

# USER-ORIENTED LANGUAGE DESIGN

CS 2110 – Spring 2016



# Decidability

# Wildcards

```
public Variable {
    boolean value;
    /** Add this to the list corresponding to value */
    public void addTo(
        List<? super Variable> trues,
        List<? super Variable> falses) {
        (value ? trues : falses).add(this);
    }
}
```

# Subtyping

```
class C<P> extends D<? super C<L<P>>>> { }
```

*Contrived*

$C\langle X \rangle \leq D\langle ? \text{ super } C\langle X \rangle \rangle$

Inheritance

$D\langle D\langle ? \text{ super } C\langle L\langle X \rangle \rangle \rangle \leq D\langle ? \text{ super } C\langle X \rangle \rangle$

**Infinite Proof  
of Subtyping**

Instantiation

Inheritance

$D\langle D\langle ? \text{ super } C\langle L\langle X \rangle \rangle \rangle \leq D\langle ? \text{ super } C\langle L\langle X \rangle \rangle \rangle$

Instantiation

$C\langle L\langle X \rangle \rangle \leq D\langle ? \text{ super } C\langle L\langle X \rangle \rangle \rangle$

# Efficiency

```
8 {
9     interface Eq<in T> {
10         bool equalTo(T other); }
11
12     interface List<out T> :
13         Eq<List<Eq<T>>> { }
14
15     interface Tree :
16         List<Tree> { }
17
```

# Restrictions

Termination  
Guaranteed

```
class C<P> extends D<D<? super C<L<P>>>> {
```

Contrived

## Inheritance Restriction

No use of ? super in the inheritance hierarchy

```
<P extends List<List<? super C<L<P>>>>>
```

Contrived

## Parameter Restriction

When constraining type parameters,

? super may only be used at covariant locations

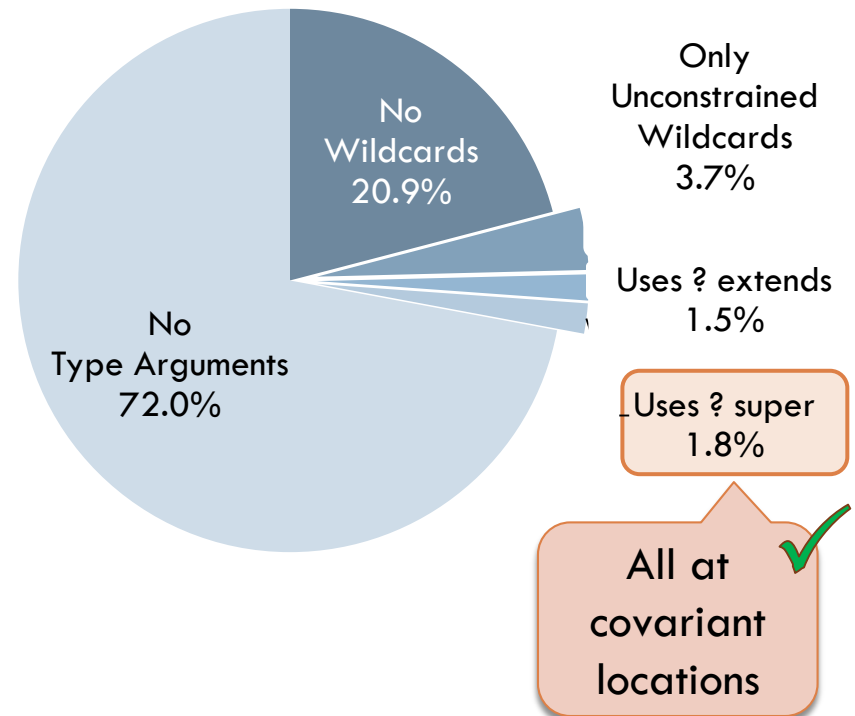
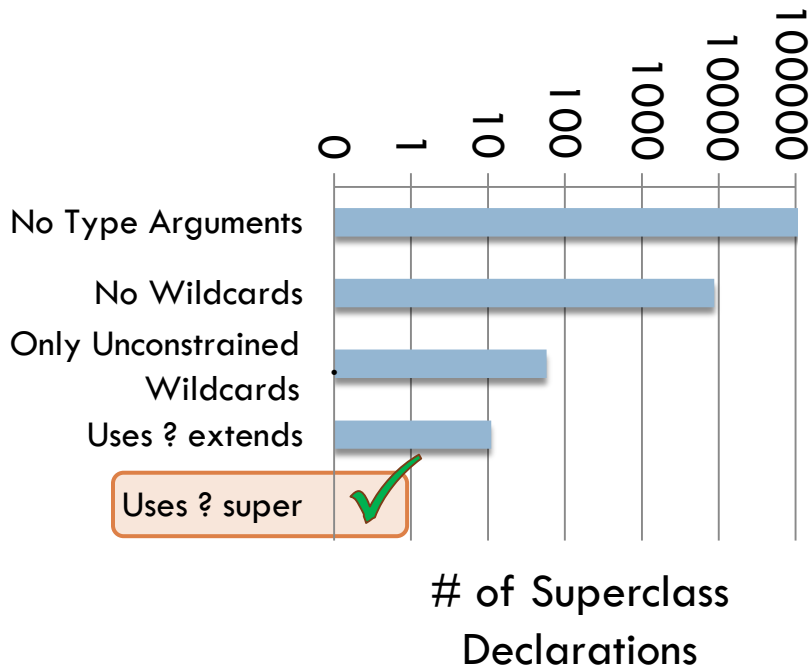
# Survey

