CS 4110

Programming Languages & Logics

Lecture 17
Definitional Translation & Continuations

Definitional Translation

We know how to encode Booleans, conditionals, natural numbers, and recursion in λ-calculus.

Can we define a *real* programming language by translating everything in it into the λ-calculus?

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Can we define a *real* programming language by translating everything in it into the λ-calculus?

In definitional translation, we define a denotational semantics where the target is a simpler programming language instead of mathematical objects.

Multi-Argument λ-calculus

Let's define a version of the λ-calculus that allows functions to take multiple arguments.

$$e ::= x \mid \lambda x_1, \ldots, x_n.\, e \mid e_0\; e_1\; \ldots\; e_n$$

Multi-Argument λ-calculus

We can define a CBV operational semantics:

$$E ::= [\cdot] \mid v_0 \ \ldots \ v_{i-1}\, E\, e_{i+1} \ \ldots \ e_n$$

$$\frac{e \to e'}{E[e] \to E[e']}$$

$$\frac{}{(\lambda x_1, \ldots, x_n.\, e_0)\, v_1 \ \ldots \ v_n \to e_0\{v_1/x_1\}\{v_2/x_2\} \ldots \{v_n/x_n\}}\ \beta$$

The evaluation contexts ensure that we evaluate multi-argument applications $e_0\, e_1 \ \ldots \ e_n$ from left to right.

Definitional Translation

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We can define a translation $\mathcal{T}[\![\cdot]\!]$ that takes an expression in the multi-argument λ-calculus and returns an equivalent expression in the pure λ-calculus.

$$\begin{aligned}
\mathcal{T}[\![x]\!] &= x \\
\mathcal{T}[\![\lambda x_1, \ldots, x_n.\, e]\!] &= \lambda x_1.\ \ldots\ \lambda x_n.\, \mathcal{T}[\![e]\!] \\
\mathcal{T}[\![e_0\, e_1\, e_2\ \ldots\ e_n]\!] &= (\ldots((\mathcal{T}[\![e_0]\!]\, \mathcal{T}[\![e_1]\!])\, \mathcal{T}[\![e_2]\!])\ldots\ \mathcal{T}[\![e_n]\!])
\end{aligned}$$

This translation *curries* the multi-argument λ-calculus.

Products (Pairs) and Let

Syntax

$$
\begin{aligned}
e ::= {} & x \\
\mid {} & \lambda x.\, e \\
\mid {} & e_1\ e_2 \\
\mid {} & (e_1, e_2) \\
\mid {} & \#1\ e \\
\mid {} & \#2\ e \\
\mid {} & \mathsf{let}\ x = e_1\ \mathsf{in}\ e_2
\end{aligned}
$$

$$
\begin{aligned}
v ::= {} & \lambda x.\, e \\
\mid {} & (v_1, v_2)
\end{aligned}
$$

Products (Pairs) and Let

Evaluation Contexts

$$
\begin{aligned}
E ::= {} & [\cdot] \\
\mid {} & E\ e \\
\mid {} & v\ E \\
\mid {} & (E, e) \\
\mid {} & (v, E) \\
\mid {} & \#1\ E \\
\mid {} & \#2\ E \\
\mid {} & \mathsf{let}\ x = E\ \mathsf{in}\ e_2
\end{aligned}
$$

