Pipelining and Hazards

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CS 3410, Spring 2015

Computer Science

Cornell University

See P&H Chapter: 4.6-4.8

Announcements

Prelim next week
- Tuesday at 7:30.
- Go to location based on netid
  - [a-g]* → MRS146: Morrison Hall 146
  - [h-l]* → RRB125: Riley-Robb Hall 125
  - [m-n]* → RRB105: Riley-Robb Hall 105
  - [o-s]* → MVRG71: M Van Rensselaer Hall G71
  - [t-z]* → MVRG73: M Van Rensselaer Hall G73

Prelim reviews
- TODAY, Tue, Feb 24 @ 7:30pm in Olin 255
- Sat, Feb 28 @ 7:30pm in Upson B17

Prelim conflicts
- Contact Deniz Altinbuken <deniz@cs.cornell.edu>

Announcements

Prelim1:

- Time: We will start at 7:30pm sharp, so come early
- Location: on previous slide
- Closed Book
  - Cannot use electronic device or outside material
- Practice prelims are online in CMS

- Material covered everything up to end of this week
  - Everything up to and including data hazards
  - Appendix B (logic, gates, FSMs, memory, ALUs)
  - Chapter 4 (pipelined [and non] MIPS processor with hazards)
  - Chapters 2 (Numbers / Arithmetic, simple MIPS instructions)
  - Chapter 1 (Performance)
  - HW1, Lab0, Lab1, Lab2, C-Lab0, C-Lab1

Goals for Today

RISC and Pipelined Processor: Putting it all together

Data Hazards

- Data dependencies
- Problem, detection, and solutions
  - (delaying, stalling, forwarding, bypass, etc)
- Hazard detection unit
- Forwarding unit

Next time

- Control Hazards

  What is the next instruction to execute if a branch is taken? Not taken?

MIPS Design Principles

Simplicity favors regularity

- 32 bit instructions

Smaller is faster

- Small register file

Make the common case fast

- Include support for constants

Good design demands good compromises

- Support for different type of interpretations/classes

Recall: MIPS instruction formats

All MIPS instructions are 32 bits long, has 3 formats

R-type

op	rs	rt	rd	shamt	func
6 bits	5 bits	5 bits	5 bits	5 bits	6 bits

I-type

op	rs	rt	immediate
6 bits	5 bits	5 bits	16 bits

J-type

op	immediate (target address)
6 bits	26 bits

Recall: MIPS Instruction Types

Arithmetic/Logical

- R-type: result and two source registers, shift amount
- I-type: 16-bit immediate with sign/zero extension

Memory Access

- load/store between registers and memory
- word, half-word and byte operations

Control flow

- conditional branches: pc-relative addresses
- jumps: fixed offsets, register absolute

Recall: MIPS Instruction Types

Arithmetic/Logical

- ADD, ADDU, SUB, SUBU, AND, OR, XOR, NOR, SLT, SLTU
- ADDI, ADDIU, ANDI, ORI, XORI, LUI, SLL, SRL, SLLV, SRLV, SRAV, SLTI, SLTIU
- MULT, DIV, MFLO, MTLO, MFHI, MTHI

Memory Access

- LW, LH, LB, LHU, LBU, LWL, LWR
- SW, SH, SB, SWL, SWR

Control flow

- BEQ, BNE, BLEZ, BLTZ, BGEZ, BGTZ
- J, JR, JAL, JALR, BEQL, BNEL, BLEZL, BGTZL

Special

- LL, SC, SYSCALL, BREAK, SYNC, COPROC

Pipelining

Principle:

Throughput increased by parallel execution

Balanced pipeline very important

Else slowest stage dominates performance

Pipelining:

- Identify *pipeline stages*
- Isolate stages from each other
- Resolve pipeline *hazards* (this and next lecture)

Basic Pipeline

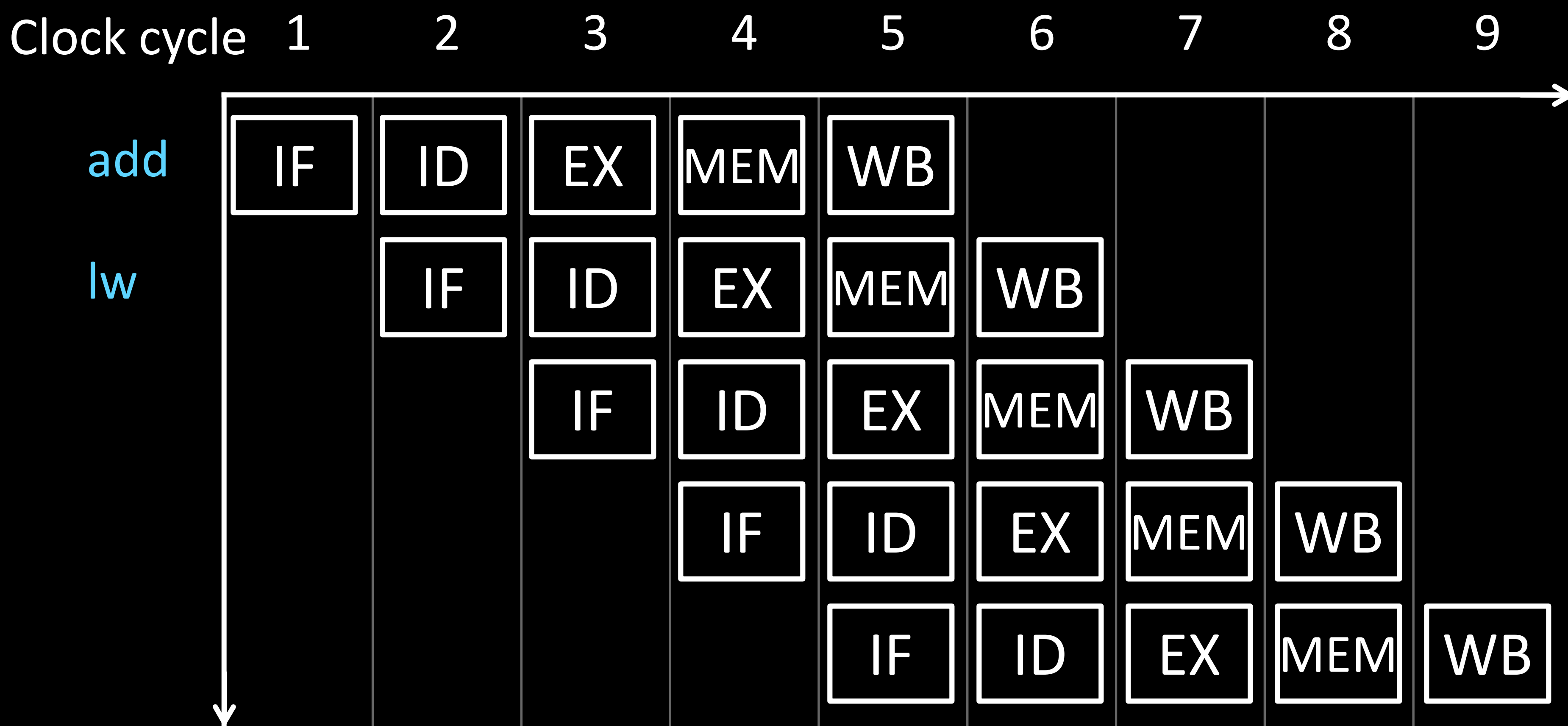
Five stage "RISC" load-store architecture

1. Instruction fetch (IF)
   - get instruction from memory, increment PC
2. Instruction Decode (ID)
   - translate opcode into control signals and read registers
3. Execute (EX)
   - perform ALU operation, compute jump/branch targets
4. Memory (MEM)
   - access memory if needed
5. Writeback (WB)
   - update register file

Pipelined Implementation

- Each instruction goes through the 5 stages
  - Each stage takes one clock cycle
    - So slowest stage determines clock cycle time

Time Graphs

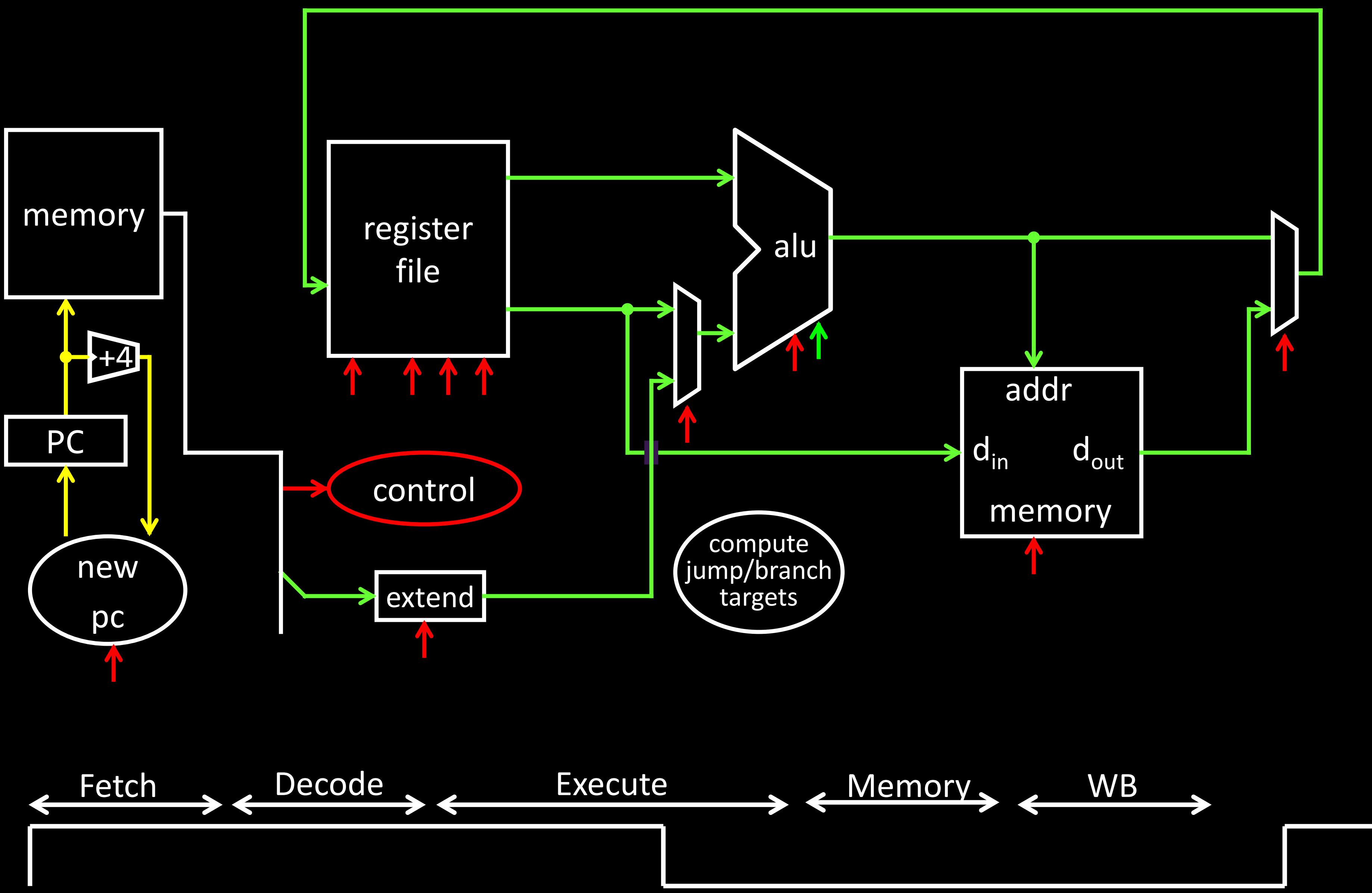

Clock cycle
1
2
3
4
5
6
7
8
9
add
lw
IF ID EX MEM WB
IF ID EX MEM WB
IF ID EX MEM WB
IF ID EX MEM WB
IF ID EX MEM WB

