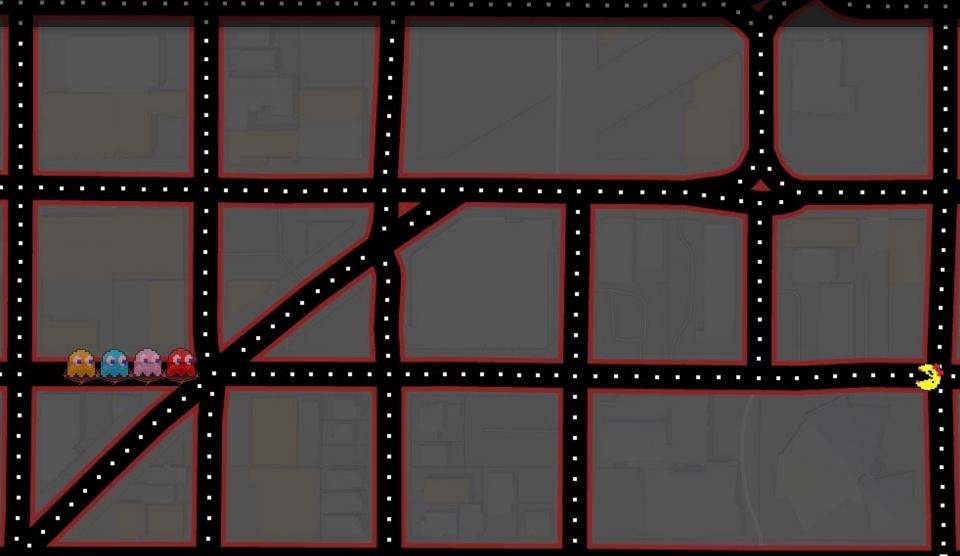
Graphs - II CS 2110, Spring 2017



Announcements

- 1. Lunch with instructors: Still many slots open!
 - See pinned Piazza note @275

- 2. Hidden Figures: Thursday at 8:30 and 11:00!
 - If you haven't signed up on the CMS schedule do so now!
 - so far: 8:30 114/330 11:00 147/330

Announcements

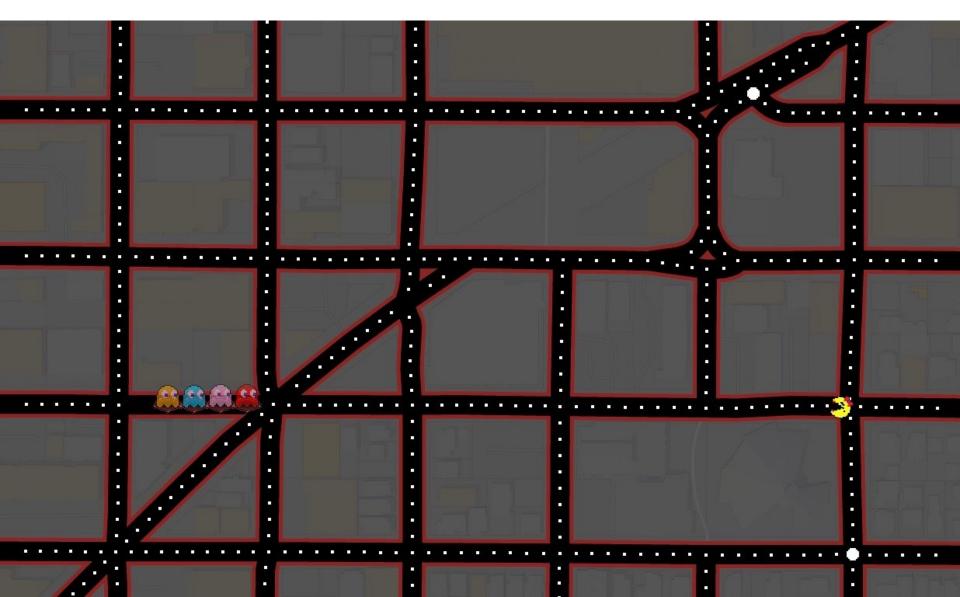
3. For next Tuesday's lecture, you **MUST** watch the tutorial on the shortest path algorithm beforehand:

<u>http://www.cs.cornell.edu/courses/cs2110/2017sp/online/shortest</u> <u>Path/shortestPath.html</u>

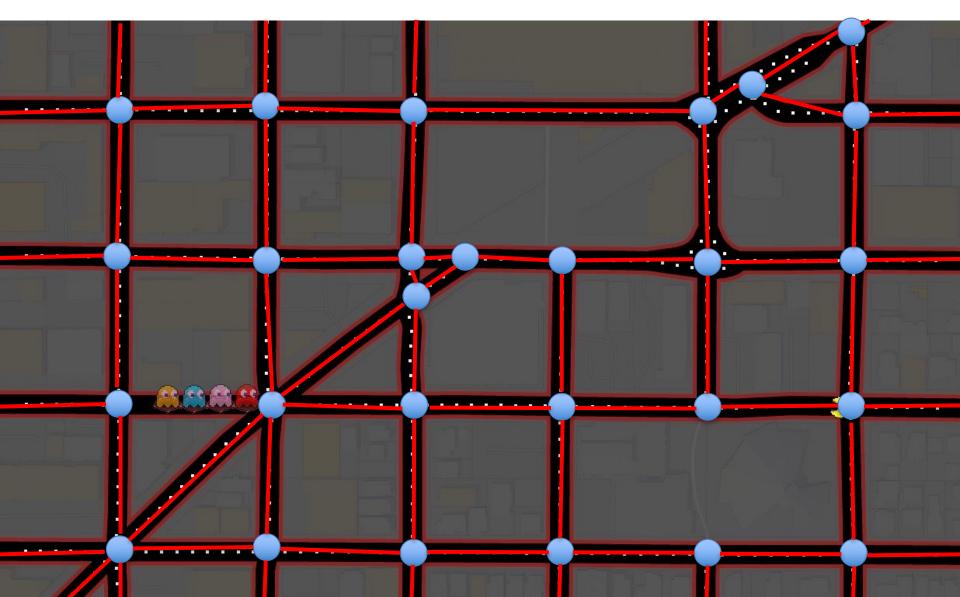
Tuesday's lecture will assume that you understand it.
 Watch the tutorial once or twice and execute the algorithm on a small graph.

4. Information about **Prelim 2** is now on the Exams page of the course website and P2Conflict will be available on the CMS today (Tuesday).

