#### Finite Automata

#### CS 2800: Discrete Structures, Spring 2015

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# A simplified model



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#### A general-purpose computer





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Church-Turing Thesis: Any "effective/mechanical/real-world" calculation can be carried out on a Turing machine



Alan Turing, 1912 – 1954

# A simple "computer"







(Binary input)



















In general, on what binary strings does this DFA return Yes?



#### Ans: All strings with an even number of 1's

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  - $F \subseteq Q$  is the set of final/accepting states







Answer: No strings





#### Answer: All strings





#### Answer: Strings of length 1

#### What does this DFA accept? Hell (any (any symbol) symbol) $q_0$ $q_1$ $q_{2}$ (any symbol)

#### Answer: Strings of length 1





## Answer: Strings containing only 1's





# Answer: Strings containing no two consecutive 1's

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#### Answer: Only the string 1





Answer: Only the string 11

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- A DFA that recognizes the language  $\{1^c\}$  (the single string of c 1's) must have at least c states
  - The parent alphabet is irrelevant (but must of course contain 1)

(Proof discussion to be completed next class)