Products (Pairs) and Let

Semantics

$$\frac{e \to e'}{E[e] \to E[e']}$$

$$\frac{}{(\lambda x.\, e)\; v \to e\{v/x\}}\;\beta$$

$$\frac{}{\#1\,(v_1, v_2) \to v_1} \qquad \frac{}{\#2\,(v_1, v_2) \to v_2}$$

$$\frac{}{\text{let}\; x = v\; \text{in}\; e \to e\{v/x\}}$$

Products (Pairs) and Let

Translation

$$
\begin{aligned}
\mathcal{T}[\![x]\!] &= x \\
\mathcal{T}[\![\lambda x.\, e]\!] &= \lambda x.\, \mathcal{T}[\![e]\!] \\
\mathcal{T}[\![e_1\ e_2]\!] &= \mathcal{T}[\![e_1]\!]\ \mathcal{T}[\![e_2]\!] \\
\mathcal{T}[\![(e_1, e_2)]\!] &= (\lambda x.\, \lambda y.\, \lambda f.\, f\, x\, y)\ \mathcal{T}[\![e_1]\!]\ \mathcal{T}[\![e_2]\!] \\
\mathcal{T}[\![\#1\ e]\!] &= \mathcal{T}[\![e]\!]\ (\lambda x.\, \lambda y.\, x) \\
\mathcal{T}[\![\#2\ e]\!] &= \mathcal{T}[\![e]\!]\ (\lambda x.\, \lambda y.\, y) \\
\mathcal{T}[\![\mathsf{let}\ x = e_1\ \mathsf{in}\ e_2]\!] &= (\lambda x.\, \mathcal{T}[\![e_2]\!])\ \mathcal{T}[\![e_1]\!]
\end{aligned}
$$

Laziness

Consider the call-by-name λ-calculus...

Syntax

$$
\begin{aligned}
e ::= {} & x \\
\mid {} & e_1\ e_2 \\
\mid {} & \lambda x.\, e
\end{aligned}
$$

$$
v ::= \lambda x.\, e
$$

Semantics

$$
\frac{e_1 \to e_1'}{e_1\ e_2 \to e_1'\ e_2}
\qquad\qquad
\frac{}{(\lambda x.\, e_1)\ e_2 \to e_1\{e_2/x\}}\ \beta
$$

Laziness

Translation

$$
\begin{aligned}
\mathcal{T}[\![x]\!] &= x\,(\lambda y.\, y) \\
\mathcal{T}[\![\lambda x.\, e]\!] &= \lambda x.\, \mathcal{T}[\![e]\!] \\
\mathcal{T}[\![e_1\ e_2]\!] &= \mathcal{T}[\![e_1]\!]\,(\lambda z.\, \mathcal{T}[\![e_2]\!]) \qquad z \text{ is not a free variable of } e_2
\end{aligned}
$$

References

Syntax

$$
\begin{aligned}
e ::= {} & x \\
\mid {} & \lambda x.\, e \\
\mid {} & e_0\ e_1
\end{aligned}
$$

$$v ::= \lambda x.\, e$$

References

Syntax

$$
\begin{aligned}
e ::=\ & x \\
\mid\ & \lambda x.\, e \\
\mid\ & e_0\ e_1 \\
\mid\ & \mathsf{ref}\ e
\end{aligned}
$$

$$v ::= \lambda x.\, e$$

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\begin{aligned}
e ::= {} & x \\
\mid {} & \lambda x.\, e \\
\mid {} & e_0\ e_1 \\
\mid {} & \mathsf{ref}\ e \\
\mid {} & !e
\end{aligned}
$$

$$v ::= \lambda x.\, e$$

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\mid {} & \mathsf{ref}\ e \\
\mid {} & !e \\
\mid {} & e_1 := e_2
\end{aligned}
$$

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v ::= \lambda x.\, e
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\mid {} & \ell
\end{aligned}
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References

Evaluation Contexts

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Evaluation Contexts

$$\begin{aligned} E ::= {} & [\cdot] \\ & \mid E\ e \\ & \mid v\ E \\ & \mid \mathrm{ref}\ E \\ & \mid\ !E \\ & \mid E := e \end{aligned}$$

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Evaluation Contexts

$$\begin{aligned} E ::= {} & [\cdot] \\ & \mid E\, e \\ & \mid v\, E \\ & \mid \text{ref}\ E \\ & \mid\ !E \\ & \mid E := e \\ & \mid v := E \end{aligned}$$

References

Semantics

$$\frac{\langle \sigma, e\rangle \rightarrow \langle \sigma', e'\rangle}{\langle \sigma, E[e]\rangle \rightarrow \langle \sigma', E[e']\rangle} \qquad \frac{}{\langle \sigma, (\lambda x.\, e)\; v\rangle \rightarrow \langle \sigma, e\{v/x\}\rangle}\;\beta$$

$$\frac{\ell \notin dom(\sigma)}{\langle \sigma, \mathsf{ref}\; v\rangle \rightarrow \langle \sigma[\ell \mapsto v], \ell\rangle} \qquad \frac{\sigma(\ell) = v}{\langle \sigma,\ !\ell\rangle \rightarrow \langle \sigma, v\rangle}$$

$$\frac{}{\langle \sigma, \ell := v\rangle \rightarrow \langle \sigma[\ell \mapsto v], v\rangle}$$

References

Translation

...left as an exercise to the reader. ;-)

Adequacy

How do we know if a translation is correct?