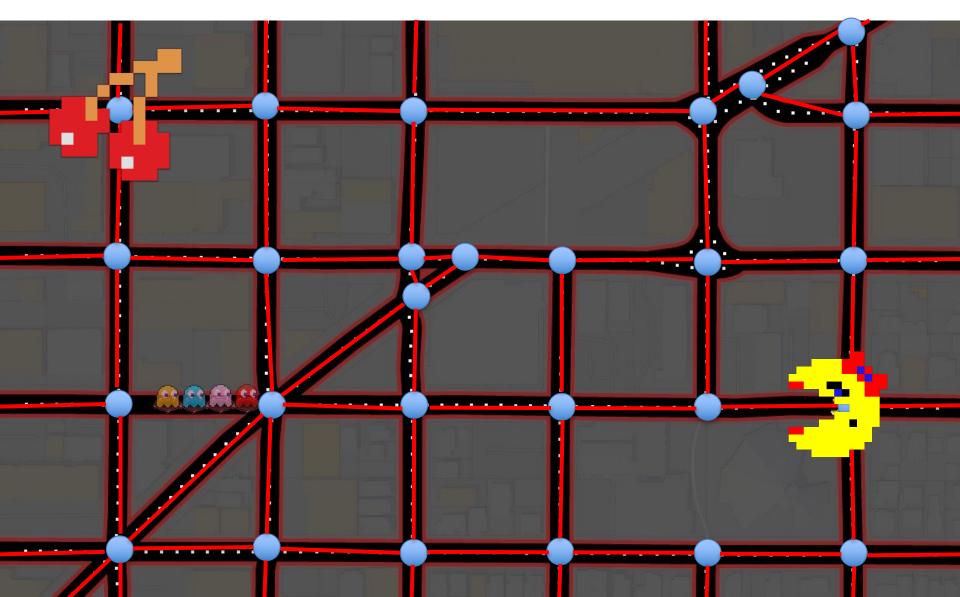
Look, a graph!



Look, a graph!



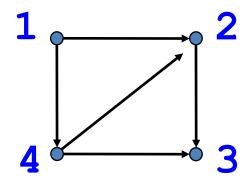
Look, a graph!



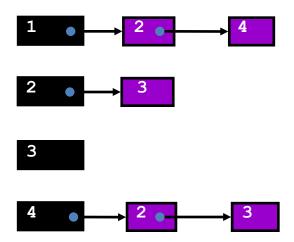
Graph Algorithms

- Search
 - Depth-first search
 - Breadth-first search
- Shortest paths
 - Dijkstra's algorithm
- Minimum spanning trees
 - Prim's algorithm
 - Kruskal's algorithm

Representations of Graphs



Adjacency List



Adjacency Matrix

	1	2	3	4
1	0	1	0	1
2	0	0	1	0
3	0	0	0	0
4	0	1	1	0

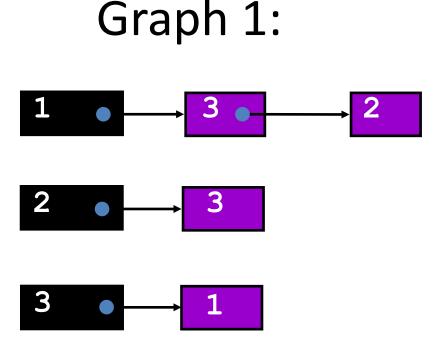
Adjacency Matrix or Adjacency List?

- Definitions:
 - n = number of vertices
 - e = number of edges
 - d(u) = degree of u = number of edges leaving u
- Adjacency Matrix
 - Uses space $O(n^2)$
 - Can iterate over all edges in time $O(n^2)$
 - Can answer "Is there an edge from u to v?" in O(1) time
 - Better for dense graphs (lots of edges)
- Adjacency List
 - Uses space O(e + n)
 - Can iterate over all edges in time O(e + n)
 - Can answer "Is there an edge from u to v?" in O(d(u)) time
 - Better for sparse graphs (fewer edges)

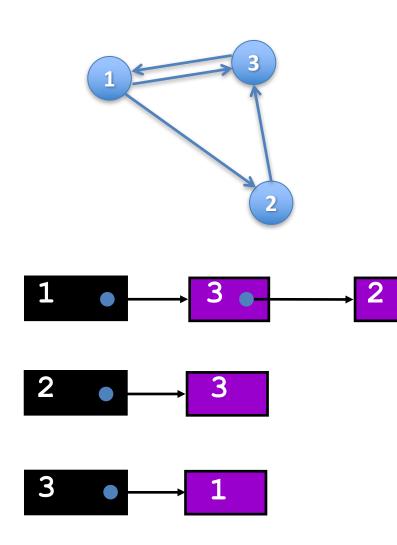
Breaking DAG

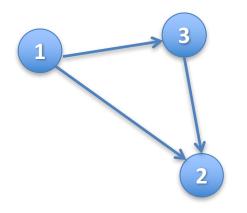
Which of the following two graphs are DAGs? Directed Acyclic Graph

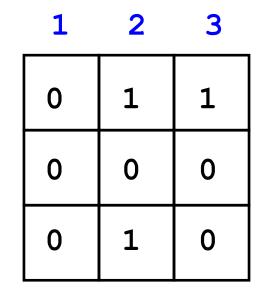
Graph 2:



