## Finite Automata

## CS 2800: Discrete Structures, Spring 2015

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## A simplified model

## Machine

## A simplified model


an alphabet $\Sigma$, e.g. $\{0,1\}$

## A general-purpose computer



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## String <br>  <br> String

Church-Turing Thesis: Any "effective/mechanical/real-world" calculation can be carried out on a Turing machine


Alan Turing, 1912-1954

## A simple "computer"



## An example



## An example


(Binary input)

## An example



Input: 01001

## An example



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## An example



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## An example



Input: 01001

## An example



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## An example



Input: 01001

## An example



Input: 01001

## An example



Input: 01001
Output: Yes!

## An example



In general, on what binary strings does this DFA return Yes?

## An example



Ans: All strings with an even number of 1's

## Deterministic Finite Automaton

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- $\delta$ is a transition function $\delta: Q \times \Sigma \rightarrow Q$
- $q_{0} \in Q$ is the start/initial state
- $F \subseteq Q$ is the set of final/accepting states


What does this DFA accept?


## What does this DFA accept?



Answer: No strings

## What does this DFA accept?



## What does this DFA accept?



Answer: All strings

## What does this DFA accept?



## What does this DFA accept?



Answer: Strings of length 1

## What does this DFA accept?



Answer: Strings of length 1

What does this DFA accept?


## What does this DFA accept?



Answer: Strings containing only 1's

What does this DFA accept?


## What does this DFA accept?



Answer: Strings containing no two consecutive 1's

## Language

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## What language does this DFA accept?



## What language does this DFA accept?



Answer: Only the string 1

What language does this DFA accept?


What language does this DFA accept?


Answer: Only the string 11

## DFA's find it difficult to count

- A DFA that recognizes the language $\left\{1^{c}\right\}$ (the single string of $c$ 1's) must have at least $c$ states


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- A DFA that recognizes the language $\left\{1^{c}\right\}$ (the single string of $c$ 1's) must have at least $c$ states
- The parent alphabet is irrelevant (but must of course contain 1)
(Proof discussion to be
completed next class)

