CS 4220: Final Exam Solutions

May 14, 2013

1. (10 points)

(a) A 3-by-3 linear system with infinity-norm condition 10^8 is solved via the MATLAB backslash operator \setminus on a computer with unit roundoff 10^{-17} . Here is the computed solution:

xHat(1) = 1234.5678901234567 xHat(2) = 1.2345678901234567 xHat(3) = .00012345678901234567

Underline the digits that are most likely correct and justify your answer by explaining how the relative error $\|\hat{x} - x\|_{\infty}/\|x\|_{\infty}$ depends on both the condition and the unit roundoff. Recall that $\|v\|_{\infty} = \max |v_i|$.

Solution (5 points)

Since

$$\frac{\|\hat{x} - x\|_{\infty}}{\|x\|_{\infty}} \approx \operatorname{cond}_{\infty}(A) \operatorname{eps}$$

we have for this example that

$$|\hat{x}_i - x_i| \le \|\hat{x} - x\|_{\infty} \le \operatorname{cond}_{\infty}(A) \operatorname{eps} \|x\|_{\infty} \approx 10^8 (10^{-17})(10^3) \approx 10^{-6}$$

So underline each component value through the sixth decimal place.

-2 if you underline 6 significant digits in each component. The basic heuristic $||\hat{x} - x||/normx \approx \text{cond}(A)eps = 10^{-d}$ says the vector \hat{x} has about d correct digits. That does not translate into d correct digits for every component, except the largest one.

(b) Assume that a and b is are initialized floating point numbers with positive value and that the message "b is small compared to a" is displayed when the following code is executed:

if a + b == a
 disp('b is small compared to a')

end

What can you say about the actual magnitude of b?

Solution (5 points)

b must be less than the spacing of the floating point numbers at a, so roughly |b| < eps a where eps is the unit roundoff.

-2 if you say $b \approx eps$. For example $fl(1 + 2^{100}) = 2^{100}$.

2. (20 points)

(a) Show how the SVD can be used to solve the linear system

$$(A^T A + \mu I)x = A^T b$$

where $A \in \mathbb{R}^{m \times n}$, $n \leq m, \mu > 0$, and $b \in \mathbb{R}^m$. Answer by completing the following code in MATLAB:

```
% A, b, mu defined
[U,S,V] = svd(A);
[m,n] = size(A);
```

Solution (10 points)

% (USV')'(USV') + muI)x = (USV')b ---->
% (VS'SV' + mu*I)x = VS'U'b ---->
% V(S'S + mu*I)V'x = VS'U'b
% (S'S + mu*I)y = btilde where btilde = S'(U'*b) and y = V'x
d = diag(S);
btilde = d.*(U(:,1:n)'*b); % 3 points for transformed rhs
y = btilde ./ (d.^2 + mu); % 4 points for solution of transformed system
x = V*y % 3 points for transforming back to get x

S is not square so -3 for things that involve S².
-3 for dimension incompatibility, e.g., diag(S)*(U'*b)
-5 for correct but with an O(n³) computation, e.g. V'*A'*b instead of V'*(A'*b)

(b) Suppose A = QR is the QR factorization of a matrix $A \in \mathbb{R}^{m \times n}$. Assume that R(k, k) = 0 for some k > 1 and that all other entries along R's diagonal are nonzero. Show how to compute a nonzero vector $x \in \mathbb{R}^n$ so that Ax = 0. You may use the \setminus operator to solve triangular systems. Answer by completing the following code in MATLAB:

```
% A defined
[Q,R] = qr(A);
[m,n] = size(A);
```

Solution (10 points)

```
% Ax = 0 means QRx = 0 means Rx = 0
% 3 points for recognizing that that need null vector for R
% R(1:k,1:k) is singular with zero in its lower right corner.
% 2 points for this observation
y = R(1:k-1,1:k-1)\R(1:k-1,k); % [y;-1] is in the nullspace of R(1:k,1:k)
% 3 points for this
x = [y;-1;zeros(n-k,1)]
% 2 points
```

-3 if correct but you set k = n.

