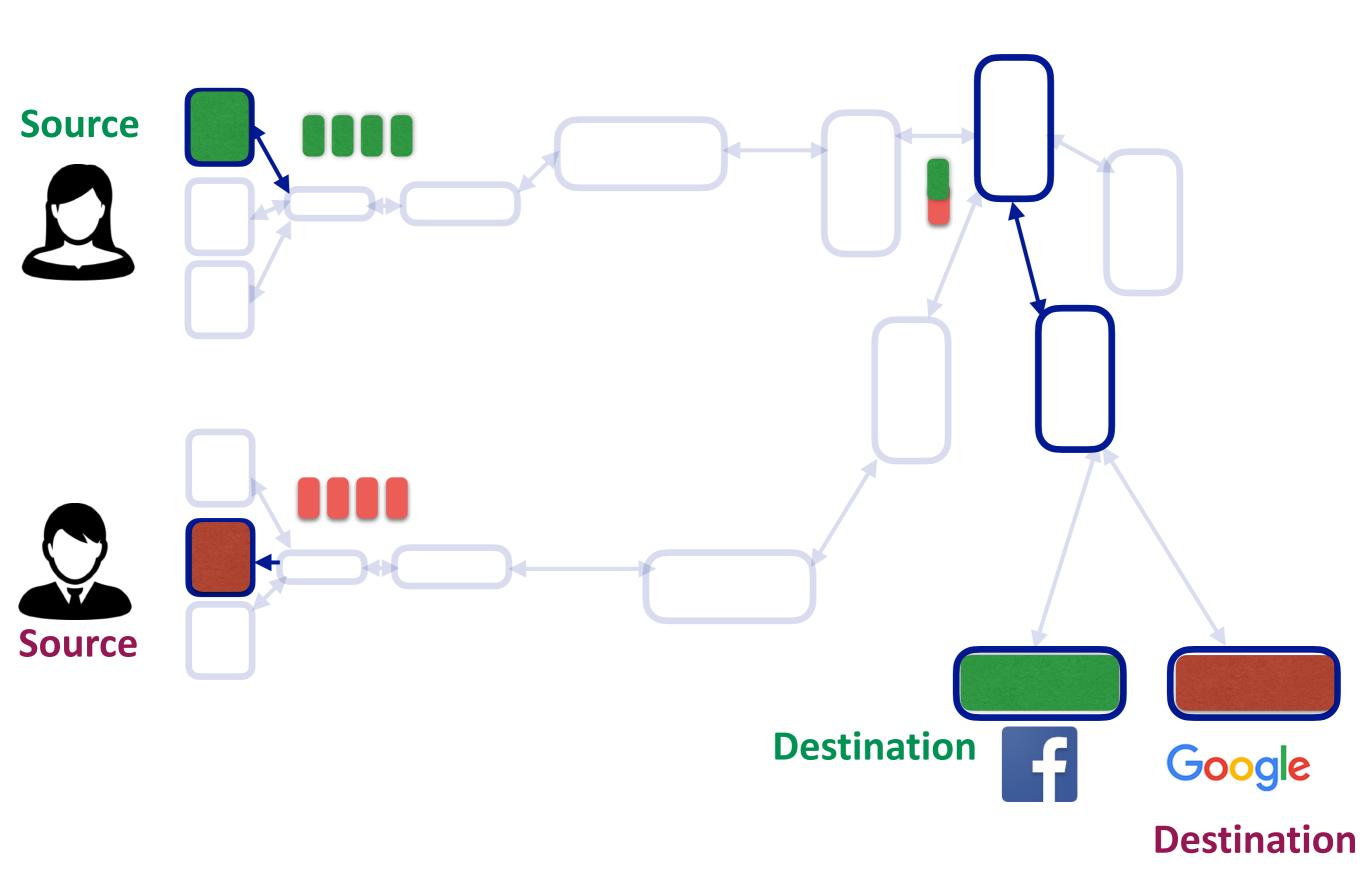
Not-so-goods

- Resource underutilization
- Blocked connections
- Connection set up overheads
- Per-connection state in switches (scalability problem)

Two approaches to sharing networks

- Second: On demand (also known as "best effort")
 - Designed specifically for the Internet
 - Break data into packets
 - Send packets when you have them
 - Hope for the best ...