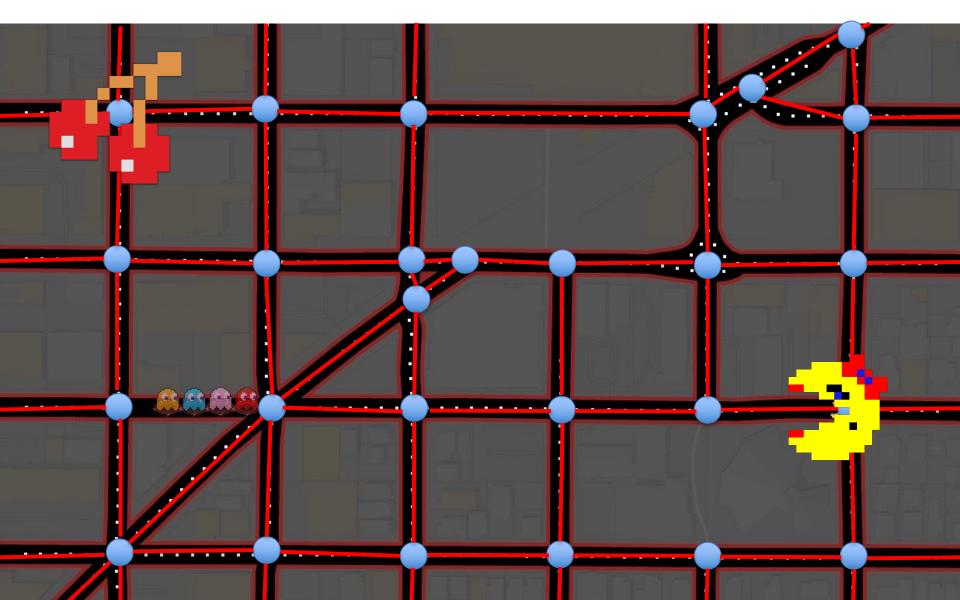
Breaking DAG





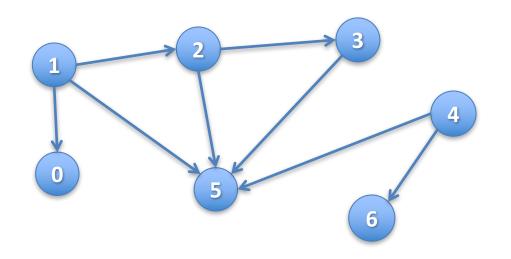


Back to Important Things:

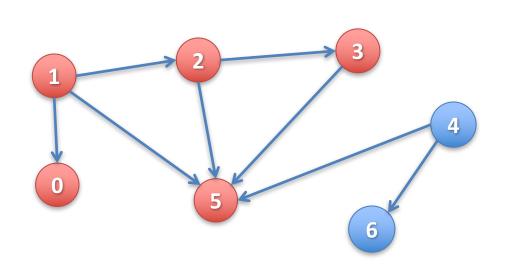


• Given a graph and one of its nodes *u*

(say node 1 below)



- Given a graph and one of its nodes u (say node 1 below)
- We want to "visit" each node reachable from u (nodes 1, 0, 2, 3, 5)



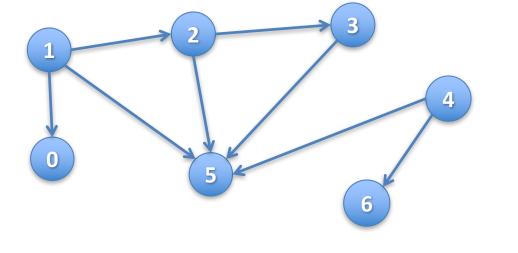
There are many paths to some nodes.

How do we visit all nodes efficiently, without doing extra work?

boolean[] visited;

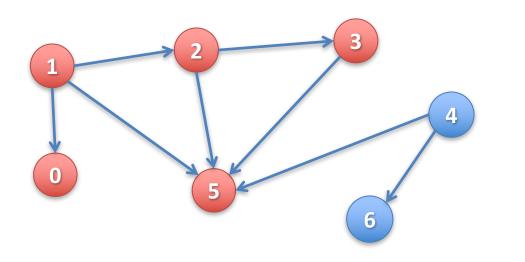
- Node u is visited means: visited[u] is true
- To visit u means to: set visited[u] to true
- v is REACHABLE* from u if there is a path (u, ..., v)
 *in which all nodes of the path are unvisited.

Suppose all nodes are unvisited.



boolean[] visited;

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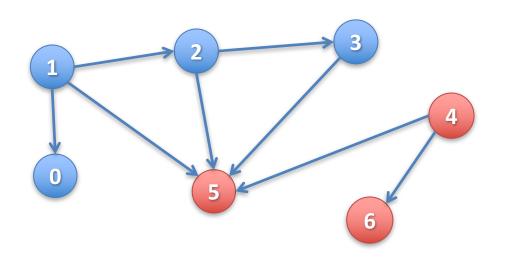


Suppose all nodes are unvisited.

Nodes REACHABLE* from node **1**: {1, 0, 2, 3, 5}

boolean[] visited;

- Node u is visited means: visited[u] is true
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Suppose all nodes are unvisited.

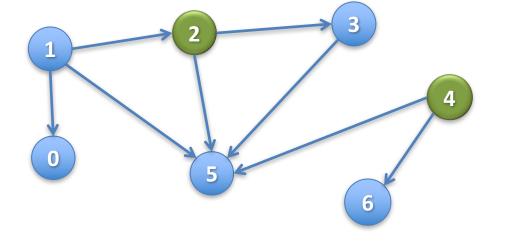
Nodes REACHABLE* from node **1**: {1, 0, 2, 3, 5}

Nodes REACHABLE* from **4**: {4, 5, 6}

boolean[] visited;

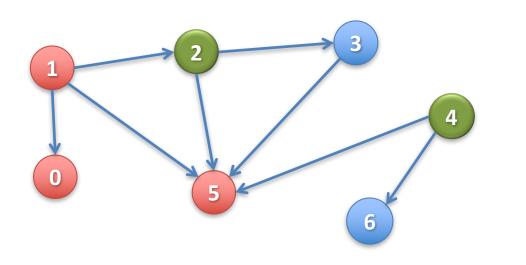
- Node u is visited means: visited[u] is true
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- v is REACHABLE* from u if there is a path (u, ..., v)
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Green: visited **Blue**: unvisited



boolean[] visited;

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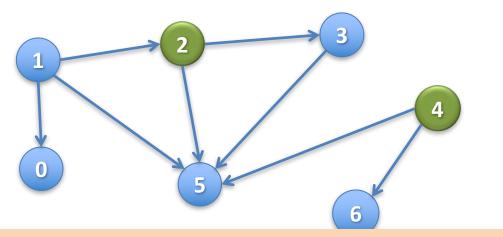


Green: visited **Blue**: unvisited

Nodes REACHABLE* from node **1**: {1, 0, 5}

boolean[] visited;

- Node u is visited means: visited[u] is true
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- v is REACHABLE* from u if there is a path (u, ..., v)
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Green: visited **Blue**: unvisited

Nodes REACHABLE * from node **1**: {1, 0, 5}

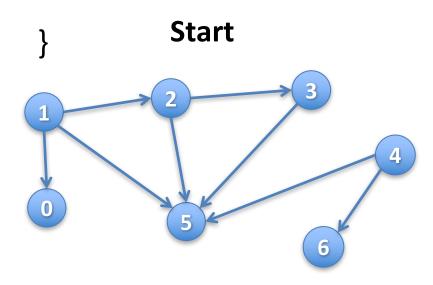
Nodes REACHABLE * from **4**: none

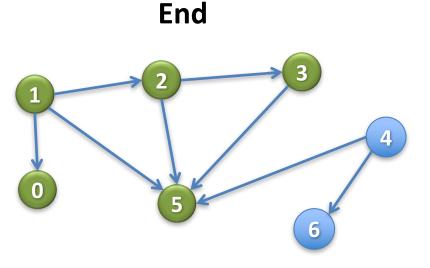
Not even **4** itself, because it's already been visited!