Pipelined Implementation

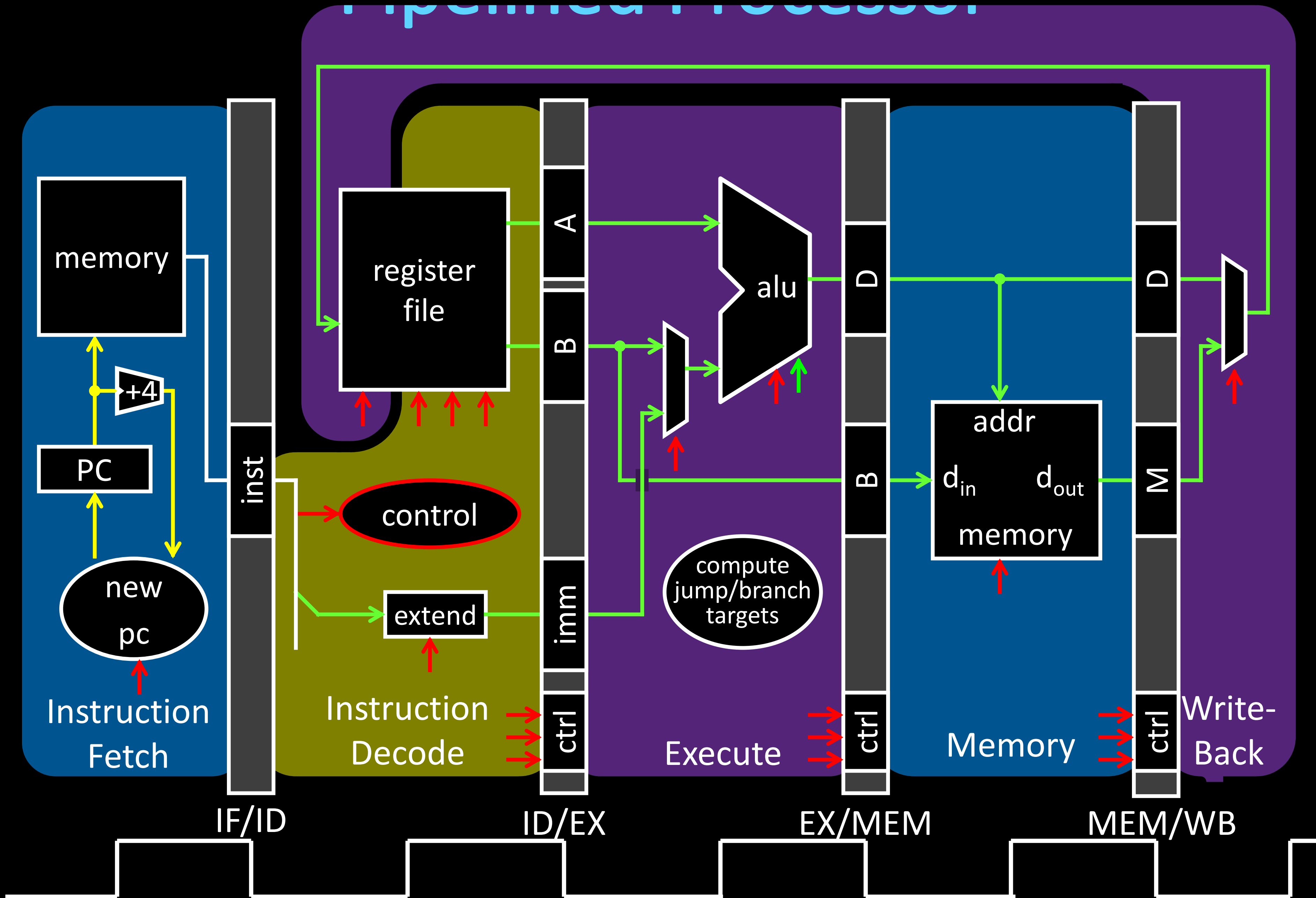
- Each instruction goes through the 5 stages
    - Each stage takes one clock cycle
        - So slowest stage determines clock cycle time

- Stages must share information. How?
    - Add pipeline registers (flip-flops) to pass results between different stages

Pipelined Processor



Pipelined Processor



Pipelined Implementation

- Each instruction goes through the 5 stages
  - Each stage takes one clock cycle
    - So slowest stage determines clock cycle time
- Stages must share information. How?
  - Add pipeline registers (flip-flops) to pass results between different stages

And is this it?

Not quite....

Hazards

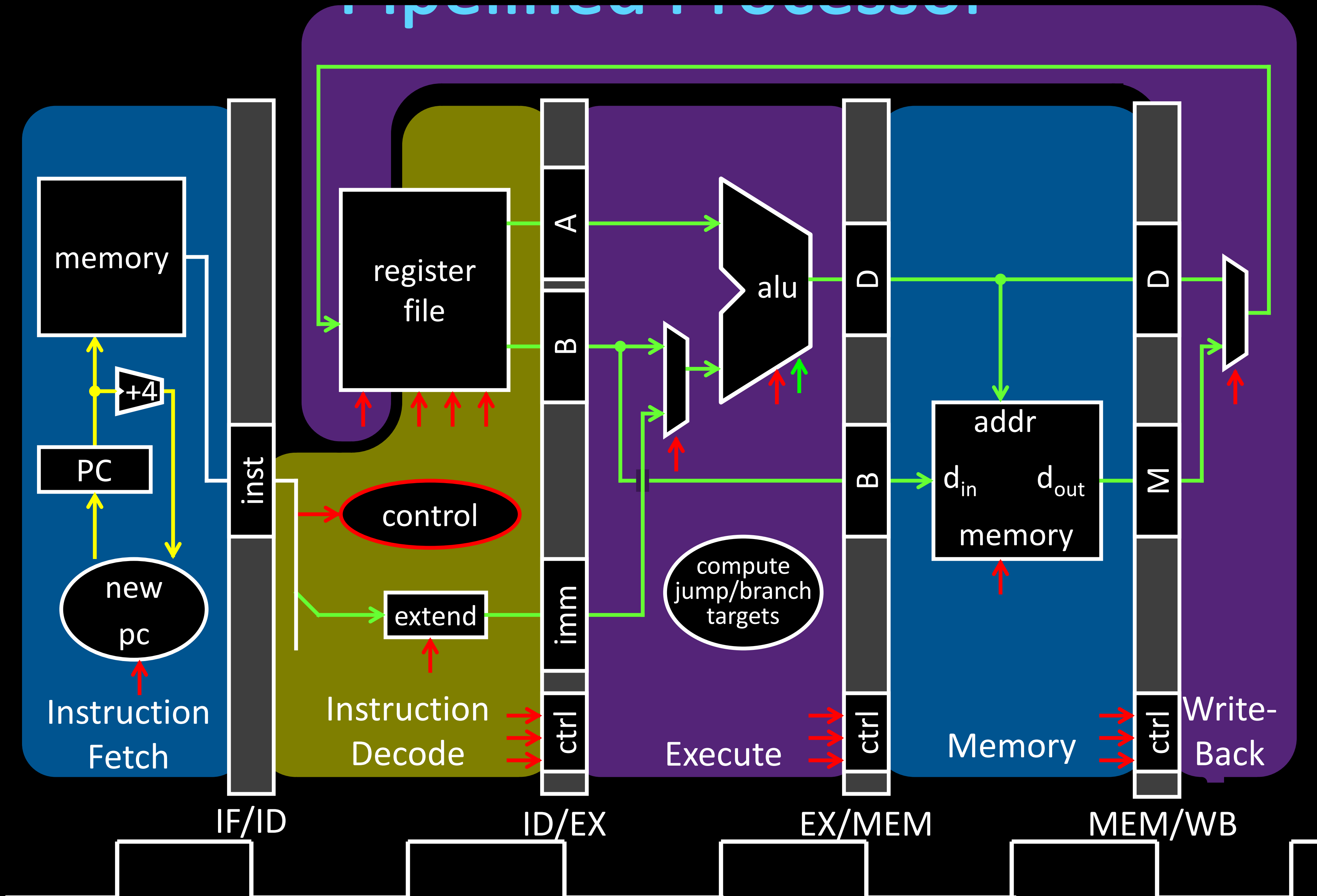
3 kinds

- Structural hazards
  - Multiple instructions want to use same unit
- Data hazards
  - Results of instruction needed before ready
- Control hazards
  - Don’t know which side of branch to take

Will get back to this

First, how to pipeline when no hazards

Pipelined Processor



Example: : Sample Code (Simple)

```
add     r3, r1, r2;
nand    r6, r4, r5;
lw      r4, 20(r2);
add     r5, r2, r5;
sw      r7, 12(r3);
```

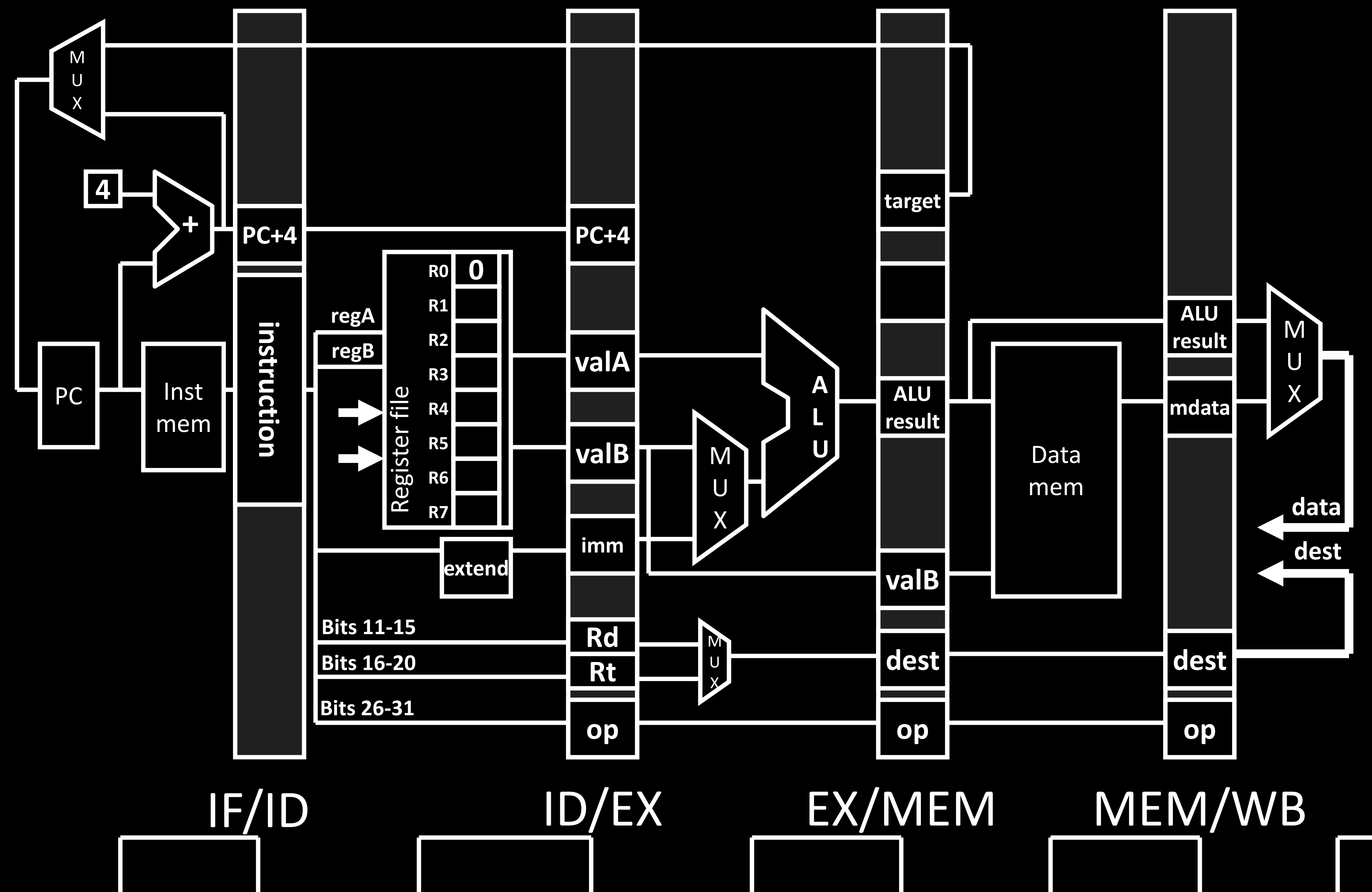
Example: Sample Code (Simple)

Assume eight-register machine

Run the following code on a pipelined datapath

```
add   r3  r1    r2    ;  reg 3 = reg 1 + reg 2
nand  r6  r4    r5    ;  reg 6 = ~(reg 4 & reg 5)
lw    r4  20 (r2)     ;  reg 4 =  Mem[reg2+20]
add   r5  r2    r5    ;  reg 5 = reg 2 + reg 5
sw    r7  12(r3)      ;  Mem[reg3+12] = reg 7
```

