CS 4110

Programming Languages & Logics

Lecture 37 Concurrency and Victory Lap

5 December 2014

Announcements

• Foster Office Hours 11am-12pm today

$x, y, z \in \mathcal{N}$

Names

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 $\pi ::= \tau \mid \overline{x}\langle y \rangle \mid x(y) \mid [x = y] \pi$ Prefixes

$$\begin{array}{rcl} x,y,z & \in & \mathcal{N} & & \text{Names} \\ \pi & ::= & \tau & \mid \bar{x}\langle y \rangle & \mid x(y) & \mid [x = y] \pi & & \text{Prefixes} \\ M,N & ::= & \mathbf{0} & \mid \pi.P & \mid M + M & & \text{Summations} \\ P,Q,R & ::= & M & \mid P_1 & \mid P_2 & \mid \nu x. P & \mid !P & & \text{Processes} \end{array}$$

 $\overline{\tau.P+M
ightarrow P}$ R-Tau

$$\overline{\tau.P+M \rightarrow P}$$
 R-Tau

$$\overline{(\overline{x}\langle y\rangle.P_1+M_1)\mid (x(z).P_2+M_2)\rightarrow P_1\mid P_2\{y/z\}}$$
 R-React

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 R-React

$$\frac{P_1 \rightarrow P'_1}{P_1 \mid P_2 \rightarrow P'_1 \mid P_2} \text{ R-Par}$$

$$\frac{P \to P'}{\nu x. \ P \to \nu x. \ P'} \text{ R-Res}$$

$$\frac{P \equiv P' \qquad P' \to Q' \qquad Q' \equiv Q}{P \to Q} \text{ R-Struct}$$

Idea: encode a boolean value *b* as a process that receives two channels *t* and *f* on the channel *l* where the boolean is "located" and then signals on the corresponding channel

$$True(I) \triangleq I(t, f).\overline{t}$$

Idea: encode a boolean value *b* as a process that receives two channels *t* and *f* on the channel *l* where the boolean is "located" and then signals on the corresponding channel

$$True(I) \triangleq l(t, f).\overline{t}$$

False(I) $\triangleq l(t, f).\overline{f}$

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$$True(I) \triangleq l(t, f).\overline{t}$$

$$False(I) \triangleq l(t, f).\overline{f}$$

$$Cond(P, Q)(I) \triangleq \nu t, f. (\overline{l}\langle t, f \rangle.(t.P + f.Q))$$

Idea: encode a natural number value *n* as a process that receives two channels *s* and *z* on the channel *c* where the number is "located" and then signals on *s n* times terminated by *z* Idea: encode a natural number value *n* as a process that receives two channels *s* and *z* on the channel *c* where the number is "located" and then signals on *s n* times terminated by *z*

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$$Zero(c) \triangleq c(s, z).\overline{z}$$

Succ(n)(c)
$$\triangleq c(s, z).\overline{n}\langle s, z \rangle.\overline{s}$$

Encoding Lists

Idea: encode a list l as a process that receives two channels c and n on the channel l where the list is "located" and then signals on c with each value of the list, terminated by n

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 $\begin{aligned} \text{Nil}(I) &\triangleq l(n,c).\overline{n} \\ \text{Cons}(H,T)(I) &\triangleq \nu h, t. \left(l(n,c).\overline{c}\langle h,t\rangle \mid H\langle h\rangle \mid T\langle t\rangle \right) \end{aligned}$

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$$\begin{array}{lll} \text{Nil}(l) &\triangleq & l(n,c).\overline{n} \\ \text{Cons}(H,T)(l) &\triangleq & \nu h, t. \left(l(n,c).\overline{c}\langle h,t\rangle \mid H\langle h\rangle \mid T\langle t\rangle \right) \\ & \text{IsNil}(L)(r) &\triangleq & \nu l, n, c. \left(L\langle l\rangle \mid \overline{l}\langle n,c\rangle.(n.\text{True}\langle r\rangle + c(h,t).\text{False}\langle r\rangle) \right) \end{array}$$