/** Visit all nodes that are REACHABLE*
from u. Precondition: u is unvisited. */
public static void dfs(int u) {

Let u be **1**

The nodes REACHABLE* from 1 are 1, 0, 2, 3, 5

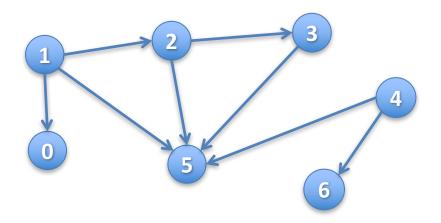




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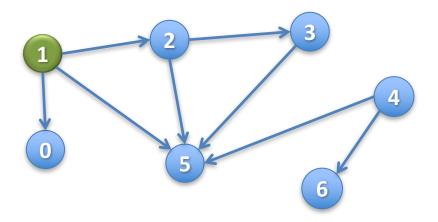


/** Visit all nodes that are REACHABLE*
from u. Precondition: u is unvisited. */
public static void dfs(int u) {

visited[u] = true;

Let u be **1**

The nodes REACHABLE* from 1 are 1, 0, 2, 3, 5

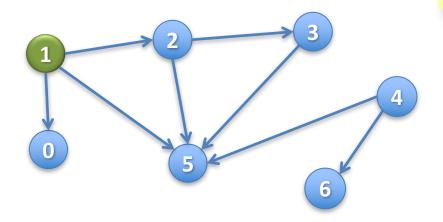


/** Visit all nodes that are REACHABLE*
from u. Precondition: u is unvisited. */
public static void dfs(int u) {

visited[u] = true;

Let u be 1 (visited)

The nodes to be visited are 0, 2, 3, 5



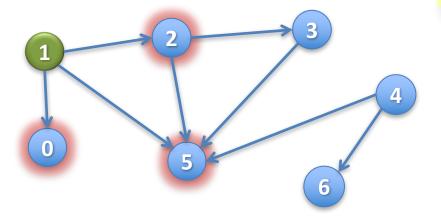
/** Visit all nodes that are REACHABLE*
from u. Precondition: u is unvisited. */
public static void dfs(int u) {
 visited[u] = true;
 for all edges (u, v) leaving u:
 if v is unvisited then dfs(v);

}

Let u be 1 (visited)

The nodes to be visited are 0, 2, 3, 5

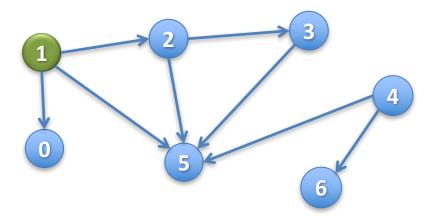
Have to do DFS on all unvisited neighbors of u!



/** Visit all nodes that are REACHABLE*
from u. Precondition: u is unvisited. */
public static void dfs(int u) {
 visited[u] = true;
 for all edges (u, v) leaving u:
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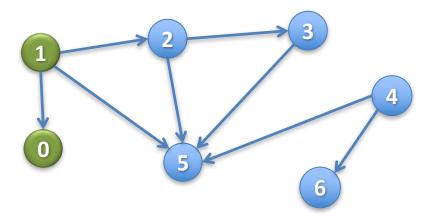
Suppose the **for** loop visits neighbors in numerical order. Then **dfs(1)** visits the nodes in this order: **1** ...



/** Visit all nodes that are REACHABLE*
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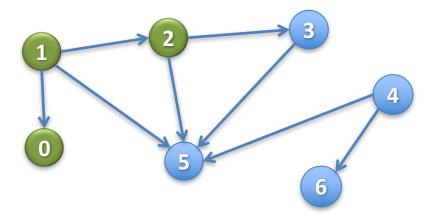
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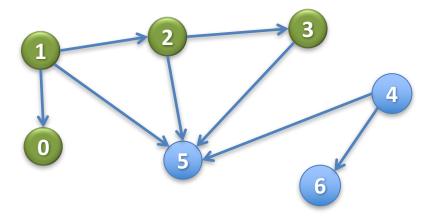
Suppose the **for** loop visits neighbors in numerical order. Then **dfs(1)** visits the nodes in this order: **1**, **0**, **2** ...



/** Visit all nodes that are REACHABLE*
from u. Precondition: u is unvisited. */
public static void dfs(int u) {
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 for all edges (u, v) leaving u:
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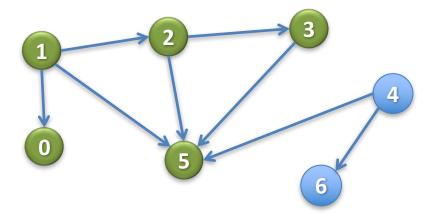
Suppose the **for** loop visits neighbors in numerical order. Then **dfs(1)** visits the nodes in this order: 1, 0, 2, 3 ...



/** Visit all nodes that are REACHABLE*
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public static void dfs(int u) {
 visited[u] = true;
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Suppose the **for** loop visits neighbors in numerical order. Then **dfs(1)** visits the nodes in this order: 1, 0, 2, 3, 5



```
/** Visit all nodes that are REACHABLE* from u. Precondition: u is unvisited. */
```

public static void dfs(int u) {

```
visited[u] = true;
```

}

```
for all edges (u, v) leaving u:
```

```
if v is unvisited then dfs(v);
```

```
Suppose n nodes are REACHABLE* along e edges (in total). What is
```

- Worst-case runtime? O(n+e)
- Worst-case space? O(n)

/** Visit all nodes that are REACHABLE*
from u. Precondition: u is unvisited. */
public static void dfs(int u) {
 visited[u] = true;
 for all edges (u, v) leaving u:
 if v is unvisited then dfs(v);
}

Example: Use different way (other than array visited) to know whether a node has been visited

Example: We really haven't said what data structures are used to implement the graph

That's all there is to basic DFS. You may have to change it to fit a particular situation.

If you don't have this spec and you do something different, it's probably wrong.

Depth-First Search in OO fashion

public class Node {
 boolean visited;
 List<Node> neighbors;

Each node of the graph is an object of type Node

/** Visit all nodes that are REACHABLE*
 * from u. Precondition: u is unvisited */
public void dfs() {
 visited= true;
 for (Node n: neighbors) {
 if (!n.visited) n.dfs();
 }
}
No need for a
parameter. The
object is the node.