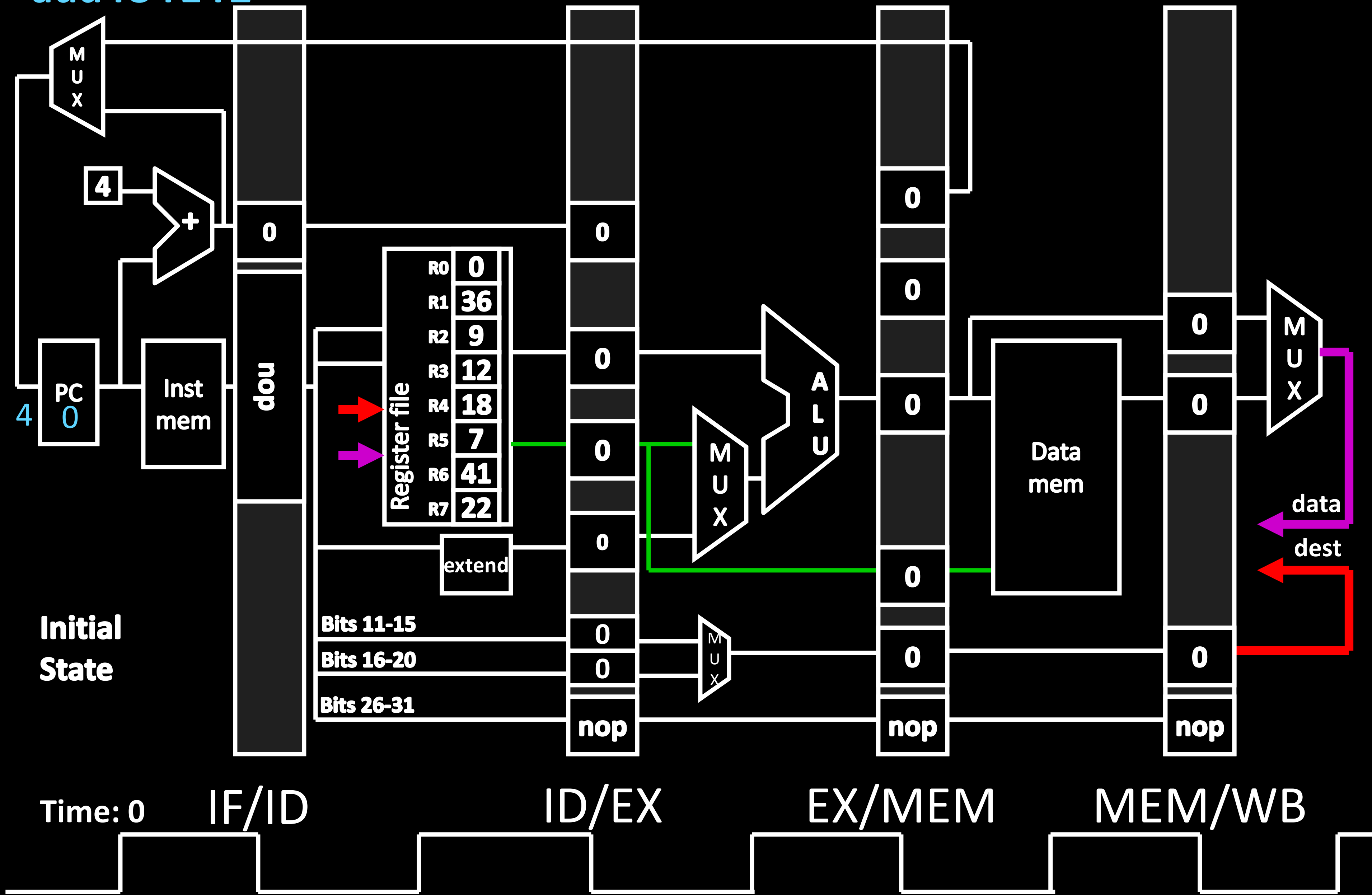
Slides thanks to Sally McKee



MUX
4
+
PC+4
PC
Inst mem
instruction
regA
regB
Register file
R0
0
R1
R2
R3
R4
R5
R6
R7
extend
Bits 11-15
Bits 16-20
Bits 26-31
PC+4
valA
valB
imm
Rd
Rt
op
MUX
ALU
MUX
target
ALU result
valB
dest
op
Data mem
ALU result
mdata
dest
op
MUX
data
dest
IF/ID
ID/EX
EX/MEM
MEM/WB

At time 1,
Fetch

add r3 r1 r2



Takeaway

Pipelining is a powerful technique to mask latencies and increase throughput

- Logically, instructions execute one at a time
- Physically, instructions execute in parallel
  - Instruction level parallelism

Abstraction promotes decoupling

- Interface (ISA) vs. implementation (Pipeline)

Hazards

See P&H Chapter: 4.7-4.8

Hazards

3 kinds

- Structural hazards
  - Multiple instructions want to use same unit
- Data hazards
  - Results of instruction needed before
- Control hazards
  - Don’t know which side of branch to take

Next Goal

What about **data dependencies** (also known as a **data hazard** in a pipelined processor)?

i.e. add r3, r1, r2

    sub r5, r3, r4

Need to detect and then fix such hazards

Data Hazards

Data Hazards

- register file reads occur in stage 2 (ID)
- register file writes occur in stage 5 (WB)
- next instructions may read values about to be written
  - i.e instruction may need values that are being computed further down the pipeline
  - *in fact, this is quite common*

Data Hazards

time →

Clock cycle

1 2 3 4 5 6 7 8 9

add r3, r1, r2

sub r5, r3, r4

lw r6, 4(r3)

or r5, r3, r5

sw r6, 12(r3)

Data Hazards

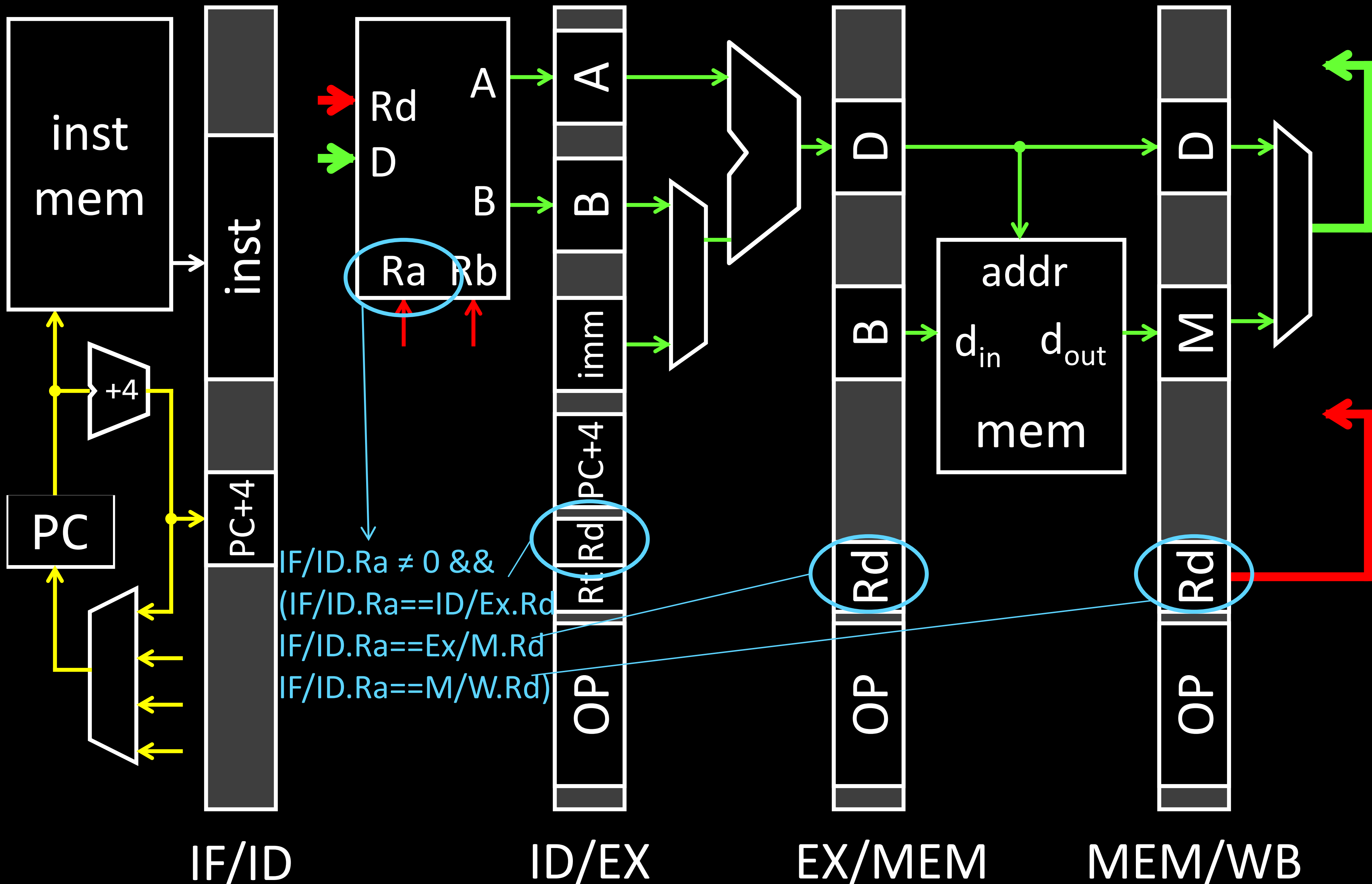
Data Hazards

- register file reads occur in stage 2 (ID)
- register file writes occur in stage 5 (WB)
- next instructions may read values about to be written

i.e. add r3, r1, r2
sub r5, r3, r4

How to detect?

Detecting Data Hazards



Data Hazards

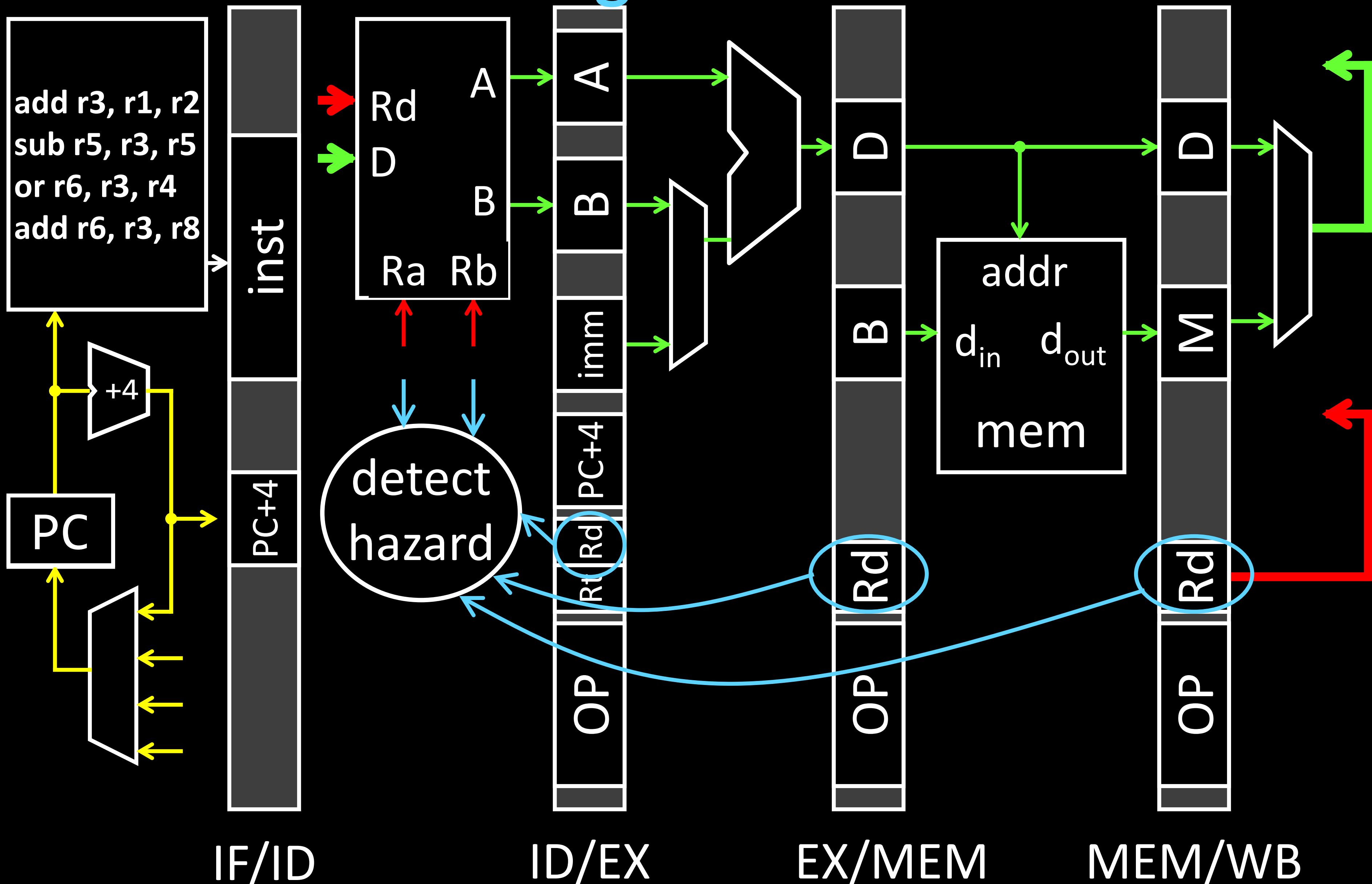
Data Hazards

- register file reads occur in stage 2 (ID)
- register file writes occur in stage 5 (WB)
- next instructions may read values about to be written

How to detect? Logic in ID stage:

```
stall = (IF/ID.Ra != 0 &&
            (IF/ID.Ra == ID/EX.Rd ||
             IF/ID.Ra == EX/M.Rd ||
             IF/ID.Ra == M/WB.Rd))
        || (same for Rb)
```

Detecting Data Hazards



Takeaway

Data hazards occur when a operand (register) depends on the result of a previous instruction that may not be computed yet. A pipelined processor needs to detect data hazards.