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Every target evaluation should represent a source evaluation...

Definition (Soundness)

$\forall e \in \mathbf{Exp}_{\text{src}}.$ if $\mathcal{T}[\![e]\!] \rightarrow^{*}_{\text{trg}} v'$ then $\exists v.\ e \rightarrow^{*}_{\text{src}} v$
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...and every source evaluation should have a target evaluation:

Definition (Completeness)

$\forall e \in \mathbf{Exp}_{\text{src}}.$ if $e \rightarrow^{*}_{\text{src}} v$ then $\exists v'.\ \mathcal{T}[\![e]\!] \rightarrow^{*}_{\text{trg}} v'$ and v' equivalent to v

Continuations

In the preceding translations, the control structure of the source language was translated directly into the corresponding control structure in the target language.

For example:

$$\mathcal{T}[\![\lambda x.\, e]\!] = \lambda x.\, \mathcal{T}[\![e]\!]$$
$$\mathcal{T}[\![e_1\ e_2]\!] = \mathcal{T}[\![e_1]\!]\ \mathcal{T}[\![e_2]\!]$$

What can go wrong with this approach?

Continuations

- A snippet of code that represents “the rest of the program”
- Can be used directly by programmers...
- ...or in program transformations by a compiler
- Make the control flow of the program explicit
- Also useful for defining the meaning of features like exceptions

Example

Consider the following expression:

$$(\lambda x.\, x)\;((1 + 2) + 3) + 4$$

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If we make all of the continuations explicit, we obtain:

$$k_0 = \lambda v.\, (\lambda x.\, x)\, v$$

Example

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If we make all of the continuations explicit, we obtain:

$$\begin{aligned} k_0 &= \lambda v.\, (\lambda x.\, x)\, v \\ k_1 &= \lambda a.\, k_0\, (a + 4) \end{aligned}$$

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Example

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If we make all of the continuations explicit, we obtain:

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k_0 &= \lambda v.\, (\lambda x.\, x)\; v \\
k_1 &= \lambda a.\, k_0\; (a + 4) \\
k_2 &= \lambda b.\, k_1\; (b + 3) \\
k_3 &= \lambda c.\, k_2\; (c + 2)
\end{aligned}$$

The original expression is equivalent to $k_3\; 1$, or:

$$(\lambda c.\, (\lambda b.\, (\lambda a.\, (\lambda v.\, (\lambda x.\, x)\; v)\; (a + 4))\; (b + 3))\; (c + 2))\; 1$$

Example (Continued)

Recall that let $x = e$ in e' is syntactic sugar for $(\lambda x.\, e')\ e$.

Hence, we can rewrite the expression with continuations more succinctly as

$$
\begin{array}{l}
\text{let } c = 1 \text{ in} \\
\text{let } b = c + 2 \text{ in} \\
\text{let } a = b + 3 \text{ in} \\
\text{let } v = a + 4 \text{ in} \\
(\lambda x.\, x)\ v
\end{array}
$$

CPS Transformation

We write $\mathcal{CPS}[\![e]\!]\,k = \ldots$ instead of $\mathcal{CPS}[\![e]\!] = \lambda k.\ \ldots$

We assume that the new variables introduced are "fresh."

CPS Transformation

$$\mathcal{CPS}[\![n]\!]\,k = k\,n$$

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CPS Transformation

$$\begin{aligned}\mathcal{CPS}[\![n]\!]\,k &= k\,n\\ \mathcal{CPS}[\![e_1 + e_2]\!]\,k &= \mathcal{CPS}[\![e_1]\!]\,(\lambda n.\,\mathcal{CPS}[\![e_2]\!]\,(\lambda m.\,k\,(n+m)))\end{aligned}$$

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\mathcal{CPS}[\![x]\!]\,k &= k\,x
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\mathcal{CPS}[\![x]\!]\,k &= k\,x \\
\mathcal{CPS}[\![\lambda x.\,e]\!]\,k &= k\,(\lambda x.\,\lambda k'.\,\mathcal{CPS}[\![e]\!]\,k')
\end{aligned}
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