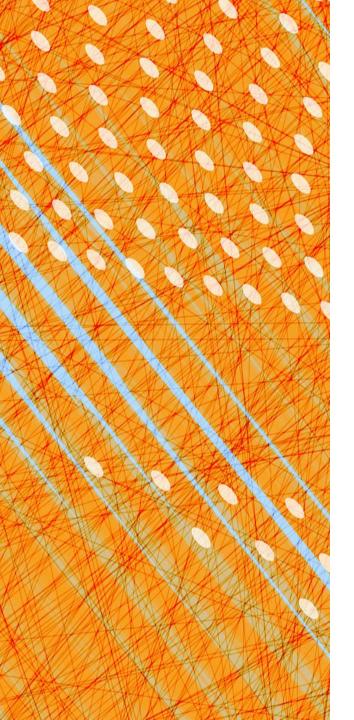
# Networking

CS 4410 Operating Systems



[R. Agarwal, L. Alvisi, A. Bracy, M. George, Kurose, Ross, E. Sirer, R. Van Renesse]



#### Introduction

Application Layer Transport Layer Link Layer Network Layer Routing

## **Basic Network Abstraction**

- A process can create "endpoints"
- Each endpoint has a unique address
- A message is a byte array
- Processes can:
  - receive messages on endpoints
  - send messages to endpoints

## Network "protocol"

Agreement between processes about the content of messages

**Syntax:** Layout of bits, bytes, fields, etc.

- message format

Semantics: what fields, messages mean

### **Example:**

HTTP "get" requests and responses

## Network Layering

## Network abstraction is usually layered

- Like Object Oriented-style inheritance
- Also like the hw/sw stack

Application
Presentation
Session
Transport
Network
Link
Physical

Application
Transport
Network
Link
Physical

Proposed 7-Layer ISO/OSI reference model (1970's)

Actual 5-Layer Internet Protocol Stack

## **OSI Layers**

Application	Network-aware applications, clients & servers	
Presentation	Translation between network and application formats (e.g., RPC packages, sockets)	
Session	Connection management	
Transport	Data transfer, reliability, packetization, retransmission.  Lets multiple apps share 1 network connection	
Network	Path determination across multiple network segments, routing, logical addressing.	
Link	Decides whose turn it is to talk, finds physical device on network.	
Physical	Exchanges bits on the media (electrical, optical, etc.)	

## Internet Protocol Stack

Application	exchanges <b>messages</b>	HTTP, FTP, DNS
Transport	Transports messages; exchanges <b>segments</b>	TCP, UDP
Network	Transports segments; exchanges datagrams	IP, ICMP (ping)
Link	Transports datagrams; exchanges <b>frames</b>	Ethernet, WiFi
Physical	Transports frames; exchanges <b>bits</b>	wires, signal encoding

## Who does what?

Application	HTTP, FTP, DNS (these^ are usually in libraries)		арр арр	
Transport	TCP, UDP		OS memory	
Network	IP, ICMP (ping)		CPU memory bus	
Link	Ethernet, WiFi		controller	
Physical	wires, signal encoding		physical transmission	
(Hard to draw firm lines here)				

- Each host has 1+ Network Interface Cards (NIC)
- Attaches into host's system buses
- Combination of hardware, software, firmware

## Layers support Modularity

#### Each layer:

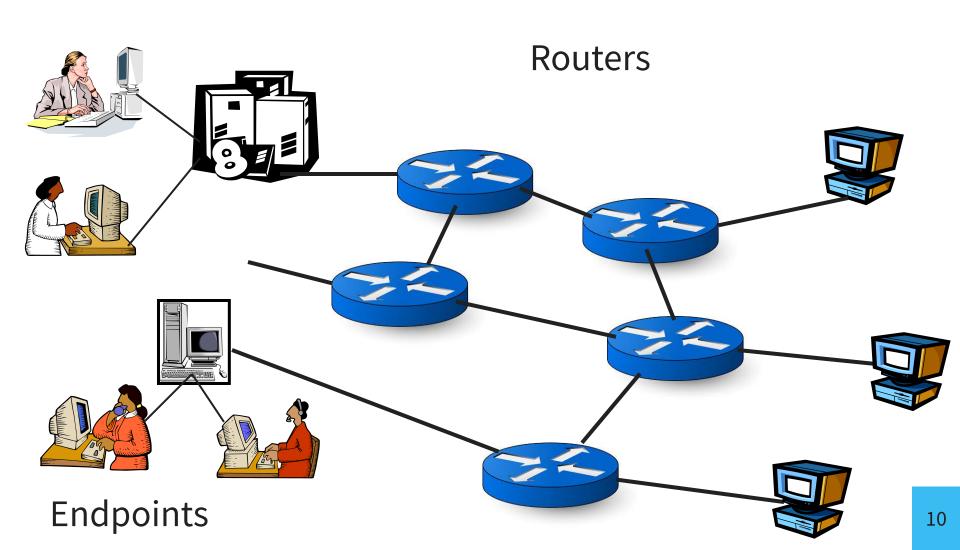
- relies on services from layer below
- exports services to layer above Can identify the relationship between distinct pieces of complex system.