Depth-First Search written iteratively

/** Visit all nodes REACHABLE* from u. Pre: u is unvisited. */
public static void dfs(int u) {

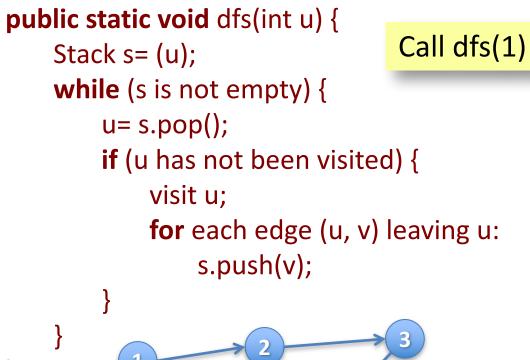
Stack s= (u); // Not Java!
// inv: all nodes that have to be visited are
// REACHABLE* from some node in s
while (s is not empty) {
 u= s.pop(); // Remove top stack node, put in u
 if (u has not been visited) {
 visit u;
 for each edge (u, v) leaving u:

s.push(v);

Depth-First Search written iteratively

6

/** Visit all nodes REACHABLE* from u. Pre: u is unvisited. */





Stack s

Depth-First Search written iteratively

/** Visit all nodes REACHABLE* from u. Pre: u is unvisited. */ public static void dfs(int u) { Call dfs(1) Iteration 0 Stack s= (u); while (s is not empty) { u= s.pop(); if (u has not been visited) { visit u; for each edge (u, v) leaving u: s.push(v); 1 Stack s 6

/** Visit all nodes REACHABLE* from u. Pre: u is unvisited. */ public static void dfs(int u) { Call dfs(1) Iteration 0 Stack s= (u); while (s is not empty) { u= s.pop(); if (u has not been visited) { visit u; for each edge (u, v) leaving u: s.push(v); Stack s 6

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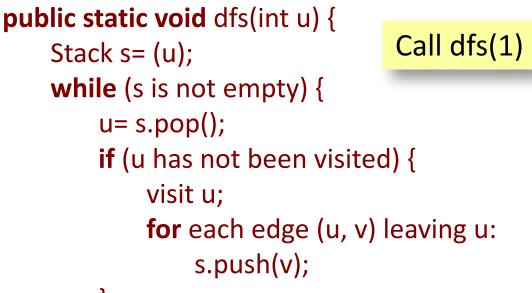
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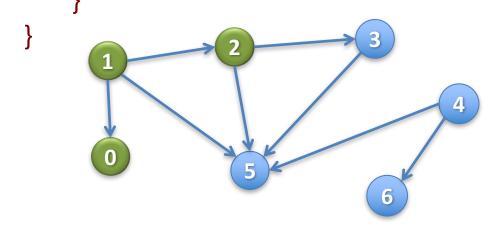
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/** Visit all nodes REACHABLE* from u. Pre: u is unvisited. . */



Iteration 2

Yes, 5 is put on the stack twice, once for each edge to it. It will be visited only once.



3 5 5

Stack s

/** Visit all nodes REACHABLE* from u. Pre: u is unvisited. . */ public static void dfs(int u) {

```
Stack s= (u);
while (s is not empty) {
    u= s.pop();
    if (u has not been visited) {
        visit u;
        for each edge (u, v) leaving u:
            s.push(v);
    }
```

That's DFS!

/** Visit all nodes REACHABLE* from u. Pre: u is unvisited. */ public static void dfs(int u) { Stack s= (u); // Not Java! // inv: all nodes that have to be visited are **REACHABLE*** from some node in s while (s is not empty) { u= s.pop(); // Remove top stack node, put in u if (u has not been visited) { visit u; for each edge (u, v) leaving u: s.push(v); }

Want to see a magic trick?

Depth-First Search

/** Visit all nodes REACHABLE* from u. Pre: u is unvisited. */
public static void dfs(int u) {
 Stack s= (u); // Not Java!
 // inv: all nodes that have to be visited are
 // REACHABLE* from some node in s
 while (s is not empty) {
 u= s.pop(); // Remove top stack node
 if (u has not been visited) {
 }
}

visit u; for each edge (u, v) leaving u: s.push(v);

/** Visit all nodes REACHABLE* from u. Pre: u is unvisited. */
public static void bfs(int u) {

Queue q= (u); // Not Java!

// inv: all nodes that have to be visited are

// REACHABLE* from some node in s

while (q is not empty) {

u= q.popFirst(); // Remove first node in queue, put in u
if (u has not been visited) {

visit u;

}

for each edge (u, v) leaving u:

q.append(v); // Add to end of queue

/** Visit all nodes REACHABLE* from u. Pre: u is unvisited. */
public static void bfs(int u) {

Queue q= (u); // Not Java!

// inv: all nodes that have to be visited are

// REACHABLE* from some node in s

while (q is not empty) {

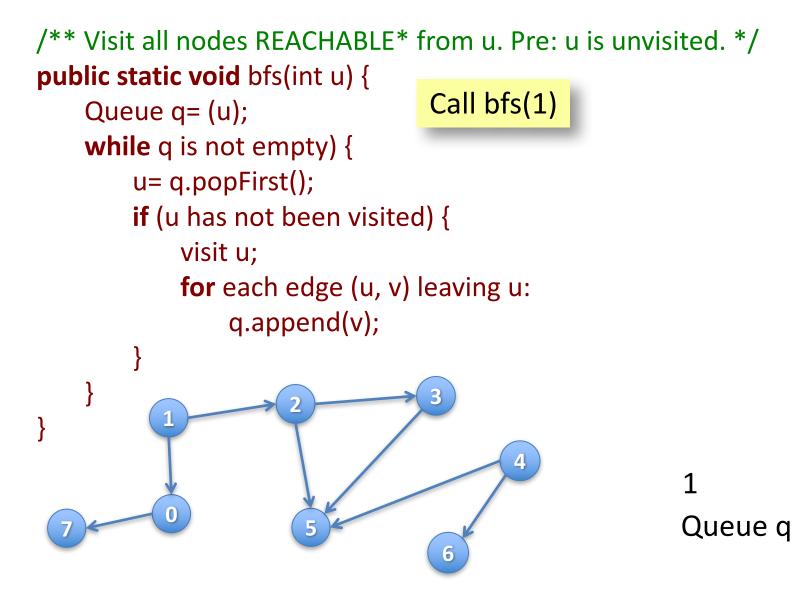
u= q.popFirst(); // Remove first node in queue, put in u
if (u has not been visited) {