Next Goal

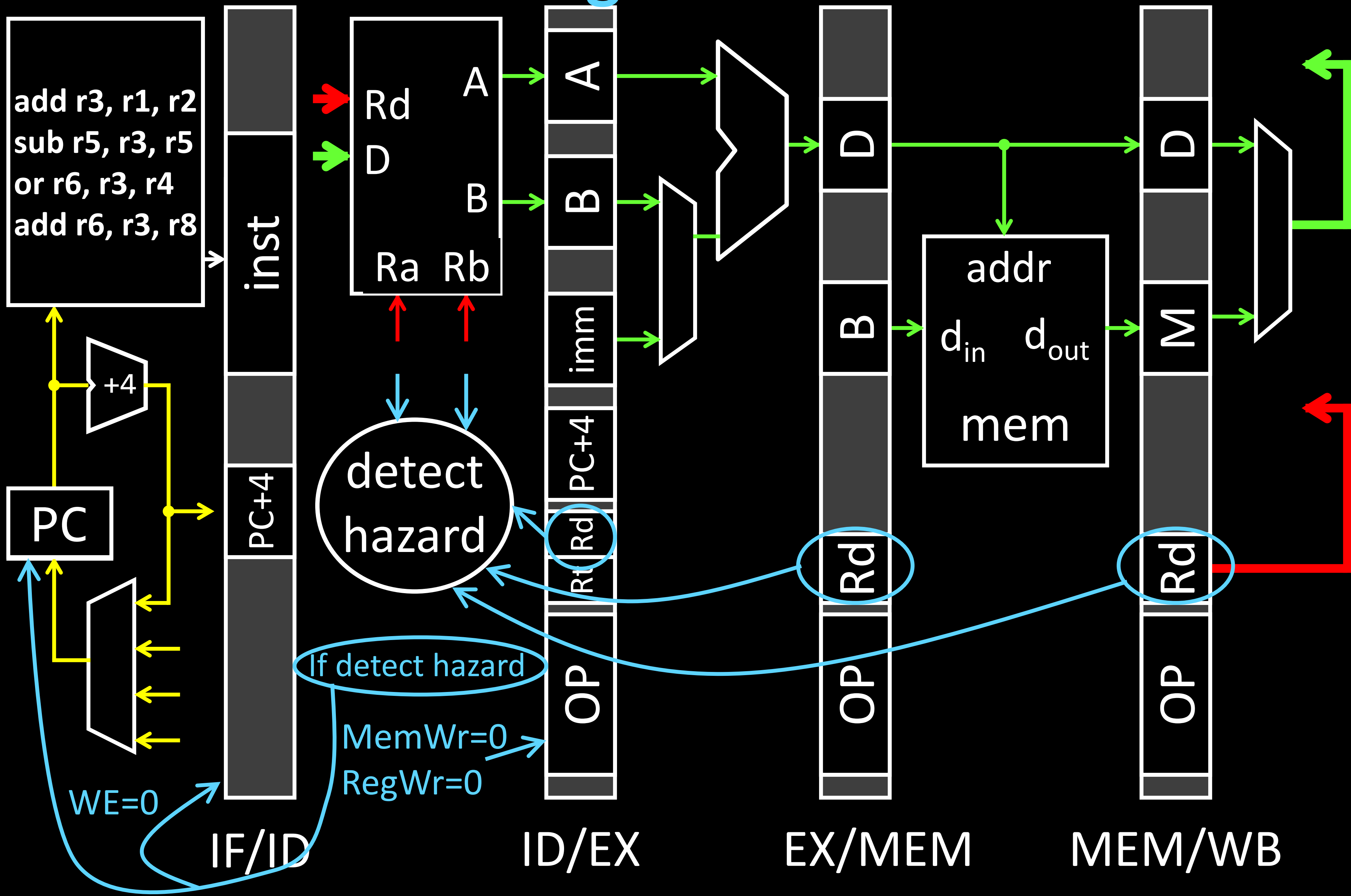
What to do if data hazard detected?

Stalling

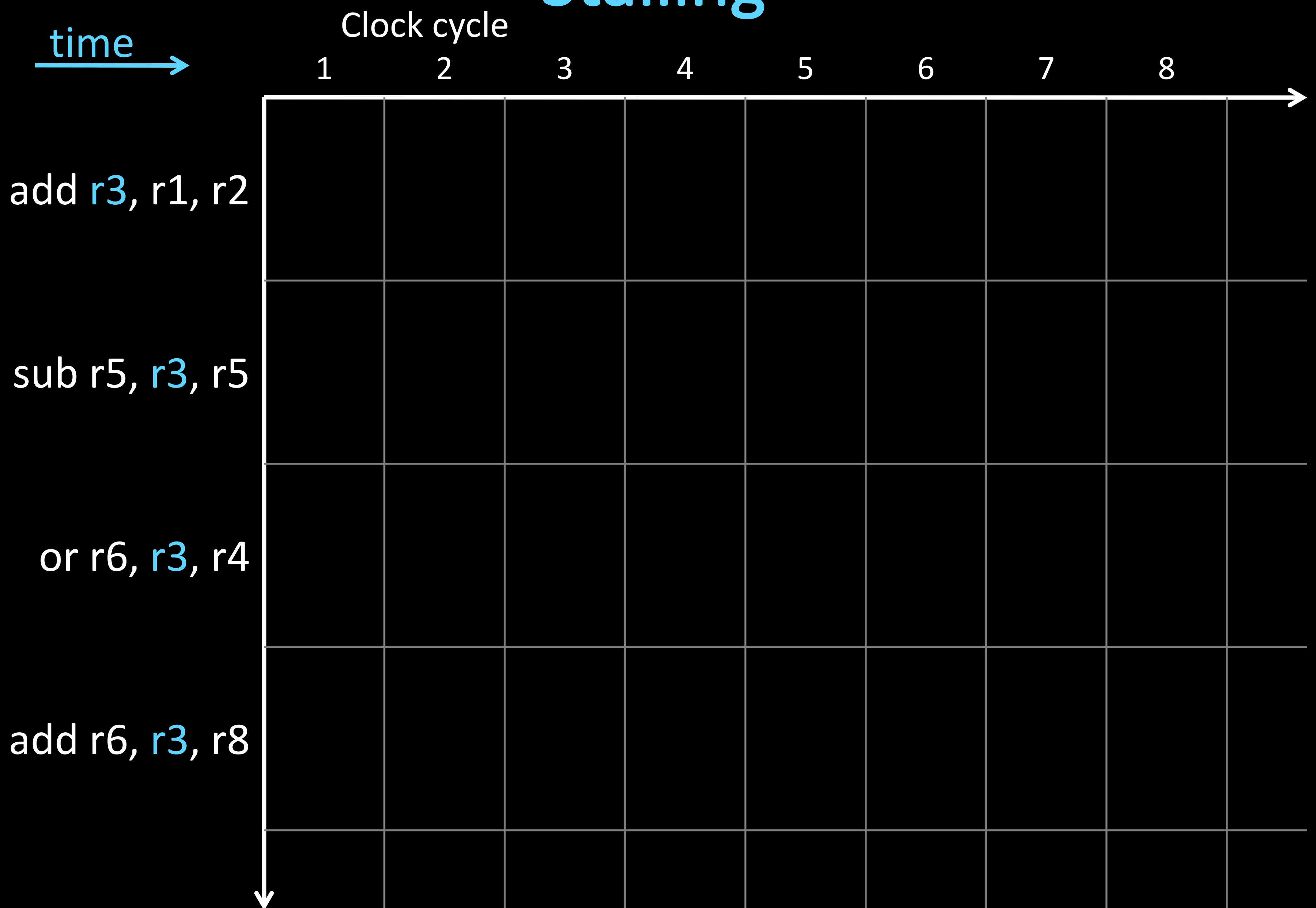
How to stall an instruction in ID stage

- prevent IF/ID pipeline register update
  - stalls the ID stage instruction
- convert ID stage instr into nop for later stages
  - innocuous “bubble” passes through pipeline
- prevent PC update
  - stalls the next (IF stage) instruction