## No Violations of Our Restrictions

9.2 Million Lines of Code Analyzed

### Wildcards in Inheritance

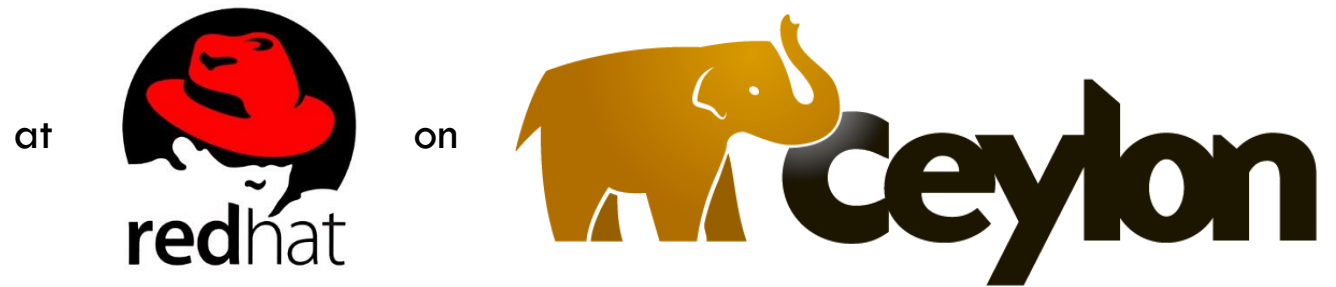
### Wildcards in Constraints



# Industry Collaborations



Gavin King



Andrey Breslav





# Materials and Shapes

- Material

- ▣ List, Integer, Property, Comparator

- Shape

- ▣ Comparable, Summable, Cloneable

No class/interface is both  
a material and a shape

13.5 Million Lines  
of Code Analyzed



Programmers are Humans

# Library Designer

- Want to provide a “separate” function
  - Inputs: middle, elems, smaller, bigger
  - Requirements:
    - Place all values in elems less than middle into smaller
    - Place all other values in elems into bigger
- Goals
  - Implement separate
  - Provide maximally flexible type signature

# Library User

- Goal
  - ▣ Place nonnegative values in “ints” into “positives”
- Context
  - ▣ “ignore” throws away all elements added to it
- Implementation
  - ▣ `separate(0, ints, ignore, positives);`