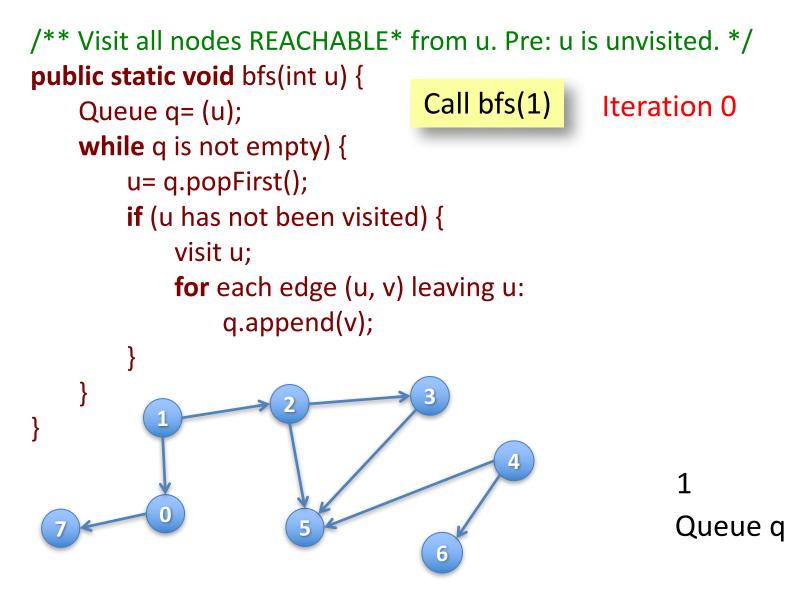
visit u;

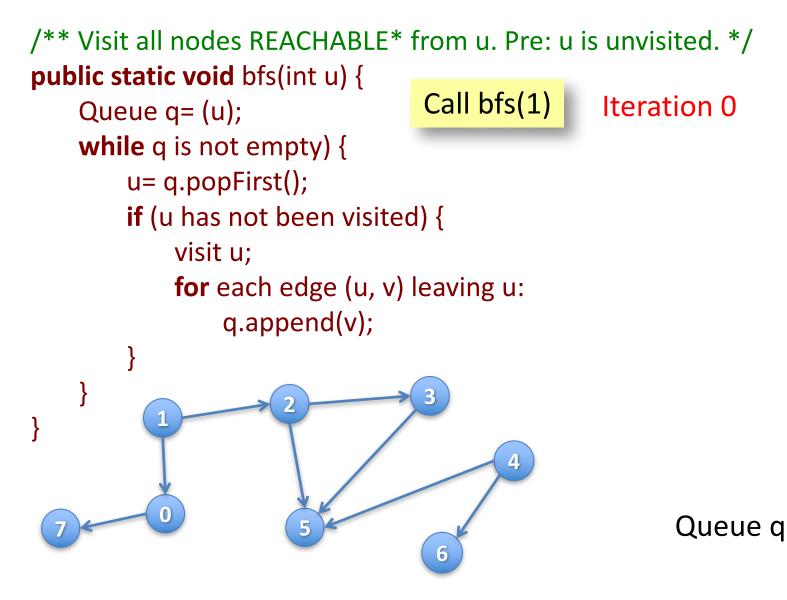
}

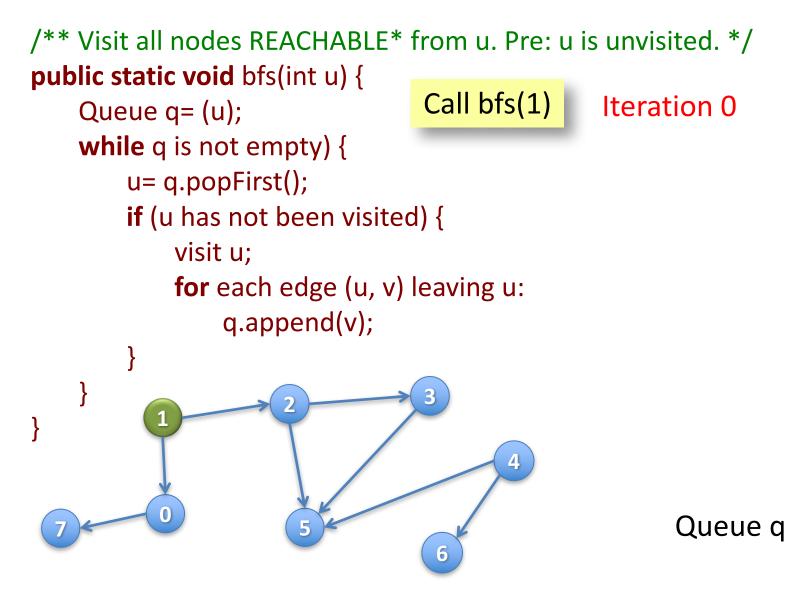
for each edge (u, v) leaving u:

