

CS 4110

# Programming Languages & Logics

Lecture 20  
Proving Type Soundness

20 October 2014



# Announcements

---

- Today: Foster office hours 4-5pm
- PS 5 starter code updated last night

# Simply-Typed Lambda Calculus

---

## Syntax

expressions       $e ::= x \mid \lambda x : \tau. e \mid e_1 e_2 \mid n \mid e_1 + e_2 \mid ()$

values             $v ::= \lambda x : \tau. e \mid n \mid ()$

types             $\tau ::= \mathbf{int} \mid \mathbf{unit} \mid \tau_1 \rightarrow \tau_2$

# Simply-Typed Lambda Calculus

## Syntax

expressions       $e ::= x \mid \lambda x : \tau. e \mid e_1 e_2 \mid n \mid e_1 + e_2 \mid ()$

values             $v ::= \lambda x : \tau. e \mid n \mid ()$

types             $\tau ::= \mathbf{int} \mid \mathbf{unit} \mid \tau_1 \rightarrow \tau_2$

## Dynamic Semantics

$E ::= [\cdot] \mid Ee \mid vE \mid E + e \mid v + E$

$$\frac{e \rightarrow e'}{E[e] \rightarrow E[e']}$$

$$(\lambda x : \tau. e) v \rightarrow e\{v/x\}$$

$$\frac{n = n_1 + n_2}{n_1 + n_2 \rightarrow n}$$

# Simply-Typed Lambda Calculus

---

## Static Semantics

# Simply-Typed Lambda Calculus

---

## Static Semantics

$$\frac{}{\Gamma \vdash n : \mathbf{int}} \text{ T-Int}$$

# Simply-Typed Lambda Calculus

---

## Static Semantics

$$\frac{}{\Gamma \vdash n : \mathbf{int}} \text{ T-Int}$$

$$\frac{}{\Gamma \vdash () : \mathbf{unit}} \text{ T-Unit}$$

# Simply-Typed Lambda Calculus

## Static Semantics

$$\frac{}{\Gamma \vdash n : \mathbf{int}} \text{ T-Int}$$

$$\frac{}{\Gamma \vdash () : \mathbf{unit}} \text{ T-Unit}$$

$$\frac{\Gamma \vdash e_1 : \mathbf{int} \quad \Gamma \vdash e_2 : \mathbf{int}}{\Gamma \vdash e_1 + e_2 : \mathbf{int}} \text{ T-Add}$$

# Simply-Typed Lambda Calculus

## Static Semantics

$$\frac{}{\Gamma \vdash n : \mathbf{int}} \text{ T-Int}$$

$$\frac{}{\Gamma \vdash () : \mathbf{unit}} \text{ T-Unit}$$

$$\frac{\Gamma \vdash e_1 : \mathbf{int} \quad \Gamma \vdash e_2 : \mathbf{int}}{\Gamma \vdash e_1 + e_2 : \mathbf{int}} \text{ T-Add}$$

$$\frac{\Gamma(x) = \tau}{\Gamma \vdash x : \tau} \text{ T-Var}$$

# Simply-Typed Lambda Calculus

## Static Semantics

$$\frac{}{\Gamma \vdash n : \mathbf{int}} \text{ T-Int}$$

$$\frac{}{\Gamma \vdash () : \mathbf{unit}} \text{ T-Unit}$$

$$\frac{\Gamma \vdash e_1 : \mathbf{int} \quad \Gamma \vdash e_2 : \mathbf{int}}{\Gamma \vdash e_1 + e_2 : \mathbf{int}} \text{ T-Add}$$

$$\frac{\Gamma(x) = \tau}{\Gamma \vdash x : \tau} \text{ T-Var}$$

$$\frac{\Gamma, x : \tau \vdash e : \tau'}{\Gamma \vdash \lambda x : \tau. e : \tau \rightarrow \tau'} \text{ T-Abs}$$

# Simply-Typed Lambda Calculus

## Static Semantics

$$\frac{}{\Gamma \vdash n : \mathbf{int}} \text{ T-Int}$$

$$\frac{}{\Gamma \vdash () : \mathbf{unit}} \text{ T-Unit}$$

$$\frac{\Gamma \vdash e_1 : \mathbf{int} \quad \Gamma \vdash e_2 : \mathbf{int}}{\Gamma \vdash e_1 + e_2 : \mathbf{int}} \text{ T-Add}$$

$$\frac{\Gamma(x) = \tau}{\Gamma \vdash x : \tau} \text{ T-Var}$$

$$\frac{\Gamma, x : \tau \vdash e : \tau'}{\Gamma \vdash \lambda x : \tau. e : \tau \rightarrow \tau'} \text{ T-Abs}$$

$$\frac{\Gamma \vdash e_1 : \tau \rightarrow \tau' \quad \Gamma \vdash e_2 : \tau}{\Gamma \vdash e_1 e_2 : \tau'} \text{ T-App}$$

# Properties

## Theorem (Type soundness)

If  $\vdash e : \tau$  and  $e \rightarrow^* e'$  and  $e' \not\rightarrow$  then  $e'$  is a value and  $\vdash e' : \tau$ .

# Properties

## Theorem (Type soundness)

If  $\vdash e : \tau$  and  $e \rightarrow^* e'$  and  $e' \not\rightarrow$  then  $e'$  is a value and  $\vdash e' : \tau$ .

## Lemma (Preservation)

If  $\vdash e : \tau$  and  $e \rightarrow e'$  then  $\vdash e' : \tau$ .

# Properties

## Theorem (Type soundness)

If  $\vdash e : \tau$  and  $e \rightarrow^* e'$  and  $e' \not\rightarrow$  then  $e'$  is a value and  $\vdash e' : \tau$ .

## Lemma (Preservation)

If  $\vdash e : \tau$  and  $e \rightarrow e'$  then  $\vdash e' : \tau$ .

## Lemma (Progress)

If  $\vdash e : \tau$  then either  $e$  is a value or there exists an  $e'$  such that  $e \rightarrow e'$ .