

★ TYPES! ★

STLC a.k.a. λ^{\rightarrow}

SYNTAX \rightarrow

$n \in \text{Nat}$

$e ::= x$

$| e_1 e_2$

$| \lambda x : \tau . e$

$| n$

$| ()$

$| e_1 + e_2$

$\tau ::= \text{int}$

$| \text{unit}$

$| \tau_1 \rightarrow \tau_2$

$(\lambda x : \text{int} . 5) ()$

$\text{int} \rightarrow \text{int}$
 $(\text{int} \rightarrow \text{int}) \rightarrow (\text{int} \rightarrow \text{int})$

unit

$\text{int} \rightarrow \text{int}$

int

$((\lambda x : \text{int} \rightarrow \text{int} . x) (\lambda x : \text{int} . 5)) ()$

$\text{int} \rightarrow \text{int}$

unit

DYNAMIC
SEMANTICS

ARE BASICALLY THE SAME

$$\frac{e \rightarrow e'}{E[e] \rightarrow E[e']}$$
$$\frac{m = n_1 + n_2}{n_1 + n_2 \rightarrow m}$$
$$\frac{}{(\lambda x : \tau. e) v \rightarrow e\{v/x\}}$$

(IGNORED)

WELL-TYPED PROGRAMS
DO NOT GET STUCK.

$$42 + \lambda x : \text{int}. x \not\rightarrow$$
$$() \not\rightarrow$$
$$(\lambda x : \text{int}. x) () \rightarrow ()$$

$$\Gamma \vdash e : \tau$$

\uparrow TYPE CONTEXT \uparrow TYPE

$$\Gamma : \text{Var} \rightarrow \text{Type}$$

$$(x : \text{int}) \vdash x : \text{int}$$

$$\emptyset \vdash \lambda x : \text{int}. x : \text{int} \rightarrow \text{int}$$

$$\Gamma [x \mapsto \tau]$$

$$\Gamma, x : \tau$$

$$\boxed{\Gamma \vdash e : \tau}$$

$$\frac{}{\Gamma \vdash () : \text{unit}}$$

$$\frac{\Gamma \vdash n : \text{int}}{\Gamma \vdash x : \tau}$$

$$\frac{\Gamma(x) = \tau}{\Gamma \vdash x : \tau}$$

~~$$(y : \text{int}) \vdash x : \tau$$~~

$$\frac{\Gamma \vdash e_1 : \text{int} \quad \Gamma \vdash e_2 : \text{int}}{\Gamma \vdash e_1 + e_2 : \text{int}}$$

(2 + 4) + 2

$$\frac{\Gamma [x \mapsto \tau] \vdash e : \tau'}{\Gamma \vdash \lambda x : \tau. e : \tau \rightarrow \tau'}$$

$$\frac{\Gamma \vdash e_1 : \tau_1 \rightarrow \tau_2 \quad \Gamma \vdash e_2 : \tau_1}{\Gamma \vdash e_1 e_2 : \tau_2}$$

$e = (\lambda x:\text{int}. x + 40) 2$

$\vdash e : \text{int}$

$\vdash e : \tau$

$\emptyset \vdash e : \tau$

"e IS WELL-TYPED"

$\exists \tau. \vdash e : \tau$