# User Types

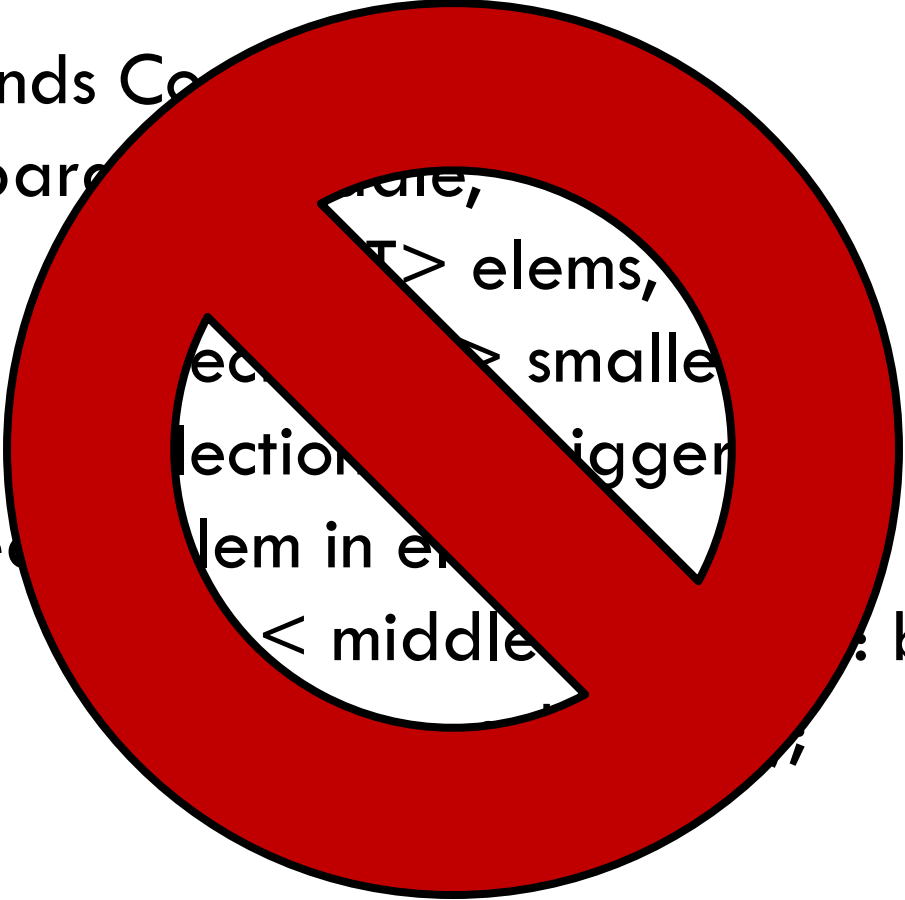
- `ints : Iterable<Integer>`
  - ▣ You can get things from iterables
- `positives : Collection<Integer>`
  - ▣ You can add things to collections
- `ignore : Collection<Object>`
  - ▣ You can add anything to it
- `Integer implements Comparable<Number>`
  - ▣ integers can be compared with any number

# Library Implementation

```
□ void separate(middle,
                elems,
                smaller,
                bigger) {
    foreach (elem in elems)
        (elem < middle ? smaller : bigger)
            .add(elem);
}
```

# Library Type

- <T extends Comparable>  
void separate



```
    int i = 0;
    while (i < elems.length) {
        if (elems[i] < smaller) {
            smaller = elems[i];
            bigger = elems[i];
        } else if (elems[i] > bigger) {
            bigger = elems[i];
        }
        i++;
    }
}
```

# Insufficient Flexibility

## □ Formals

### □ <T extends

Comparable<T>>

void separate(T middle,

Iterable<T> elems,

Collection<T> smaller,

Collection<T> bigger)

## Actuals

Integer



0



ints



ignore



positives





# Wildcards

```
□ <T extends Comparable<? super T>>
void separate(T middle,
              Iterable<? extends T> elems,
              Collection<? super T> smaller,
              Collection<? super T> bigger) {
    foreach (elem in elems)
        (elem < middle ? smaller : bigger)
            .add(elem);
}
```

