

CS 3110

Lecture 22: Mechanized Logic

Prof. Clarkson

Fall 2014

Today's music: "Mr. Roboto" by Styx

The problem's plain to see: too much technology.

Machines to save our lives. Machines dehumanize.

Review

Current topic:

- How to reason about correctness of code
- Started with informal arguments
- Developed formal logic

Today:

- A proof assistant called Coq

Question #1

How much of PS5 have you finished?

- A. None
- B. About 25%
- C. About 50%
- D. About 75%
- E. I'm done!!!

Review: Proof rules of IPC, part 1

Rule name	Rule
\wedge intro	if $\mathbf{F} \vdash \mathbf{f1}$ and $\mathbf{F} \vdash \mathbf{f2}$ then $\mathbf{F} \vdash \mathbf{f1} \wedge \mathbf{f2}$
\wedge elim L	if $\mathbf{F} \vdash \mathbf{f1} \wedge \mathbf{f2}$ then $\mathbf{F} \vdash \mathbf{f1}$
\wedge elim R	if $\mathbf{F} \vdash \mathbf{f1} \wedge \mathbf{f2}$ then $\mathbf{F} \vdash \mathbf{f2}$
\Rightarrow elim	if $\mathbf{F} \vdash \mathbf{f}$ and $\mathbf{F} \vdash \mathbf{f} \Rightarrow \mathbf{g}$ then $\mathbf{F} \vdash \mathbf{g}$
\Rightarrow intro	if $\mathbf{F}, \mathbf{f} \vdash \mathbf{g}$ then $\mathbf{F} \vdash \mathbf{f} \Rightarrow \mathbf{g}$
assump	$\mathbf{f} \vdash \mathbf{f}$
weak	if $\mathbf{F} \vdash \mathbf{f}$ then $\mathbf{F}, \mathbf{g} \vdash \mathbf{f}$
set assump	$\mathbf{F}, \mathbf{f} \vdash \mathbf{f}$

Review: Proof rules of IPC, part 2

Rule name	Rule
\vee intro L	if $\mathbf{F} \vdash f1$ then $\mathbf{F} \vdash f1 \vee f2$
\vee intro R	if $\mathbf{F} \vdash f2$ then $\mathbf{F} \vdash f1 \vee f2$
\vee elim	if $\mathbf{F} \vdash f1 \vee f2$ and $\mathbf{F} \vdash f1 \Rightarrow g$ and $\mathbf{F} \vdash f2 \Rightarrow g$ then $\mathbf{F} \vdash g$
true intro	$\mathbf{F} \vdash \mathbf{true}$
false elim	if $\mathbf{F} \vdash \mathbf{false}$ then $\mathbf{F} \vdash f$
\sim intro	if $\mathbf{F} \vdash f \Rightarrow \mathbf{false}$ then $\mathbf{F} \vdash \sim f$
\sim elim	if $\mathbf{F} \vdash \sim f$ then $\mathbf{F} \vdash f \Rightarrow \mathbf{false}$

Review: Proof rules of IQC

Rule name	Rule
---	All rules of IPC
forall intro	if $\mathbf{F} \vdash f(\mathbf{x})$ and \mathbf{x} not in $FV(\mathbf{F})$ then $\mathbf{F} \vdash \text{forall } \mathbf{x}, f(\mathbf{x})$
forall elim	if $\mathbf{F} \vdash \text{forall } \mathbf{x}, f(\mathbf{x})$ then $\mathbf{F} \vdash f(\mathbf{t})$
exists intro	if $\mathbf{F} \vdash f(\mathbf{t})$ then $\mathbf{F} \vdash \text{exists } \mathbf{x}, f(\mathbf{x})$
exists elim	if $\mathbf{F} \vdash \text{exists } \mathbf{x}, f(\mathbf{x})$ and $\mathbf{F} \vdash f(\mathbf{x}) \Rightarrow \mathbf{g}$ and \mathbf{x} not in $FV(\mathbf{F}, \mathbf{g})$ then $\mathbf{F} \vdash \mathbf{g}$