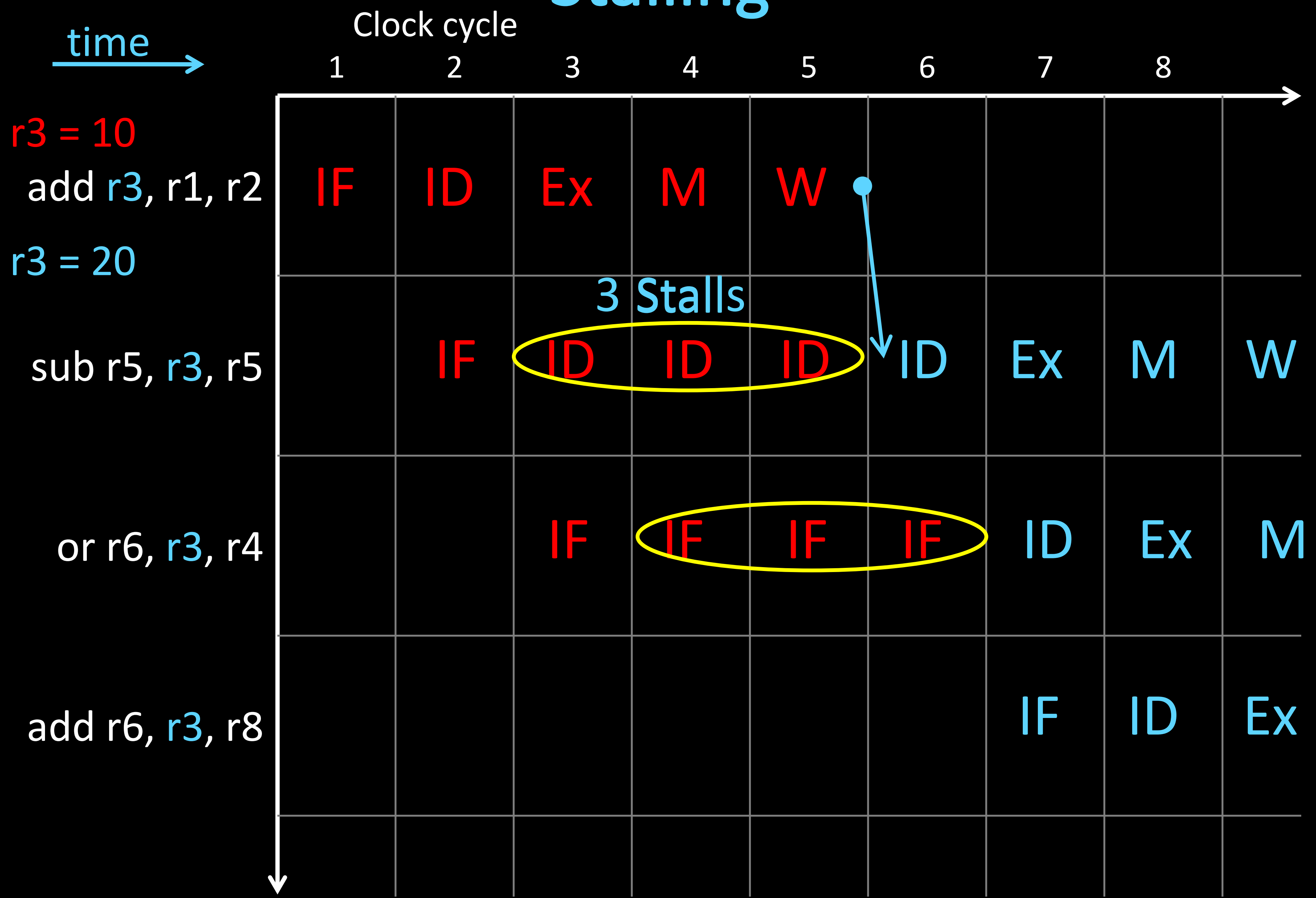
Detecting Data Hazards



Stalling

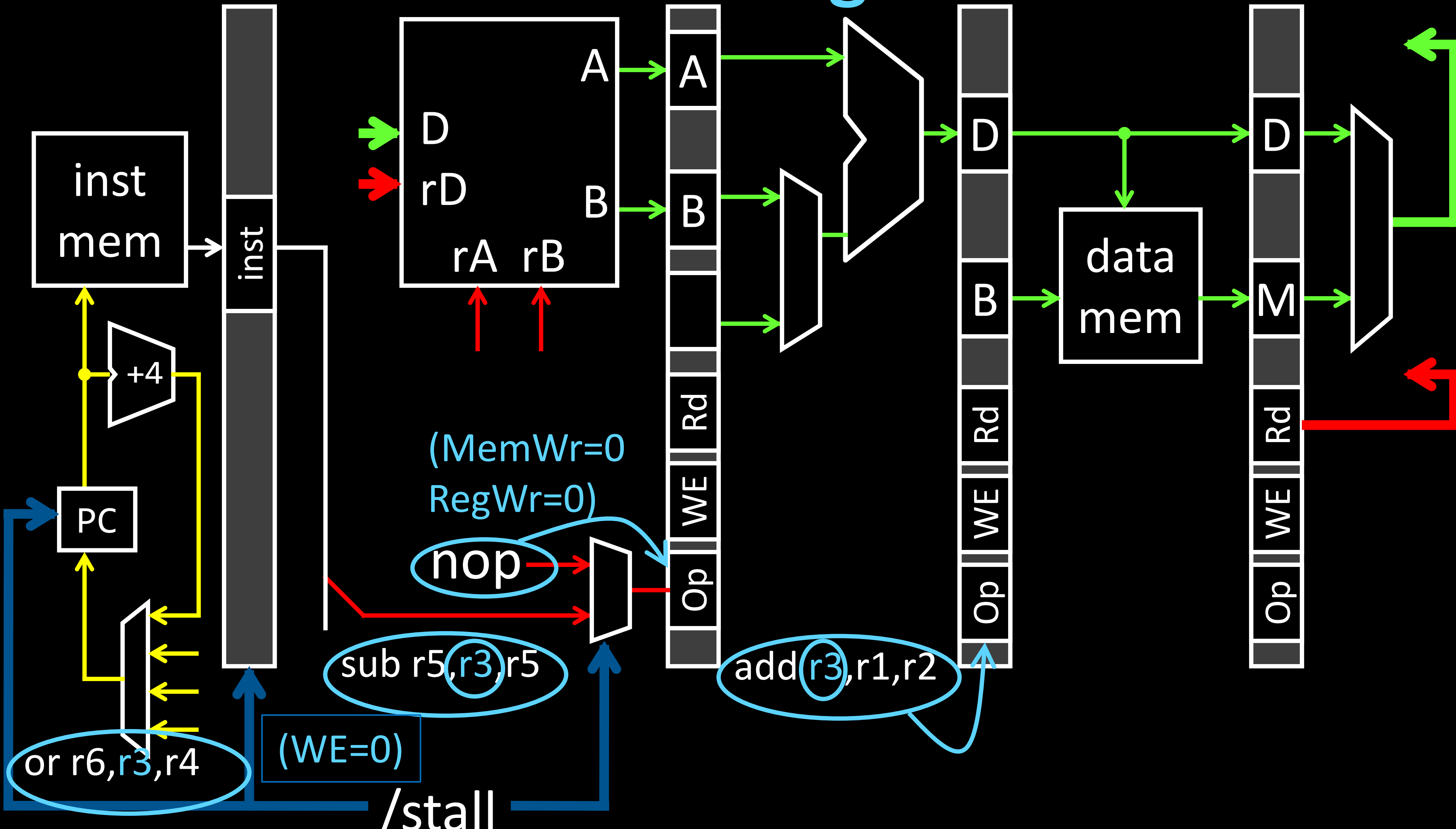

time
Clock cycle
1
2
3
4
5
6
7
8
add r3, r1, r2
sub r5, r3, r5
or r6, r3, r4
add r6, r3, r8

Stalling



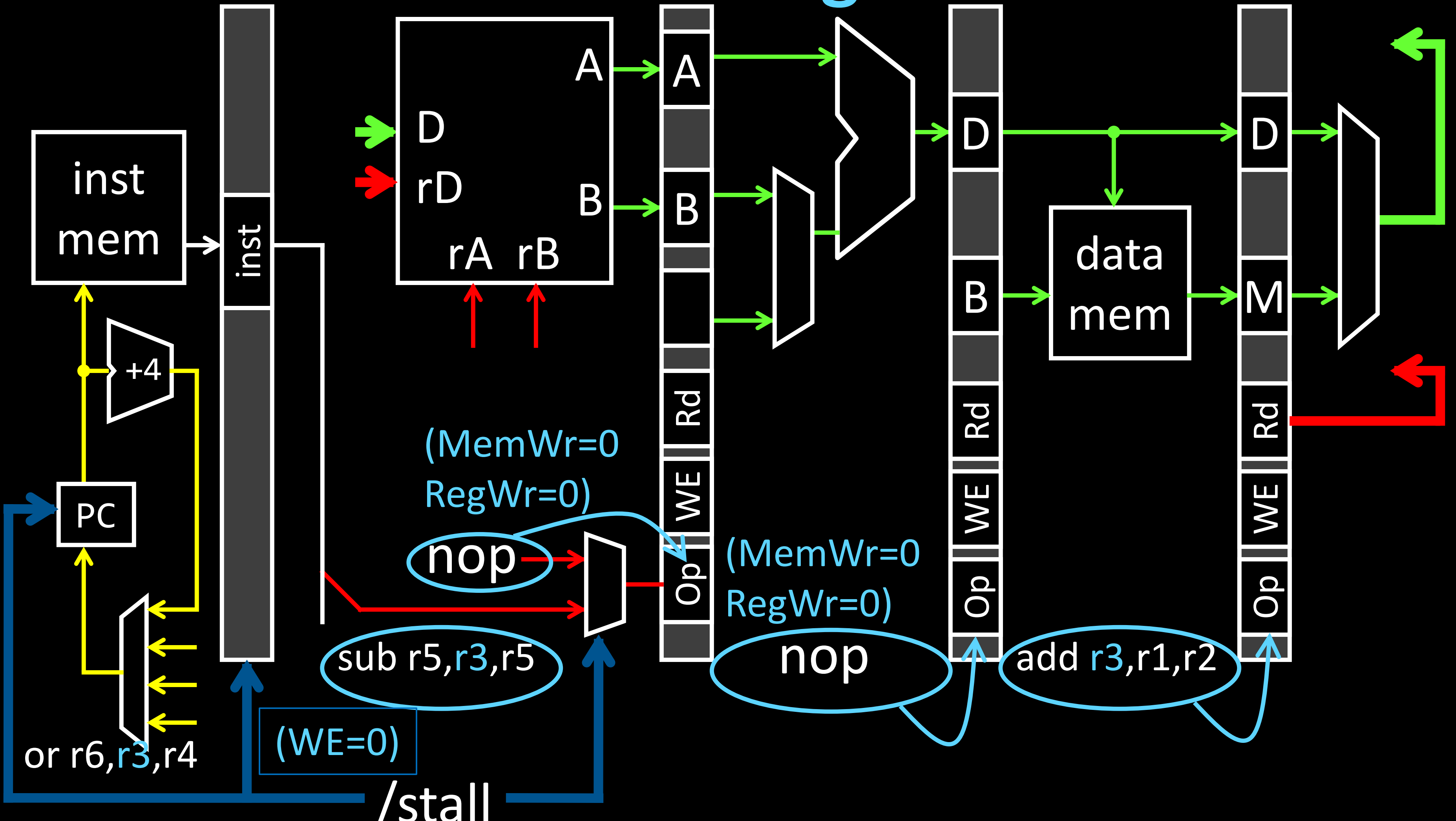
time
Clock cycle
1
2
3
4
5
6
7
8
r3 = 10
add r3, r1, r2
IF
ID
Ex
M
W
r3 = 20
3 Stalls
sub r5, r3, r5
IF
ID
ID
ID
ID
Ex
M
W
or r6, r3, r4
IF
IF
IF
IF
ID
Ex
M
add r6, r3, r8
IF
ID
Ex

Stalling



NOP = If(IF/ID.rA ≠ 0 &&
(IF/ID.rA==ID/Ex.Rd
IF/ID.rA==Ex/M.Rd
IF/ID.rA==M/W.Rd))

Stalling



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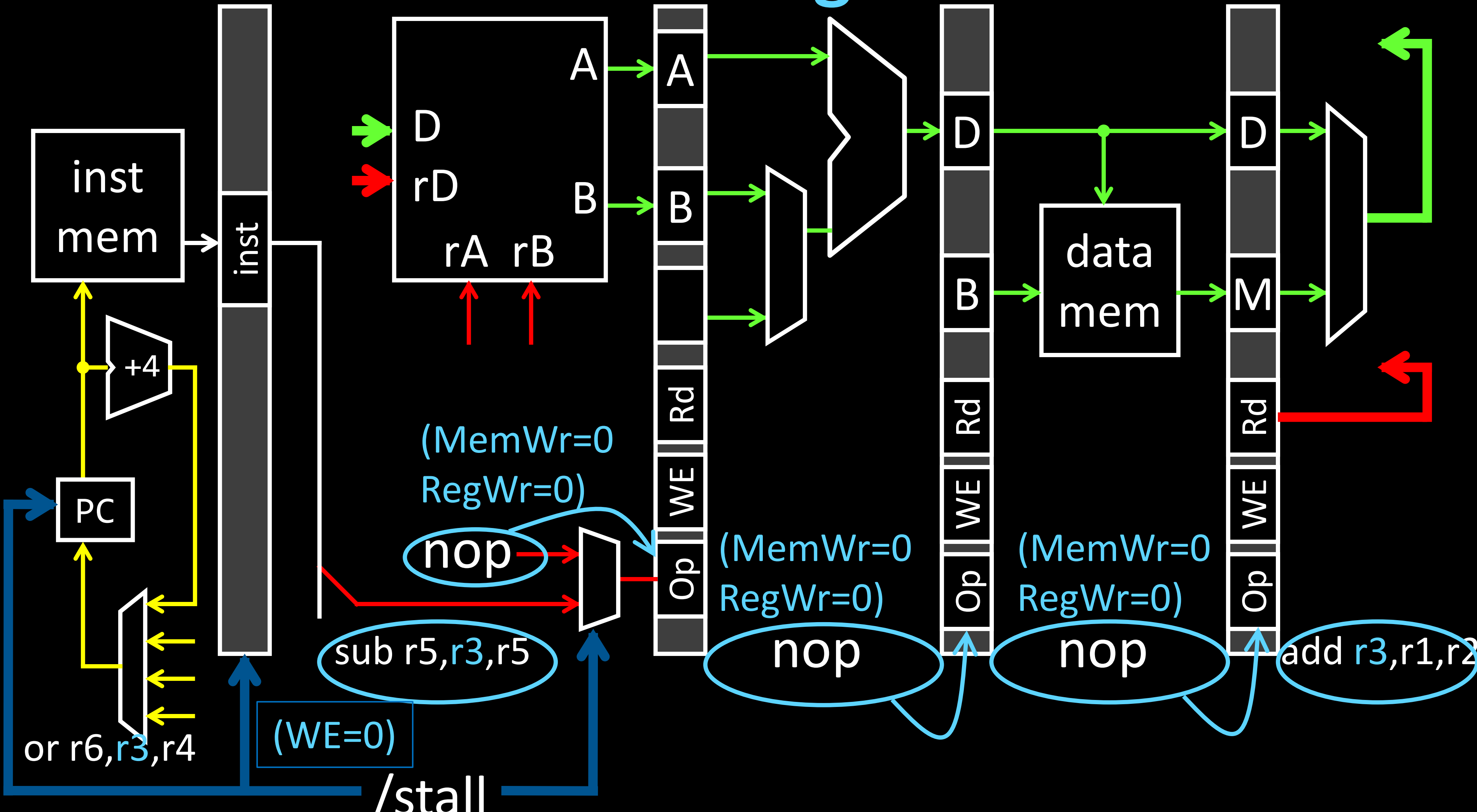

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Stalling

How to stall an instruction in ID stage

- prevent IF/ID pipeline register update
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- convert ID stage instr into nop for later stages
  - innocuous “bubble” passes through pipeline
- prevent PC update
  - stalls the next (IF stage) instruction

Takeaway

Data hazards occur when a operand (register) depends on the result of a previous instruction that may not be computed yet. A pipelined processor needs to detect data hazards.

