CONCURRENCY 2

CS 2110 – Spring 2016

Consistency



What is printed?

0, 1, and 2 can be printed!



Thread 1 on Core 1

Thread 2 on Core 2

Write 2 to x in local cache

Write 3 to y in local cache

3 gets pushed to y in memory

Read 3 from y in memory Read 1 from x in memory Write 1 to a Print 1

2 gets pushed to x in memory

Not sequentially consistent!

Harsh Reality

Sequential Consistency

- There is an interleaving of the parallel operations that explains the observations and events
- Currently unknown how to implement efficiently
- Volatile keyword
 - Java fields can be declared volatile
 - Writing to a volatile variable ensures all local changes are made visible to other threads
 - x and y would have to be made volatile to fix code



What is the value of x?

Can be both 1 and 2!

java.util.concurrent.atomic

class AtomicInteger, AtomicReference<T>, ...

Represents a value

- method set(newValue)
 - has the effect of writing to a volatile variable
- method get()
 - returns the current value
- effectively an extension of volatile
- but what about atomicity???

Compare and Set (CAS)

- boolean compareAndSet(expectedValue, newValue)
 - If value doesn't equal expectedValue, return false
 - if equal, store newValue in value and return true
 - executes as a single atomic action!
 - supported by many processors
 - without requiring locks!

AtomicInteger n = new AtomicInteger(5); n.compareAndSet(3, 6); // return false – no change n.compareAndSet(5, 7); // returns true – now is 7

Incrementing with CAS

```
/** Increment n by one. Other threads use n too. */
public static void increment(AtomicInteger n) {
    int i = n.get();
    while (n.compareAndSet(i, i+1))
        i = n.get();
```

// AtomicInteger has increment methods doing this

Lock-Free Data Structures

- Usable by many concurrent threads
- □ using only atomic actions no locks!
- compare and swap is god here
- but it only atomically updates one variable at a time!

Let's implement one!