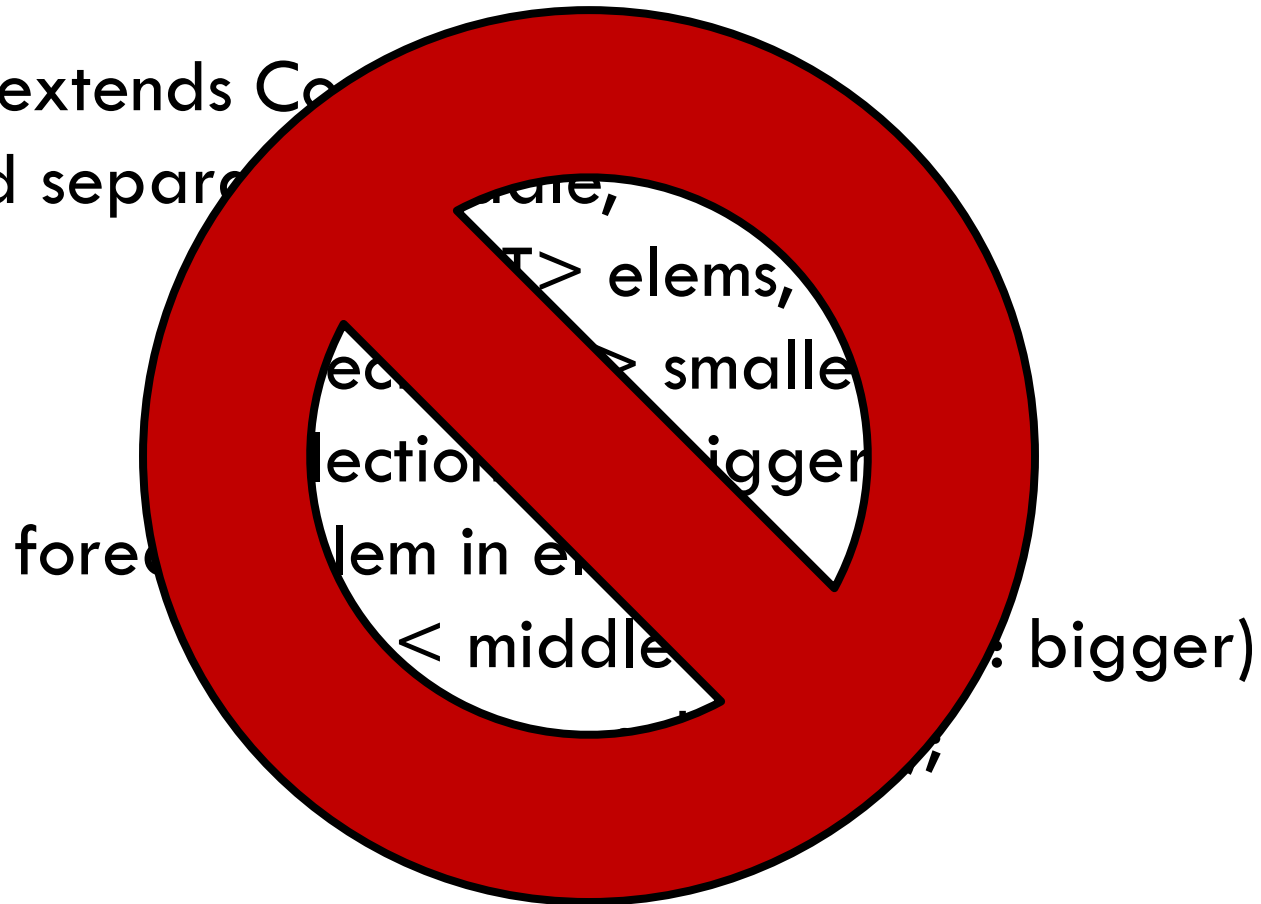
# Excessive Annotations

□ <T>

```
void flatten(  
Iterable<? extends Iterable<? extends T>>  
colls,  
           Collection<? super T> into) {  
    for (Iterable<? extends T> coll : colls)  
        for (T elem : coll)  
            into.add(elem);  
}
```

# Declaration-Site Variance

- <T extends Comparable>  
void separate



# Insufficient Flexibility

## □ Formals

### □ <T extends

Comparable<T>>

void separate(T middle,

Iterable<T> elems,

Collection<T> smaller,

Collection<T> bigger)

## Actuals

Integer



0



ints



ignore



positives



# Declaration-Site Variance Retry

```
□ <T extends Comparable<T>,
    U super T, V super T>
void separate(T middle,
    Iterable<T> elems,
    Collection<U> smaller,
    Collection<V> bigger) {
    foreach (elem in elems)
        (elem < middle ? smaller : bigger)
            .add(elem);
}
```