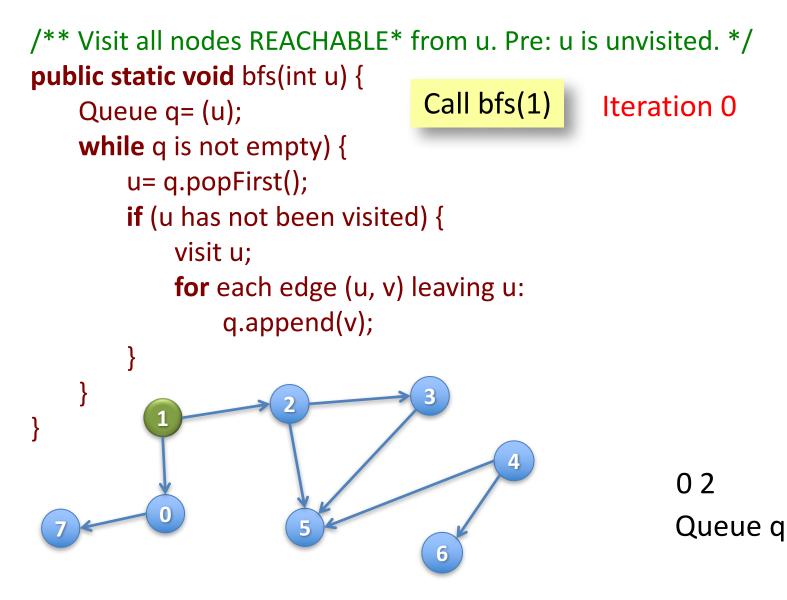
q.append(v); // Add to end of queue

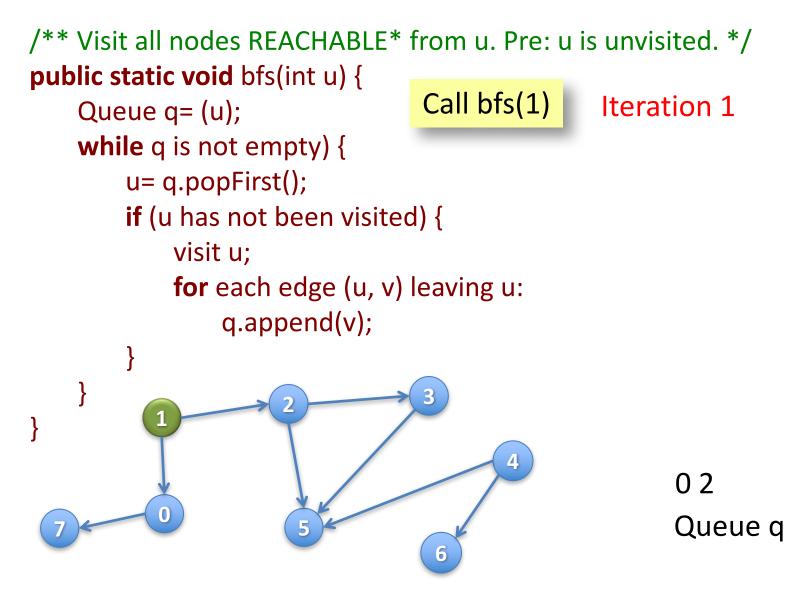


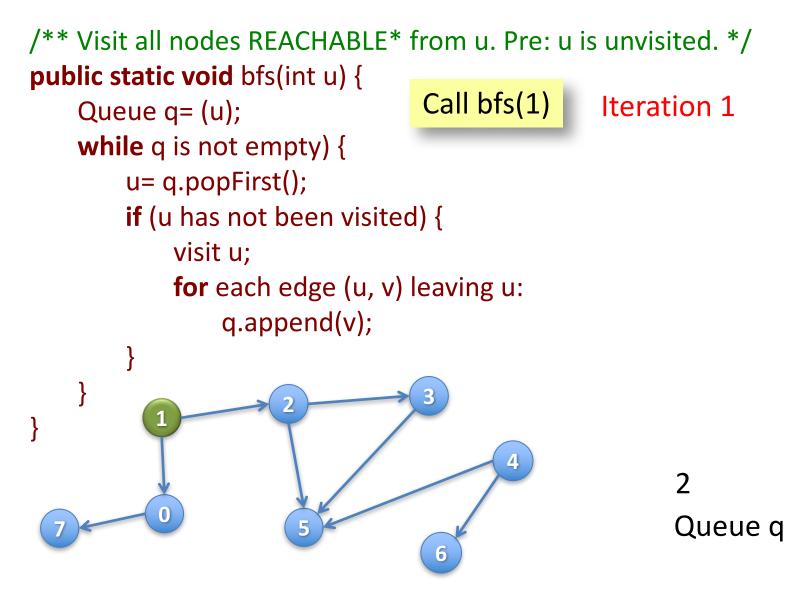


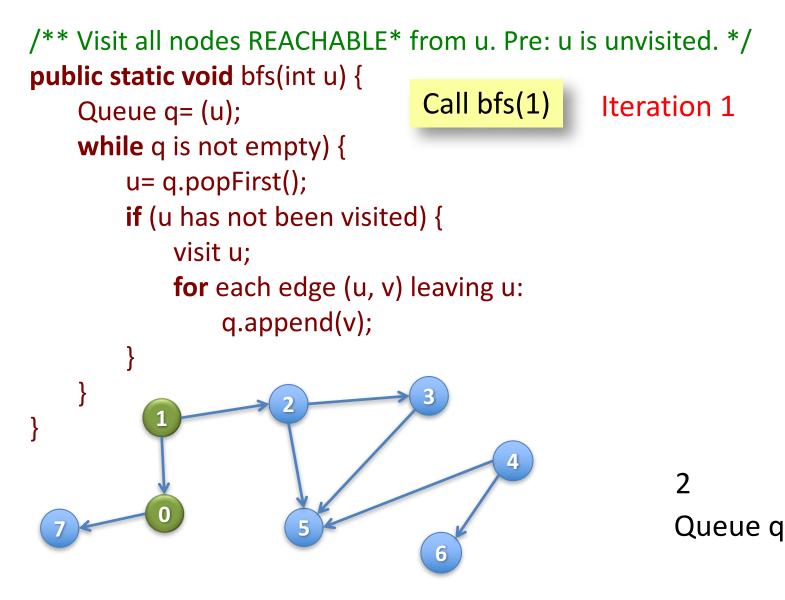


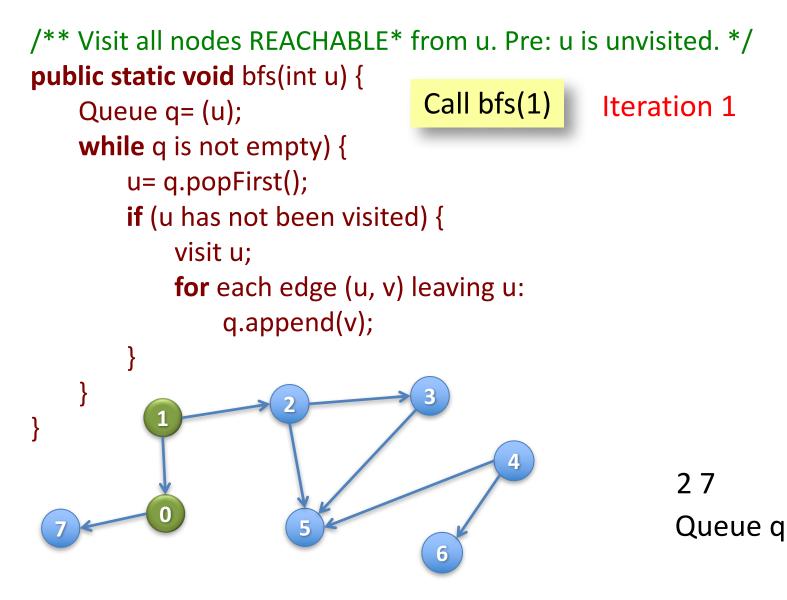


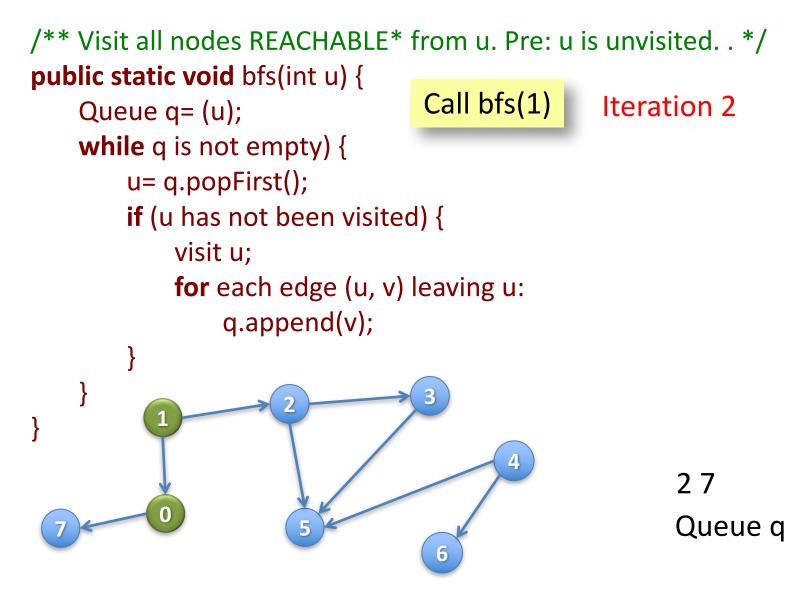


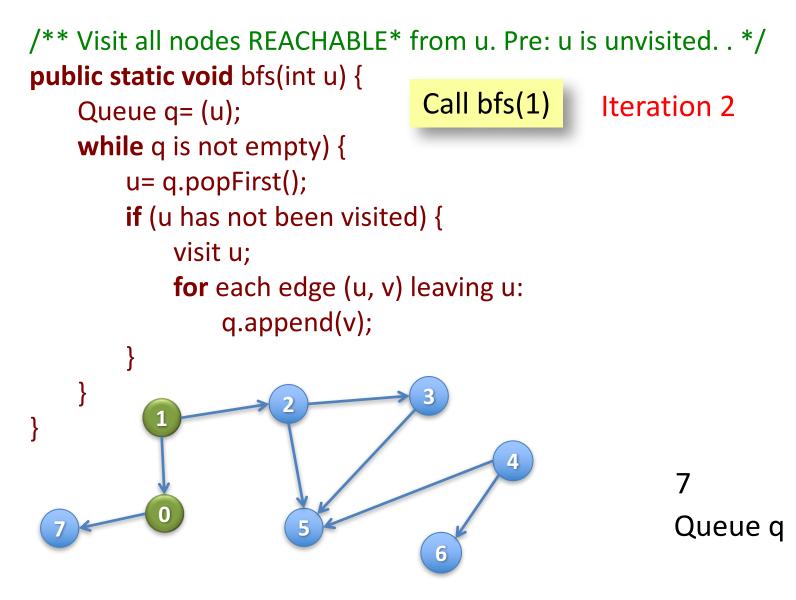