3. (15 points) If $C \in \mathbb{R}^{n \times n}$ is symmetric and positive definite and $u \in \mathbb{R}^n$, then

$$(C + uu^T)^{-1} = C^{-1} + \alpha v v^T$$

where $v = C^{-1}u$ and $\alpha = -1/(1 + u^T C^{-1}u)$. By making effective use of the Cholesky factorization and the above math fact, complete the following function so that it performs as specified:

```
function [x,z] = DoubleSolve(A,u,b)
% A is a symmetric positive definite n-by-n matrix.
% u and b are column n-vectors.
% x and z are column n-vectors with the property that Ax = b and
% (A + u*u')z = b.
```

You may use the \setminus operator to solve triangular systems.

Solution (15 points)

```
% Solve Ax = b using Cholesky for 5 points..
G = chol(A, 'lower');
x = G'\(G\b);
% Solve Av = u for 4 points
v = G'\(G\u);
% 6 points for using the math fact without inverse computation...
alfa = -1/(1+u'*v);
% z = (inv(A) + alfa*v*v')b = inv(A)*b + alfa*(v'*b)*v
z = x + alfa*(v'*b)*v
```

```
-4 if you have unnecessary solves with G
-2 for (v*v')*x instead of (v'*x)*v
-10 for any kind of inverse computation, e.g., inv(A) or A\ eye(n,n).
```

4. (15 points) Short answer.

(a) Assume that $A \in \mathbb{R}^{n \times n}$ is symmetric with distinct eigenvalues. Describe in English what the power method would try to compute when applied to the matrix $(A - \mu I)^{-1}$.

Solution (5 points)

If λ is the closest eigenvalue of A to μ and $Ax = \lambda x$, then $\lambda_* = 1/(\lambda - \mu)$ is the dominant eigenvalue of $(A - \mu I)^{-1}$ and $x_* = x$ is the corresponding eigenvector. The power method will try to find λ_* and x_* . If μ is equidistant to two eigenvalues then a problem.

(b) After k exact-arithmetic steps with starting unit 2-norm vector q_1 we assume that the Lanczos method computes part of the decomposition $Q^T A Q = T$ where Q is orthogonal with $Q(:, 1) = q_1$ and T is tridiagonal. Explain. What makes the method a "sparse matrix friendly"? What is the method typically used for?

Solution (5 points)

Will have Q(:, 1:k) and T(1:k, 1:k). (2 points)

Uses only matrix-vector products and the last two Lanczos vectors (2 points)

The extremal eigenvalues of T_k are good approximations for the extremal eigenvalues of A. (1 point)

(c) Any symmetric positive definite matrix A has a Cholesky factorization $A = GG^T$ where G is lower triangular. Why does MATLAB sparse Cholesky software compute the factorization $PAP^T = GG^T$ where P is a permutation matrix?

Solution (5 points)

P is chosen to minimize fill-in in the Cholesky factor

5. (15 points) For each of the following methods, draw a picture that communicates the main idea behind a step. No formulas are necessary. Just a labeled sketch that graphically indicates how the next iterate is obtained. (Such a picture for Newton's method would show the linear model and label its zero.)

(a) The Secant method for finding a zero of $f: \mathbb{R} \to \mathbb{R}$.

The picture should show show the secant line that goes through $(x_k, f(x_k) \text{ and } (x_{k-1}, f(x_{k-1})))$. (3 points). The intersection of that line with the x-axis defones x_{k+1} .

(b) The Golden Section search method for finding a minimum of $f:\mathbb{R} \to \mathbb{R}$ on [L, R] assuming that f'' is always positive.

The picture should show a function f(x) with f''(x) > 0 across the search interval [L, R]. (2 points). It should show two sample points c and d in the interval and a comparison based on comparison of f(c) and f(d). (1 point). It should show reduction of the search interval and the reuse of either f(c) of f(d) in th enext step (2 points)

(c) The steepest descent method with exact line search for finding a minimum of $f:\mathbb{R}^2 \to \mathbb{R}$. (Draw contours.)