Stalling, preventing a dependent instruction from advancing, is one way to resolve data hazards.

Stalling introduces NOPs (“bubbles”) into a pipeline. Introduce NOPs by (1) preventing the PC from updating, (2) preventing writes to IF/ID registers from changing, and (3) preventing writes to memory and register file.

*Bubbles in pipeline significantly decrease performance.

Next Goal: Resolving Data Hazards via Forwarding

What to do if data hazard detected?

A) Wait/Stall

B) Reorder in Software (SW)

C) Forward/Bypass

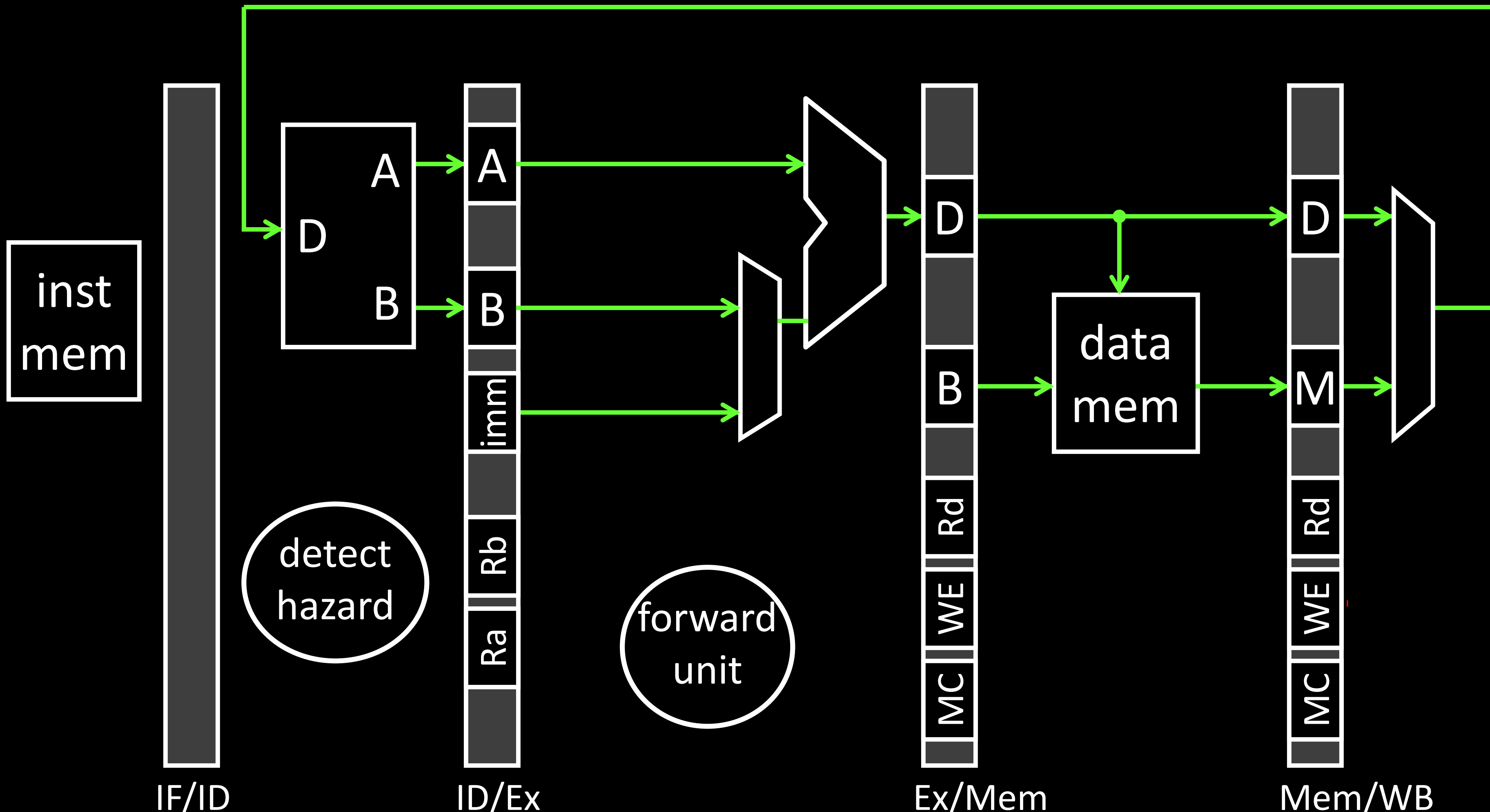
Forwarding

Forwarding bypasses some pipelined stages forwarding a result to a dependent instruction operand (register).

Three types of forwarding/bypass

- Forwarding from Ex/Mem registers to Ex stage (M→Ex)
- Forwarding from Mem/WB register to Ex stage (W→Ex)
- RegisterFile Bypass

Forwarding Datapath



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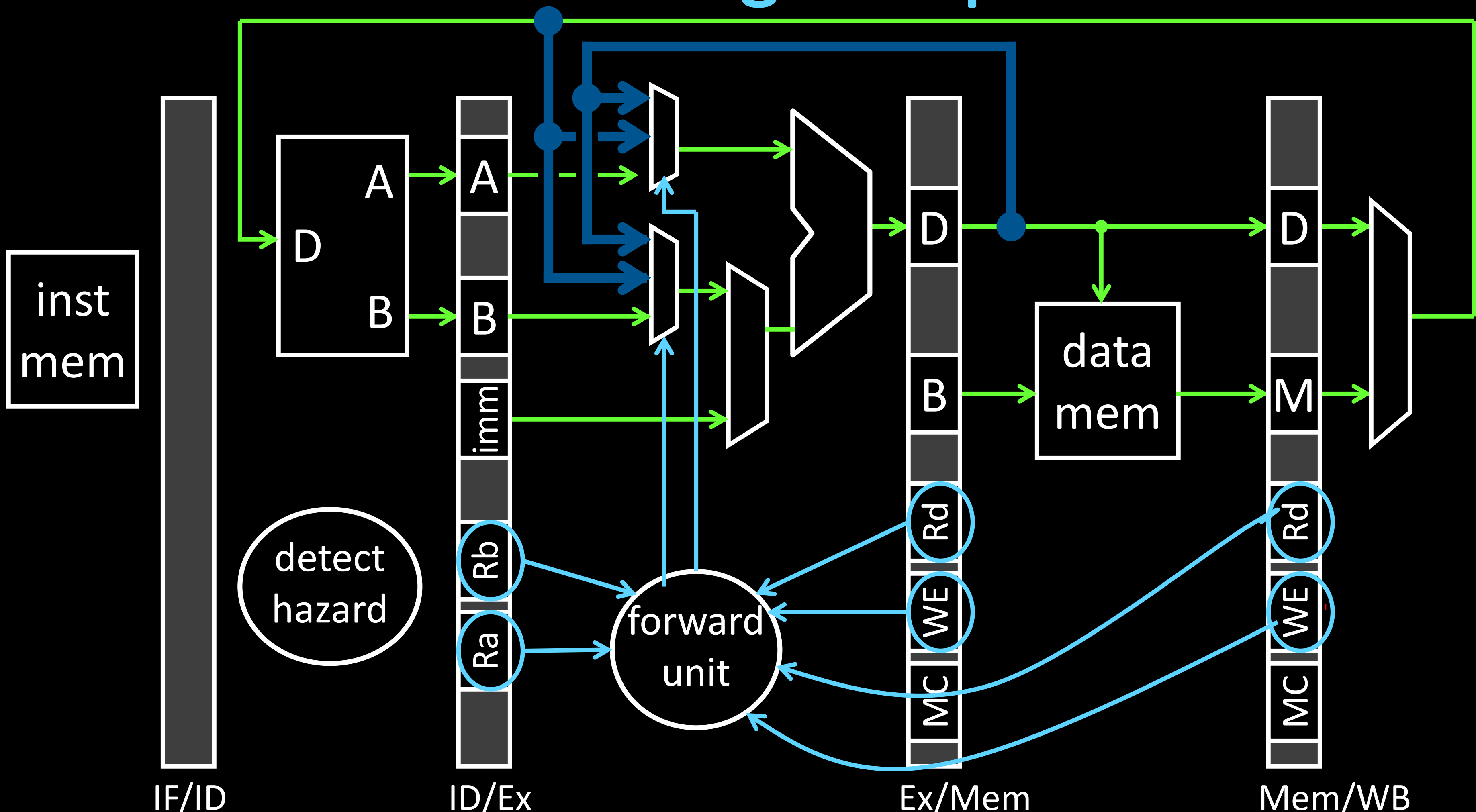


Three types of forwarding/bypass

- Forwarding from Ex/Mem registers to Ex stage (M→Ex)
- Forwarding from Mem/WB register to Ex stage (W → Ex)
- RegisterFile Bypass

Forwarding Datapath 1

Ex/MEM to EX Bypass

- EX needs ALU result that is still in MEM stage
- Resolve:

  Add a bypass from EX/MEM.D to start of EX

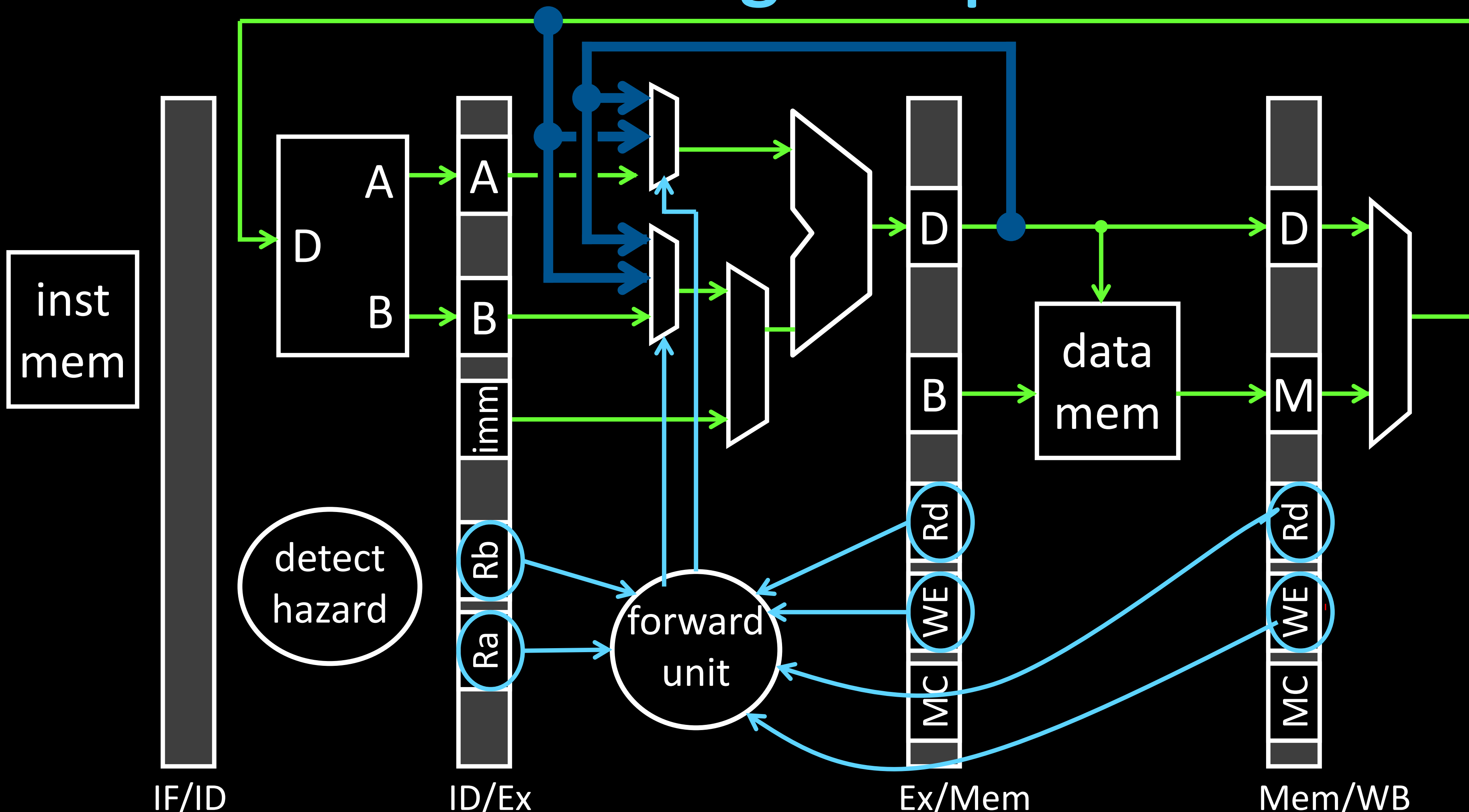
How to detect? Logic in Ex Stage:

forward = (Ex/M.WE && EX/M.Rd != 0 &&

ID/Ex.Ra == Ex/M.Rd)

|| (same for Rb)