Pattern Matching

We can encode pattern matching on lists

case l of Nil? \Rightarrow P Cons?(h, t) \Rightarrow Q

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 $\nu n, c. \overline{l} \langle n, c \rangle n.P + c(h, t).Q$

$$\begin{array}{ll} Copy\langle l,m\rangle &\triangleq & \text{case } l \text{ of} \\ & \text{Nil?} \Rightarrow \text{Nil}\langle m \rangle \\ & \text{Cons?}(h,t) \Rightarrow \nu t'. \left(m(n,c).\overline{c}\langle h,t' \rangle \mid \text{Copy}\langle t,t' \rangle\right) \end{array}$$

$$\begin{array}{ll} Copy\langle I,m\rangle & \triangleq & \text{case } I \text{ of} \\ & \text{Nil?} \Rightarrow \text{Nil}\langle m \rangle \\ & \text{Cons?}(h,t) \Rightarrow \nu t'. \left(m(n,c).\overline{c}\langle h,t' \rangle \mid \text{Copy}\langle t,t' \rangle\right) \end{array}$$

$$\begin{aligned} \text{Join}\langle k, l, m \rangle &\triangleq \text{ case } k \text{ of} \\ &\text{Nil}? \Rightarrow \text{Copy}\langle l, m \rangle \\ &\text{Cons}?(h, t) \Rightarrow \nu t'. \left(m(n, c).\overline{c}\langle h, t' \rangle \mid \text{Join}\langle t, l, t' \rangle \right) \end{aligned}$$

We can put a ! in front of processes to turn them into servers create arbitrary numbers of the original process

$$\begin{aligned} \text{Nil}(I) &\triangleq !!(n,c).\overline{n} \\ \text{Cons}(H,T)(I) &\triangleq \nu h, t. (!!(n,c).\overline{c}\langle h, t \rangle \mid H\langle h \rangle \mid T\langle t \rangle) \end{aligned}$$

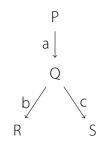
This causes the list to still exist after sending or receiving a message

Encoding λ -calculus

$$\begin{split} \llbracket x \rrbracket(u) &\triangleq \overline{x} \langle u \rangle \\ \llbracket \lambda x. e \rrbracket(u) &\triangleq u(x, y). \llbracket e \rrbracket(y) \\ \llbracket e_1 \ e_2 \rrbracket(u) &\triangleq \nu y. (\llbracket e_1 \rrbracket(y) \mid \nu x. (\overline{y} \langle x, u \rangle \mid !x(w). \llbracket e_2 \rrbracket(w))) \end{split}$$

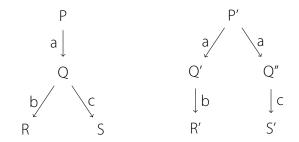
When are two processes equal?

Perhaps the most important contributions of research on π calculus has been the development of the notion of *bisimulation*



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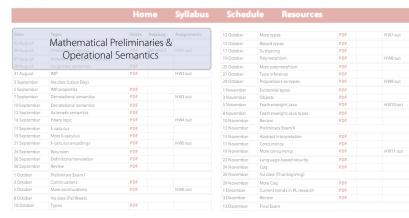