Theories

- IQC reaches its full power when augmented with *theories*
- Collections of
 - names of relations and functions, and
 - new proof rules for those

Theory of equality

- Relation: **equals** (t_1, t_2)
 - normally written $t_1=t_2$
- Proof rules:
 - reflexivity: $t=t$
 - symmetry: if $t_1=t_2$ then $t_2=t_1$
 - transitivity: if $t_1=t_2$ and $t_2=t_3$ then $t_1=t_3$
 - eq-fn: if $t_1=u_1$ and...and $t_n=u_n$ then
 $fn(t_1, \dots, t_n) = fn(u_1, \dots, u_n)$
 - eq-rel: if $t_1=u_1$ and...and $t_n=u_n$ then
 $R(t_1, \dots, t_n) = R(u_1, \dots, u_n)$

Theory of rings

- *Ring*: mathematical structure that abstracts addition and multiplication
 - see Math 4320
- Relies on theory of equality
- Functions:
 - **plus**(**t1**, **t2**) and **mult**(**t1**, **t2**) and **neg**(**t**)
 - written **t1+t2** and **t1*t2** and **-t**
 - **zero** and **one**
 - written **0** and **1**

Theory of rings

- Proof rules (all are axioms):
 - forall $a\ b\ c$, $(a+b)+c = a+(b+c)$
 - forall $a\ b$, $a+b = b+a$
 - forall a , $0+a = a$
 - forall a , $a + (-a) = 0$
 - forall $a\ b\ c$, $a*(b+c) = (a*b)+(a*c)$
 - forall $a\ b\ c$, $(b+c)*a = (b*a)+(c*a)$
 - forall $a\ b\ c$, $(a*b)*c = a*(b*c)$
 - forall $a\ b$, $a*b = b*a$
 - forall a , $1*a = a$
- Syntactic sugar:
forall $a\ b$, f
means **forall a , (forall b , f)**

Prelim 2

- One week from today
- Covers everything from Oct 2 through Nov 12 (inclusive)
 - People with Thursday recitations, note that today's recitation is included
- Sample prelim posted on Piazza
- Review session in recitation day before prelim
- Cancel lecture on day of prelim
- You can take prelim at your choice of 5:30-7:00 pm or 7:30-9:00 pm; no need to reserve in advance
- Three rooms, will be assigned by netid next week
- Closed book
 - But you may have one page of notes
 - 8.5x11" two-sided 😊

Why formal logic?

- Humans make mistakes in writing proofs
- Humans make mistakes in checking proofs
- Formal logic:
 - Reduces proof to symbolic manipulation
 - Maybe a machine could check that manipulation
- Analogy:
 - Compiler type checks program
 - Proof checker uses proof rules we've given to check proof

Mechanized proof

- Automated theorem provers
 - You give tool a theorem
 - Tools finds a proof or a counterexample
 - Or runs out of time
 - e.g., Z3, developed at Microsoft
 - Ships with the Windows 7 device driver developer's kit
- Proof assistants
 - You give tool a theorem
 - You and tool cooperatively find proof
 - Human guides the construction
 - Machine does the low-level details
 - e.g., Coq, Isabelle/HOL, NuPRL
 - NuPRL: Prof. Constable (Cornell)
 - Coq: used to verify compiler, OS kernel, etc.