#### Interfaces between layers:

- Hide implementation details
- Ease maintenance, updates
  - change of implementation of layer's service transparent to rest of system

## Internet, The Big Picture

How about an analogy?



# Application

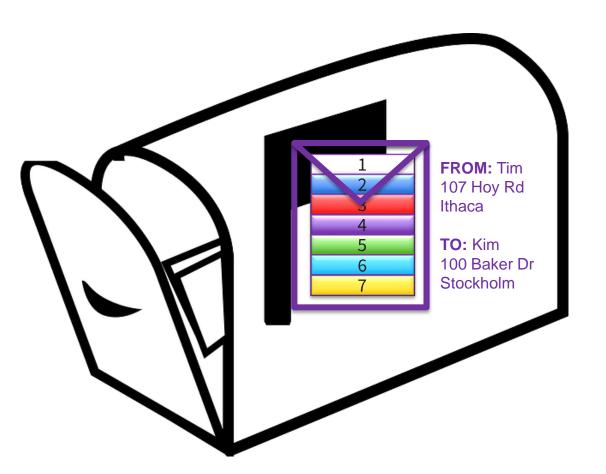












Transport Layer



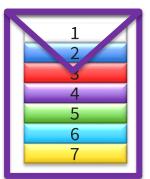
Ithaca Postman



Stockholm Postman



**SPEEDY** 



FROM: Tim 107 Hoy Rd Ithaca

**TO:** Kim 100 Baker Dr Stockholm

#### **Ithaca Sorting Office**



# Network Layer



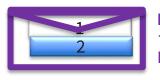
Ithaca Sorting Office



Stockholm Sorting Office



#### **Ithaca Sorting Office**



FROM: Tim TO: Kim
107 Hoy Rd 100 Baker Dr
Ithaca Stockholm



FROM: Tim TO: Kim
107 Hoy Rd 100 Baker Dr
Ithaca Stockholm



FROM: Tim TO: Kim
107 Hoy Rd 100 Baker Dr
Ithaca Stockholm



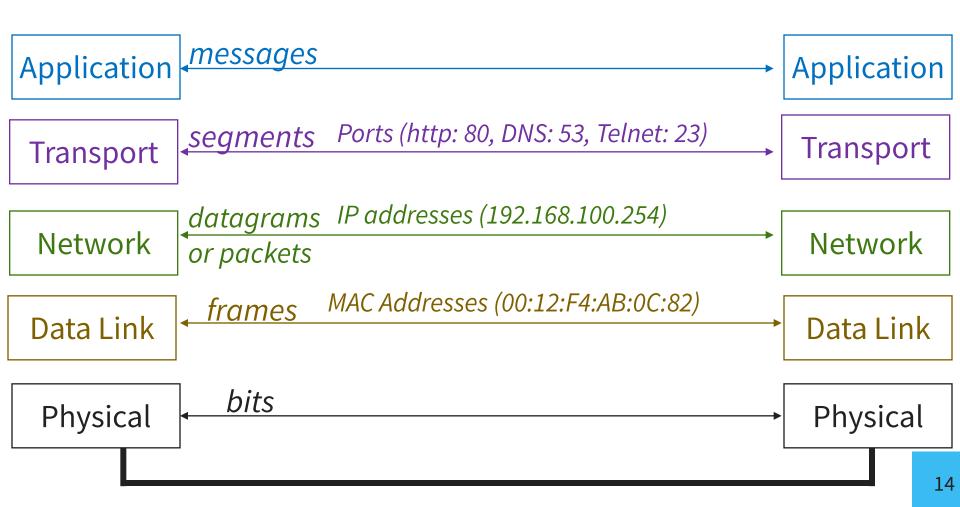
FROM: Tim TO: Kim
107 Hoy Rd 100 Baker Dr
Ithaca Stockholm