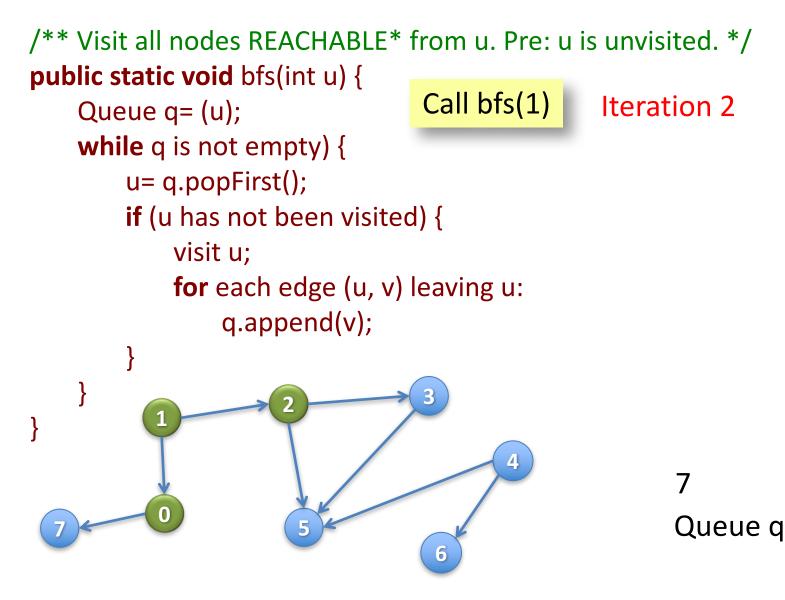


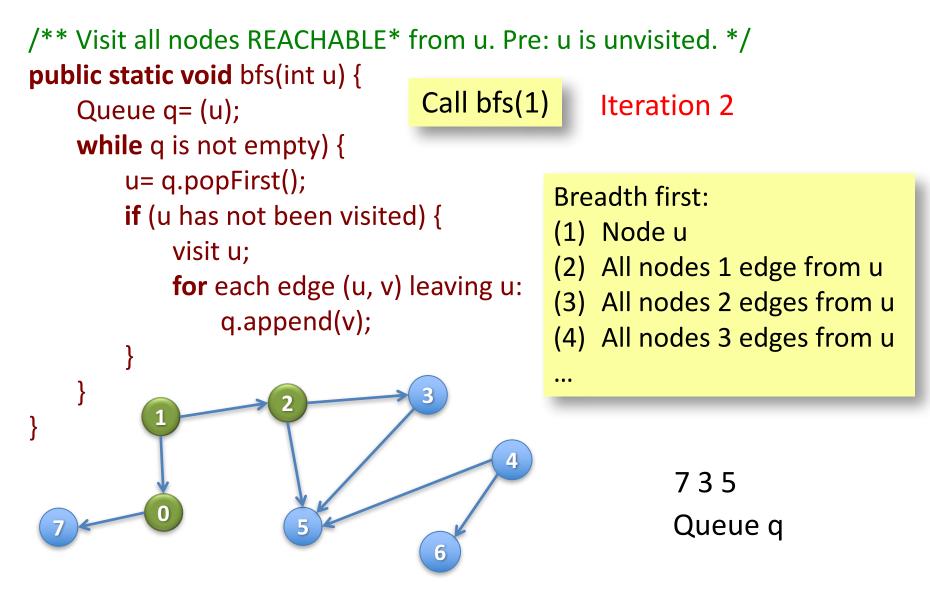












Some food for thought:

- BFS(root) on a tree corresponds to which tree traversal?
- Write out the order nodes are visited in this undirected graph, when calling:
 - BFS(5)
 - DFS(5)
 - DFS(0)

(if there are ties, visit the lower # first)