Should show the negative gradient direction (2 points) and give a clue that one tries to minimize f in that direction.

6. (15 points)

(a) Consider the following MATLAB script for approximating the square root of a positive real number A:

```
L = A;
W = 1;
for k=1:10
L = (L+W)/2;
W = A/L;
end
```

Explain why this is essentially an instance of Newton's method.

Solution (10 points)

Equivalent to

```
L = A;
for k=1:10
L = (L+A/L)/2;
end
```

This is Newton's method applied to $f(L) = L^2 - A$:

$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x_k)} = x_k - \frac{x_k^2 - A}{2x_k} = \frac{1}{2} \left(x_k + \frac{A}{x_k} \right)$$

At least -5 if you thing that $f(x) = \sqrt{x}$.

(b) Newton's method is applied to a function $f:\mathbb{R} \to \mathbb{R}$ that has a single root x_* . Here is a heuristic result that relates the error at step k + 1 to the error at step k:

$$|x_{k+1} - x_*| \approx \frac{|f''(x_k)|}{2|f'(x_k)|} |x_k - x_*|^2.$$

What additional information is required before one can be confident that the iteration converges quadratically for a particular initial guess x_0 ?

Solution

The derivative must be bounded away from zero.

 x_0 must be close enough. How close depends on the second derivative (the bigger the closer) and the first derivative (the small the closer).

Can also say that |f''(x)/f'x| bounded by M and $|x_0 - x_*| \cdot M < 1$

7. (15 points) Suppose $t, x, y \in \mathbb{R}^m$ are given and that we wish to use lsqnonlin to determine $a^T = [a_1, \ldots, a_k]$ and $\rho^T = [\rho_1, \ldots, \rho_k]$ so that

$$\phi(a,\rho) = \frac{1}{2} \sum_{i=1}^{m} \left(x_i - \sum_{j=1}^{k} a_j \cos\left(\frac{2\pi t_i}{\rho_j}\right) \right)^2 + \frac{1}{2} \sum_{i=1}^{m} \left(y_i - \sum_{j=1}^{k} a_j \sin\left(\frac{2\pi t_i}{\rho_j}\right) \right)^2$$

is minimized. To use lsqnonlin we must design a vector-valued function $R(a, \rho)$ with the property that

$$\phi(a,\rho) = \frac{1}{2}R(a,\rho)^T R(a,\rho)$$

Express $R(a, \rho)$ in the form

$$R(a, \rho) = \text{Matrix} \times \{\text{vector}\} - \{\text{vector}\}$$

Clearly define the matrix and the two vectors.

Solution

Since

$$\sum_{i=1}^{m} \left(x_i - \sum_{j=1}^{k} a_j \cos\left(\frac{2\pi t_i}{\rho_j}\right) \right)^2 = \|Ca - x\|_2^2$$

where $C = (c_{ij}) \in \mathbb{R}^{m \times k}$ and $c_{ij} = \cos(2\pi t_i/\rho_j)$ and

$$\sum_{i=1}^{m} \left(y_i - \sum_{j=1}^{k} a_j \sin\left(\frac{2\pi t_i}{\rho_j}\right) \right)^2 = \|Sa - y\|_2^2$$

where $C = (c_{ij}) \in \mathbb{R}^{m \times k}$ and $c_{ij} = \cos(2\pi t_i/\rho_j)$ we have

$$\phi(a,\rho) = \frac{1}{2} \|Ca - x\|_2^2 + \frac{1}{2} \|Sa - y\|_2^2 = \frac{1}{2} \|\begin{bmatrix}C\\S\end{bmatrix} a - \begin{bmatrix}x\\y\end{bmatrix} \|_2^2$$

Thus,

$$R = \left[\begin{array}{c} C \\ S \end{array} \right] a - \left[\begin{array}{c} x \\ y \end{array} \right]$$

5 points for the matrix

 $3~{\rm points}$ for the rhs

2 points for the vector that is multiplied by the matrix.

Up to -7 for stuff like [C S][a; a] = x + y