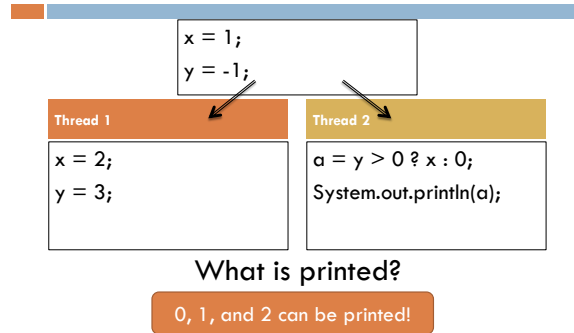


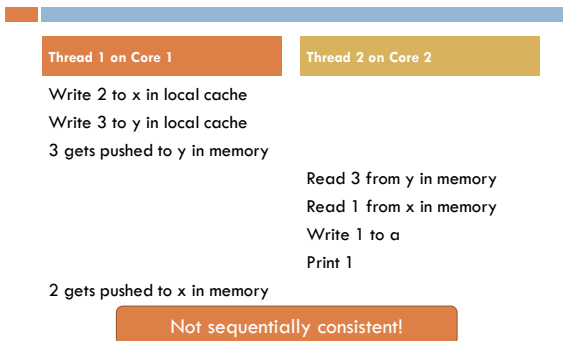
CONCURRENCY 2

CS 2110 – Spring 2016

Consistency



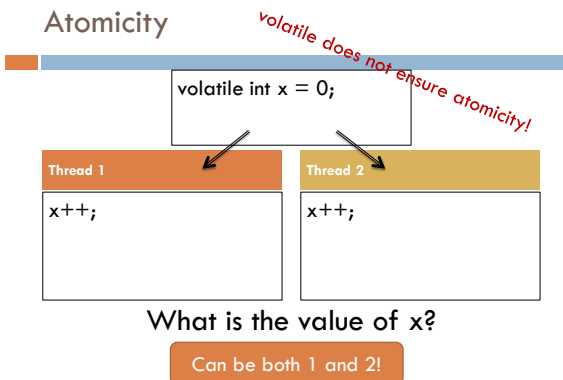
Consistency



Harsh Reality

- Sequential Consistency
 - There is an interleaving of the parallel operations that explains the observations and events
 - Currently unknown how to implement efficiently
- Volatile keyword
 - Java fields can be declared volatile
 - Writing to a volatile variable ensures all local changes are made visible to other threads
 - x and y would have to be made volatile to fix code

Atomicity



java.util.concurrent.atomic

- class AtomicInteger, AtomicReference<T>, ...
 - Represents a value
- method set(newValue)
 - has the effect of writing to a volatile variable
- method get()
 - returns the current value
- effectively an extension of volatile
- but what about atomicity???

Compare and Set (CAS)

- boolean `compareAndSet(expectedValue, newValue)`
 - If value doesn't equal `expectedValue`, return false
 - if equal, store `newValue` in value and return true
 - executes as a single atomic action!
 - supported by many processors
 - without requiring locks!

```
AtomicInteger n = new AtomicInteger(5);
n.compareAndSet(3, 6); // return false – no change
n.compareAndSet(5, 7); // returns true – now is 7
```

Incrementing with CAS

```
/** Increment n by one. Other threads use n too. */
public static void increment(AtomicInteger n) {
    int i = n.get();
    while (n.compareAndSet(i, i+1))
        i = n.get();
}

// AtomicInteger has increment methods doing this
```

Lock-Free Data Structures

- Usable by many concurrent threads
- using only atomic actions – no locks!
- compare and swap is god here
- but it only atomically updates one variable at a time!

Let's implement one!