Coq

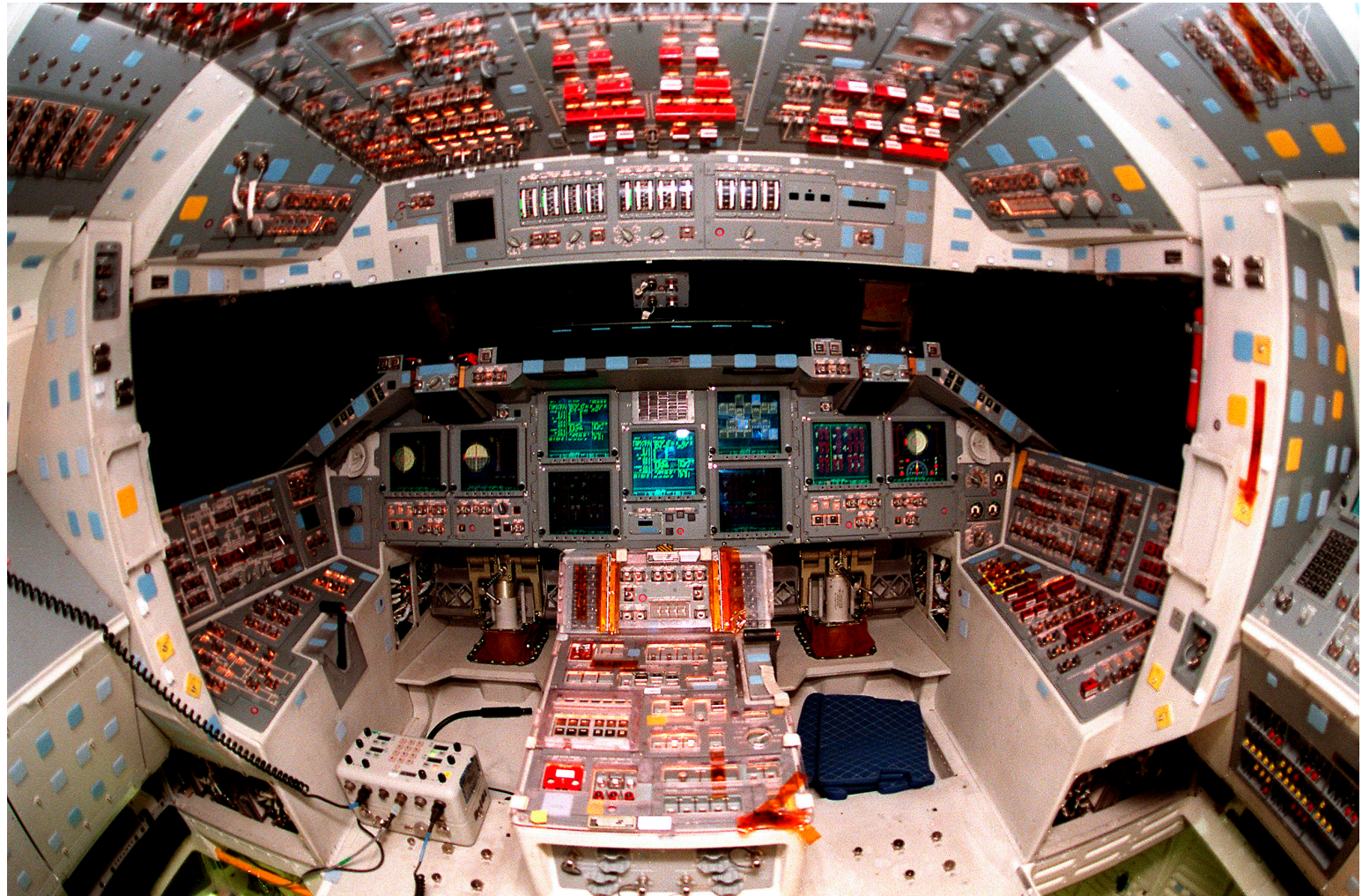


- **1984:** Coquand and Huet first begin implementing a new theorem prover Coq based on *calculus of inductive constructions*
- **1992:** Coq ported to Caml
- **2012:** Coq version 8.4
 - Implemented in OCaml
 - Can produce verified OCaml code



Thierry Coquand
1961 –

Coq's full system



Subset of Coq we'll use



Coq3110.v

- We went through the file up through and including implication and forall.

Please hold still for 1 more minute

WRAP-UP FOR TODAY

Upcoming events

- **PS5 due tonight**
- Prelim 2 in one week

This is mechanized.

THIS IS 3110