Forwarding Datapath



Three types of forwarding/bypass

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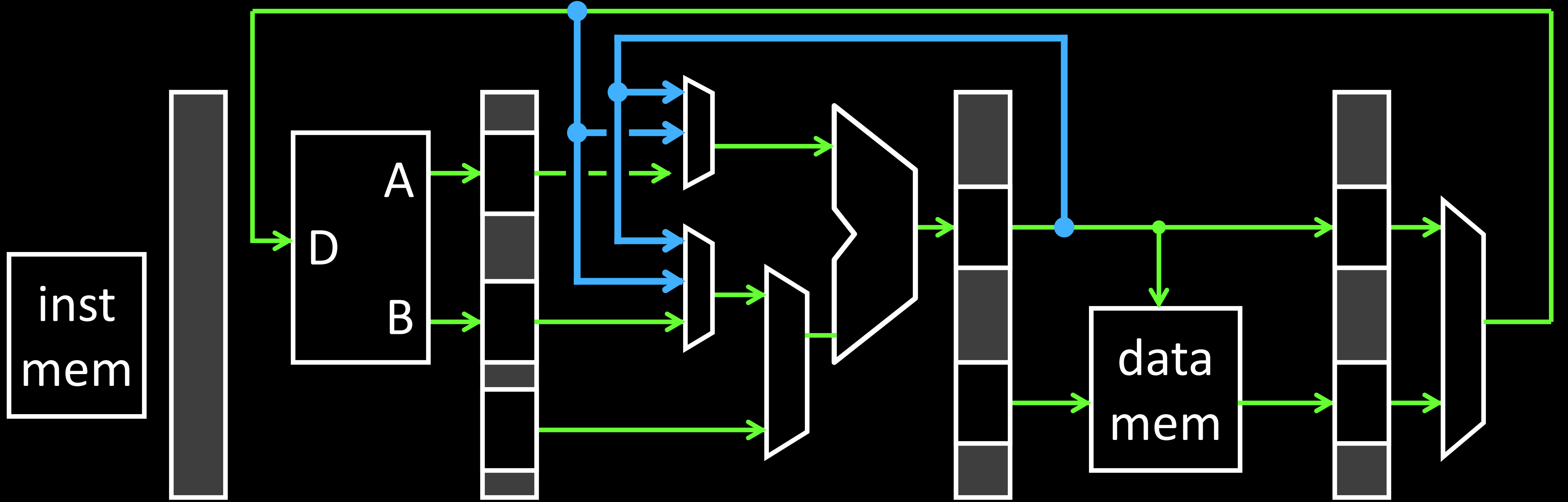
Forwarding Datapath 1



add r3, r1, r2

sub r5, r3, r1

Forwarding Datapath 2

Mem/WB to EX Bypass

- EX needs value being written by WB
- Resolve:

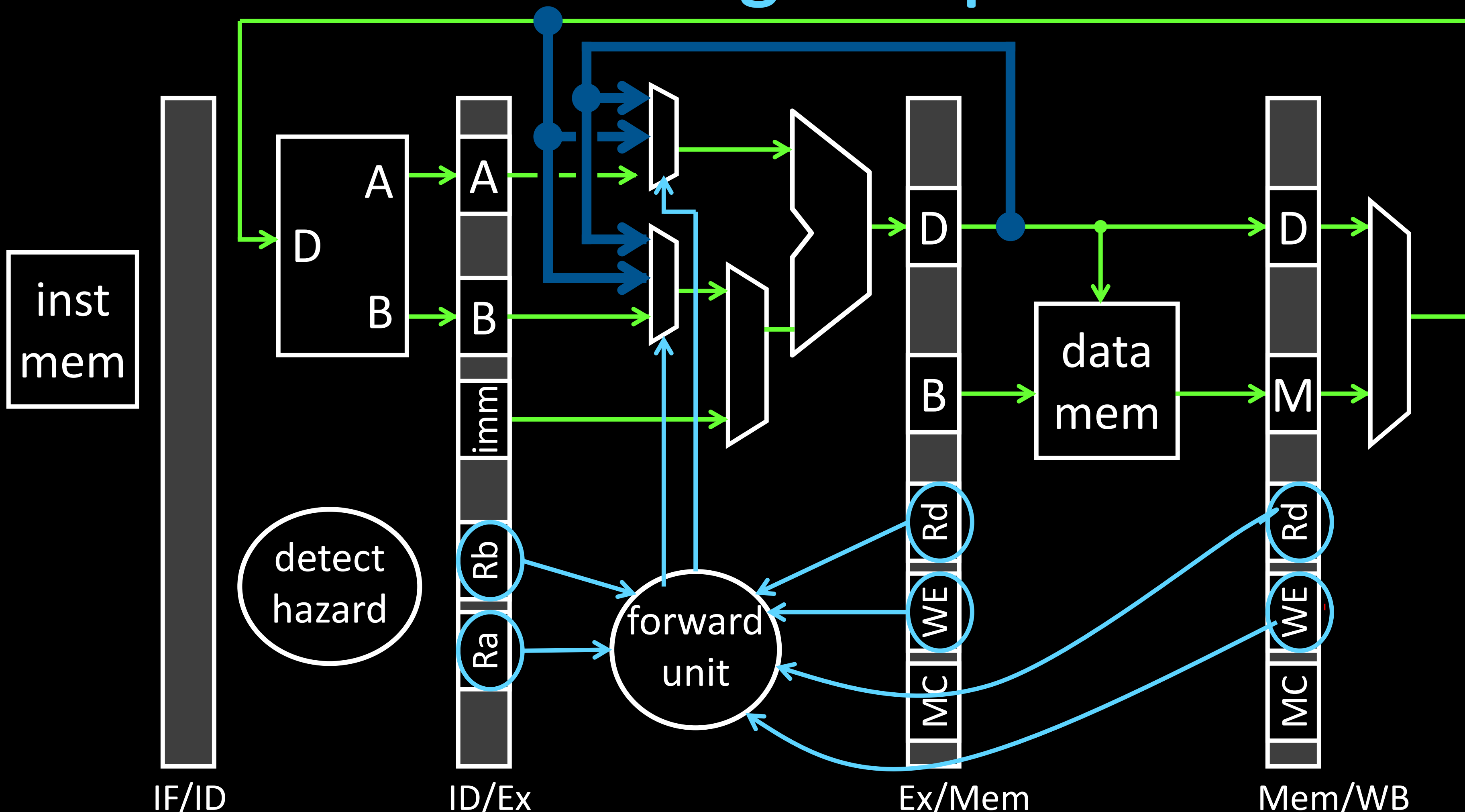
  Add bypass from WB final value to start of EX

How to detect? Logic in Ex Stage:

```
forward = (M/WB.WE && M/WB.Rd != 0 &&
           ID/Ex.Ra == M/WB.Rd &&
           not (Ex/M.WE && Ex/M.Rd != 0 &&
                ID/Ex.Ra == Ex/M.Rd)
        || (same for Rb)
```

Check pg. 311

Forwarding Datapath



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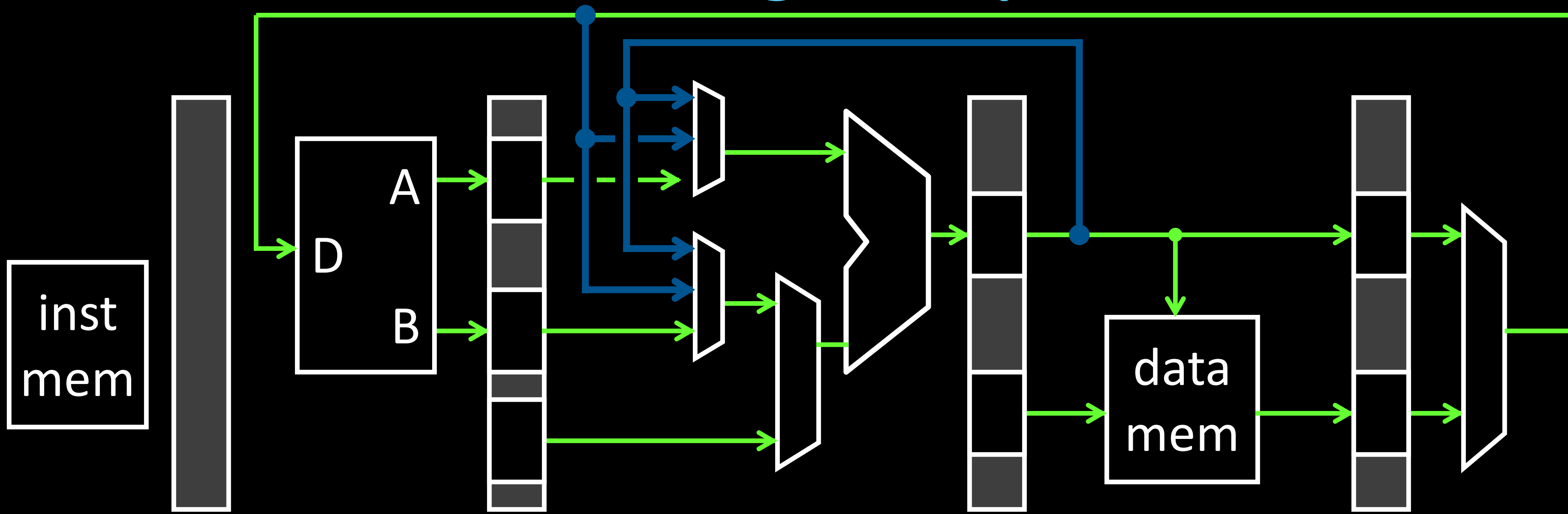
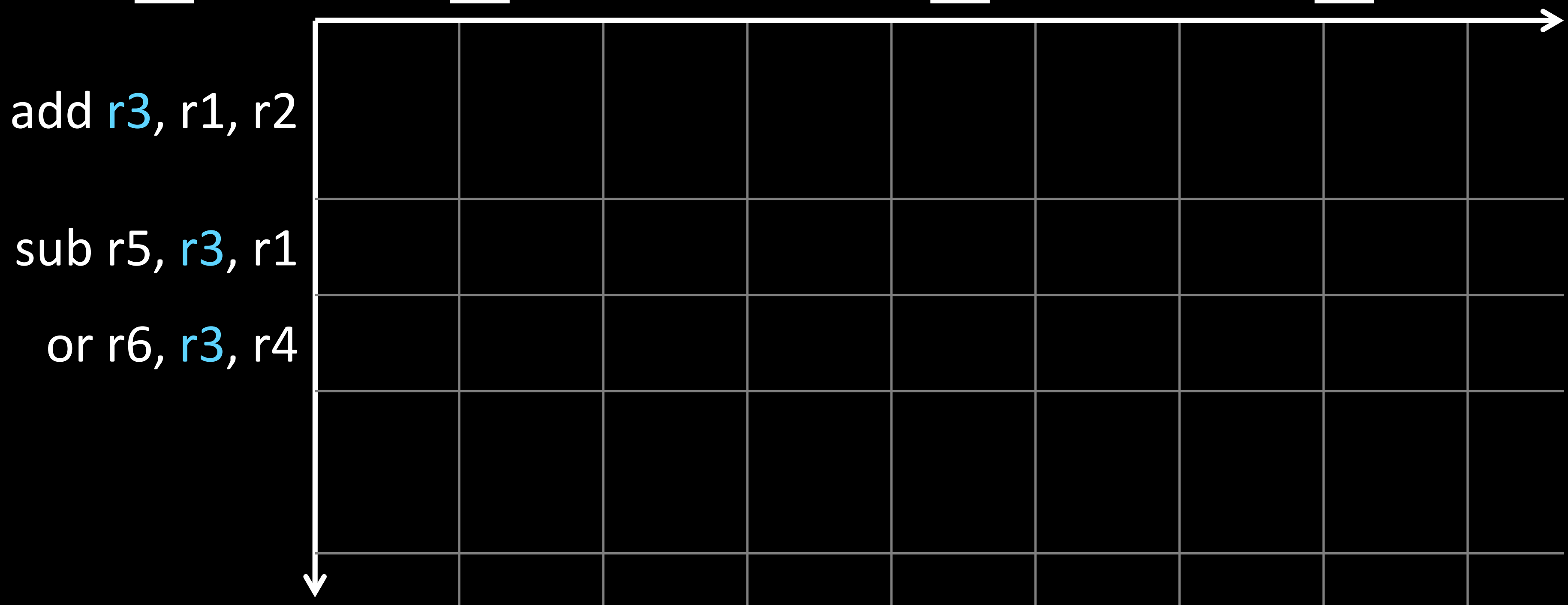
- RegisterFile Bypass