Packet switching: an example



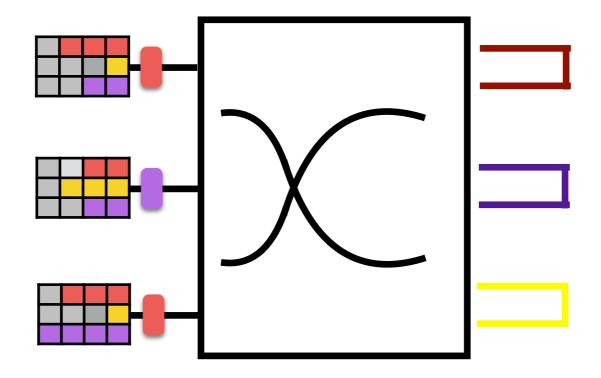
Packets

- Packets carry data (are bag of bits):
 - Header: meaningful to network (and network stack)
 - can be multiple headers
 - Body: meaningful only to application
 - More discussion in next lecture
- Body can be bits in a file, image, whatever
 - can have its own application "header"
- What information goes in the header?

What must headers contain to enable network functionality?

- Packets must describe where it should be sent
 - Requires an address for the destination host
 - can be multiple headers
- Packets must describe where its coming from
 - why?
 - Acknowledgments, etc.
- Thats the only way a router/switch can know what to do with the packet

Packet switching: what does a switch look like



Packet switching summary

Goods:

- No resource underutilization
 - A source can send more if others don't use resources
- No blocked connection problem
- No per-connection state
- No set-up cost

Not-so-goods:

- Unpredictable bandwidth availability
- Unpredictable delay/latency
- Packet header overhead

Circuits vs packets

- Pros for circuits:
 - Better application performance (reserved bandwidth)
 - More predictable and understandable (w/o failures)
- Pros for packets:
 - Better resource utilization
 - Easier recovery from failures
 - Faster startup to first packet delivered

Summary of network sharing

Statistical multiplexing

- Statistical multiplexing: combining demands to share resources efficiently
- Long history in computer science
 - Processes on an OS (vs every process has own core)
 - Cloud computing (vs every one has own datacenter)
- Based on the premise that:
 - Peak of aggregate load is << aggregate of peak load
- Therefore, it is better to share resources than to strictly partition them ...

Two approaches to sharing networks

Both embody statistical multiplexing

- Reservation: sharing at <u>connection</u> level
 - Resources shared between connections currently in system
 - Reserve the peak demand for a flow
- On-demand: sharing at <u>packet</u> level
 - Resources shared between packets currently in system
 - Resources given out on packet-by-packet basis
 - No reservation of resources

Understanding delay/latency

Packet Delay/Latency

- Consists of four components
 - Transmission delay (hardware properties)
 - Propagation delay (hardware properties, distance)
 - Queueing delay (traffic, switch internals)
 - Processing delay (end hosts)
- First, consider transmission and propagation delays
- Then queueing delay
- Ignore processing delays

Transmission delay

- How long does it take to push all the bits of a packet into a link?
- Packet size / Link Bandwidth
- Example:
 - Packet size = 1500Byte
 - Bandwidth = 100Mbps
 - 1500*8/100*1024*1024 seconds

Propagation delay

- How long does it take to move one bit from one end of a link to the other?
- Link length / Propagation speed of link
 - Propagation speed ~ some fraction of speed of light
- Example:
 - Length = 30,000 meters
 - Delay = 30*1000/3*100,000,000 second = 100us

Group Exercise 2:

How long does it take for a packet on a link?

Constraints:

- Packet size = 1000Byte
- Rate = 100Mbps
- Length = 30,000m