Programming Languages and Logics MWF 9:05-9:55 Gates G01



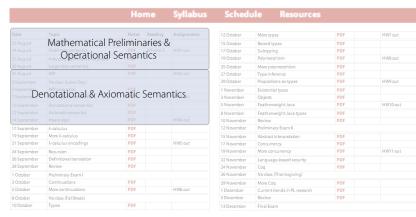
		Home		Syllabus	Schedule Resources			
Date	Topic	Notes	Reading	Assignments	12 October	More types	PDF	HW7 out
22 August	Introduction	PDF	Winskel 1		15 October	Record types	PDF	
24 August	Small-step semantics	PDF	Winskel 2	HW1 out	17 October	Subtyping	PDF	
27 August	Inductive definitions and proofs	PDF			19 October	Polymorphism	PDF	HW8 out
29 August	Large-step semantics	PDF			25 October	More polymorphism	PDF	
31 August	IMP	PDF		HW2 out	27 October	Type inference	PDF	
3 September	No class (Labor Day)				29 October	Proposition s-as-types	PDF	HW9 out
5 September	IMP properties	PDF			1 November	Existential types	PDF	
7 September	Denotational semantics	PDF		HW3 out	3 November	Objects	PDF	
10 September	Denotational semantics	PDF			5 November	Featherweight Java	PDF	HW10 out
12 September	Axiomatic semantics	PDF			8 November	Featherweight Java types	PDF	
14 September	Hoare logic	PDF		HW4 out	10 November	Review	PDF	
17 September	λ-calculus PDF				12 November	Preliminary Exam II		
19 September	More λ -calculus	PDF			15 November	Abstract interpretation	PDF	
21 September	λ-calculus encodings	PDF		HWS out	17 November	Concurrency	PDF	
24 September	Recursion	PDF			19 November	More concurrency	PDF	HW11 out
26 September	Definitional translation	PDF			22 November	Language-based security	PDF	
28 September	Review	PDF			24 November	Coq	PDF	
1 October	Preliminary Exam I				26 November	No class (Thanksgiving)		
3 October	Continuations	PDF			29 November	More Cog	PDF	
5 October	More continuations	PDF		HW6 out	1 December	Current trends in PL research	PDF	
8 October	No class (Fall Break)				3 December	Review	PDF	
10 October	Types PDF				13 December	Final Exam		

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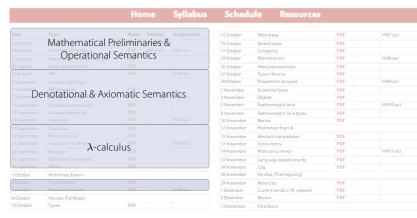


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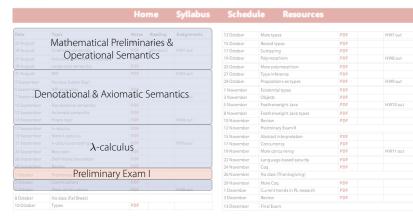


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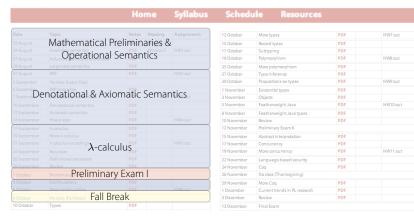


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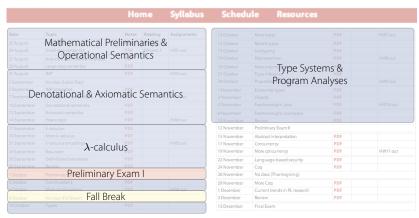


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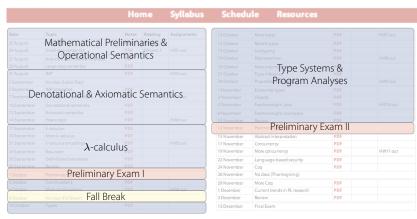




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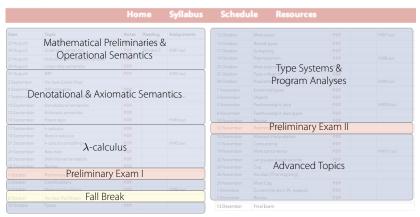


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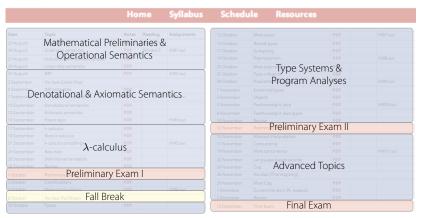


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Programming Languages and Logics MWF 9:05-9:55





Department of Computer Science

Programming Languages and Logics MWF 9:05-9:55 Gates G01





• Mathematical Preliminaries (inductive definitions)

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Expect to solve probems just like the ones we've seen throughout the course...

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Expect to solve probems just like the ones we've seen throughout the course...

...and to apply the skills you've acquired to new problems too!

Final Logistics

- Date: Friday, December 12th
- Time: 9-11:30am
- Where: Gates G01
- Practice: Available today
- Review: Next week?

• CS 6110 – Advanced Programming Languages

- CS 6110 Advanced Programming Languages
- CS 611X Certified Software Systems

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- CS 611X Certified Software Systems
- CS 7190 Seminar in Programming Languages

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Thank you, and stay in touch!