Forwarding Datapath 2





Register File Bypass

Register File Bypass

- Reading a value that is currently being written

Detect:

((Ra == MEM/WB.Rd) or (Rb == MEM/WB.Rd))
and (WB is writing a register)

Resolve:

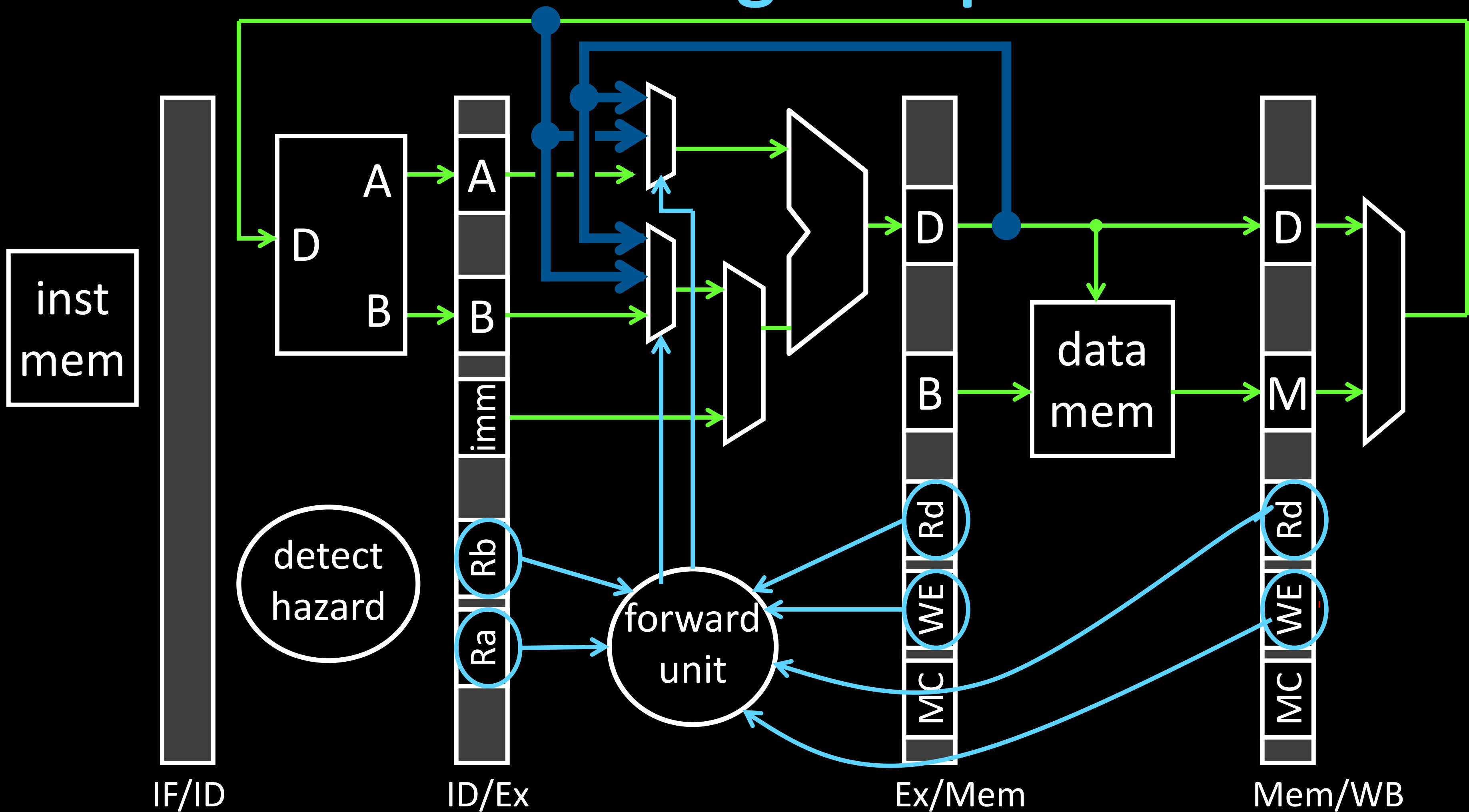
Add a bypass around register file (WB to ID)

Better: (Hack) just negate register file clock

– writes happen at end of first half of each clock cycle

– reads happen during second half of each clock cycle

Forwarding Datapath



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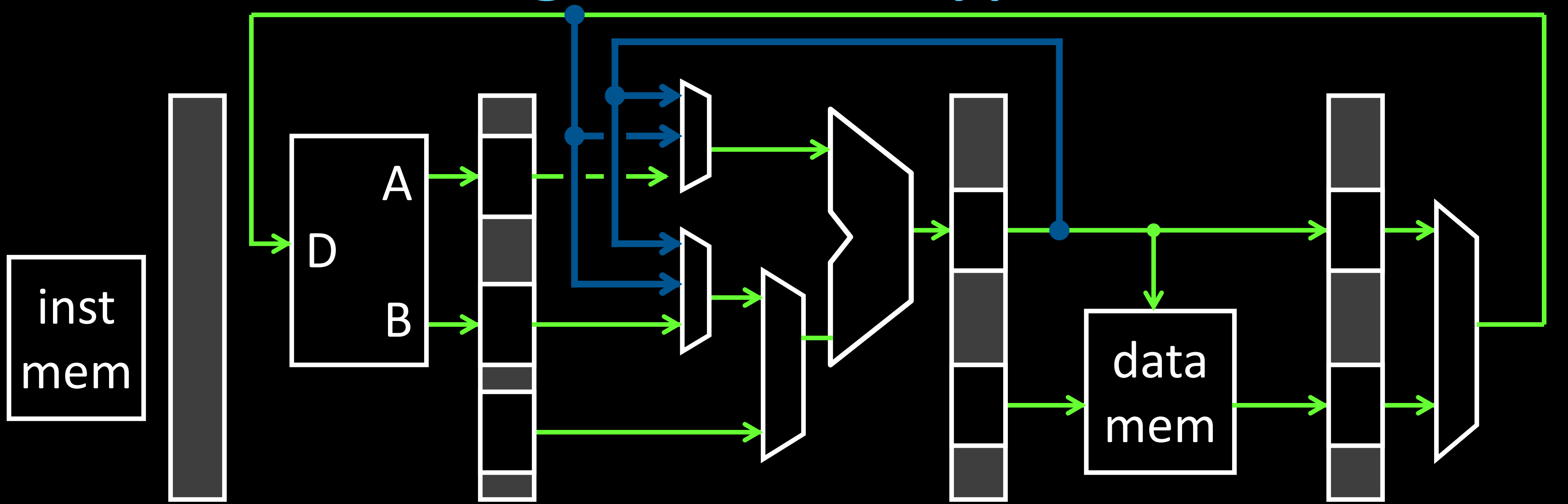
Register File Bypass



add r3, r1, r2

sub r5, r3, r1

or r6, r3, r4

add r6, r3, r8

Forwarding Example

time →

Clock cycle

1 2 3 4 5 6 7 8

r3 = 10

add r3, r1, r2

r3 = 20

sub r5, r3, r5

or r6, r3, r4

add r6, r3, r8

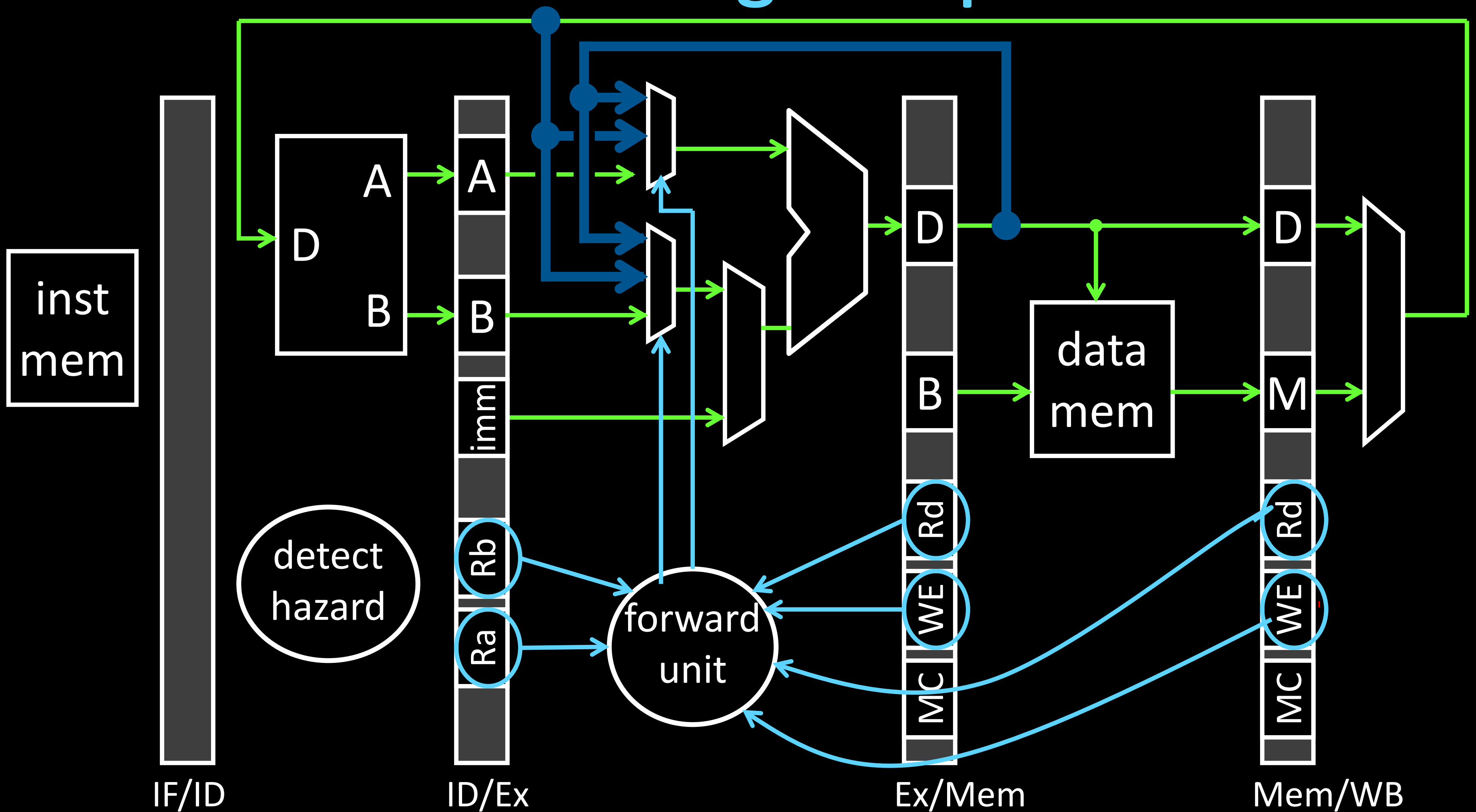
Forwarding Example 2

time →

Clock cycle

	1	2	3	4	5	6	7	8
add r3, r1, r2	IF	ID	Ex	M	W			
sub r5, r3, r4		IF	ID	Ex	M	W		
lw r6, 4(r3)								
or r5, r3, r5								
sw r6, 12(r3)								

Forwarding Datapath



Three types of forwarding/bypass

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- Forwarding from Mem/WB register to Ex stage (W → Ex)
- Register File Bypass

Takeaway

Data hazards occur when a operand (register) depends on the result of a previous instruction that may not be computed yet. A pipelined processor needs to detect data hazards.

Stalling, preventing a dependent instruction from advancing, is one way to resolve data hazards. Stalling introduces NOPs ("bubbles") into a pipeline. Introduce NOPs by (1) preventing the PC from updating, (2) preventing writes to IF/ID registers from changing, and (3) preventing writes to memory and register file. Bubbles (nops) in pipeline significantly decrease performance.

Forwarding bypasses some pipelined stages forwarding a result to a dependent instruction operand (register). Better performance than stalling.

Data Hazard Recap

Stall

- Pause current and all subsequent instructions

Forward/Bypass

- Try to steal correct value from elsewhere in pipeline
- Otherwise, fall back to stalling or require a delay slot

Tradeoffs?