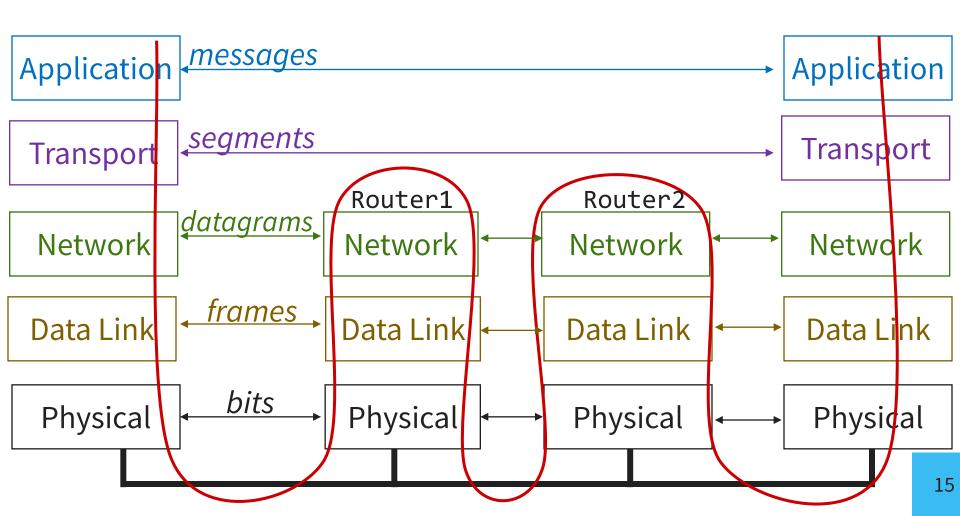




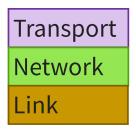
# The Big Picture



# The Big Picture



# Encapsulation

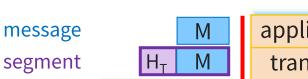


Headers

src & dst ports + ...

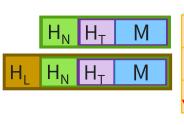
src & dest IP addr + ...

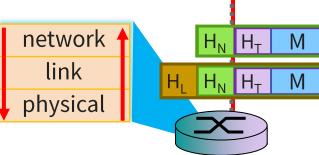
src & dest MAC addr + ...



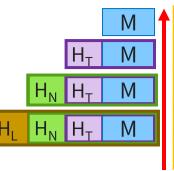
datagram  $H_N H_T M$ frame  $H_L H_N H_T M$  application transport network link physical

source





#### destination



application transport network link physical router

## End-to-End Argument

- Occam's Razor for Internet architecture
- Application-specific properties are best provided by the applications, not the network
  - Guaranteed, or ordered, packet delivery, duplicate suppression, security, *etc.*
- Internet performs the simplest packet routing and delivery service it can
  - Packets are sent on a best-effort basis
  - Higher-level applications do the rest

## End-to-End Example

#### Should the network guarantee packet delivery?

**Consider:** a file transfer program (read file from disk, send it, receiver reads packets & writes them to disk)

- **Q:** If network guarantees delivery, wouldn't applications be simpler? (no retransmissions!)
- A: no, still need to check that file was written to remote disk intact

#### A check is necessary if nodes can fail.

→ Applications need to be written to perform their own retransmits

Why burden the network with properties that can, and must, be implemented at the periphery?

## The Missing Layers

#### **Presentation**

translation between network & application formats (e.g., RPC packages, sockets). Allows communicating applications to interpret the meaning of data exchanged:

- data conversion
- character code translation
- compression
- encryption

#### **Session**

synchronization of data exchange:

- supports checkpointing and recovery schemes
- establish, manage, and tear down connections

Need these services?
Put them in your application.