# Mixed-Site Variance

```
□ <T extends Comparable<T>>
void separate(T middle,
              Iterable<T> elems,
              Collection<in T> smaller,
              Collection<in T> bigger) {
    foreach (elem in elems)
        (elem < middle ? smaller : bigger)
            .add(elem);
}
```

# Use+Declaration-Site

- The two address orthogonal roles
  - ▣ declaration-site for class/interface designers
  - ▣ use-site for class/interface users
- How do the two interact?
  - ▣ given interface `Iterator<out T> {...}`
  - ▣ what does `Iterator<in Number>` mean?

# Expectations vs. Designs

	Java	Scala	Default	Layer	Join	Mixed
Has declaration-site variance	Red	Green	Green	Green	Green	Green
$C\langle\tau\rangle \prec: C\langle\text{out } \tau\rangle$	Green	Green	Red	Green	Green	Green
$C\langle\text{out } \tau\rangle \prec: C\langle\text{out } \tau'\rangle$ implies $\tau \prec: \tau'$	Red	Red	Green	Red	Red	Green
instance of $C\langle\text{out } \tau\rangle$ is instance of $C\langle\tau'\rangle$ for some $\tau' \prec: \tau$	Green	Green	Green	Green	Red	Green
instance of $C\langle\text{out } \tau\rangle$ allocated as $C\langle\tau'\rangle$ for some $\tau' \prec: \tau$	Green	Red	Green	Red	Red	Green
$C\langle\text{in } \tau \text{ out } \tau'\rangle$ is expressible	Red	Green	Green	Green	Red	Green
$C\langle\text{in } \tau \text{ out } \tau\rangle$ is valid implies $C\langle\tau\rangle$ is valid	Grey	Red	Green	Green	Grey	Green
$\text{Out}\langle\text{in } \tau\rangle$ is valid	Grey	Green	Green	Green	Green	Red
implicit constraints used in subtyping	Green	Red	Red	Red	Red	Red





# Principal Types

# Principal Type

- The principal type of an expression
  - ▣ a type for that expression that is better than all other types for that expression
- “Hello” has principal type String
  - ▣ “Hello” also has type Object, CharSequence, ...
  - ▣ String is a subtype of all those types
- A language has principal types
  - ▣ if every possible expression has a principal type

# Java

- `assertEquals(5, Integer.valueOf(5))`
  - ambiguous!
  - Is it two ints or two Integers?
- But the expression `5` is also an Integer
- And `Integer.valueOf(5)` is also an int
- Neither expression has a principal type

# Ambiguous Semantics

```
<P> List<P> singleton(P elem) {return ...}
<Q extends Comparable<?>> Q
{return javac
```

"Comparable" ✓

```
String typeName(Comparable<?> obj) {return obj.toString();}
```

"String" ✓

```
String typeName(Smith & Cartwright)
```

"Integer" ✓

```
String typeName(Calendar c) {return c.toString();}
```

```
String ambiguous(Socrates)
```

P ↦ Object

"Calendar" ✓

Q ↦ Calendar

# Use-Site Inferability Check

- `<T> List<T> singletonList(T) {...}`
- `var objs = singletonList("Hello");`
- `objs.add(5);`
  - ▣ fails to type check
  - ▣ `objs` is inferred to be an `List<String>`
  - ▣ needs to be an `List<Object>`

# Declaration-Site Inferability

- `<T> List<T> singletonList(T)`
  - T is not inferable because Array is invariant
  - `singletonList("Hello")`
    - could have type `List<String>` or `List<Object>`
    - no principal type
- `<T> Iterable<? extends T> singletonIterable(T)`
  - T is inferable because `? extends` is covariant
  - `singletonIterable("Hello")`
    - has type `Iterable<? extends String>`
    - which is subtype of `Iterable<? extends Object>`



# Gradual Types

# Goal

- Mix static and dynamic type systems
  - e.g. Java with JavaScript
- Requirements
  - no implicit insertions of wrappers
  - dynamic code is just static code minus types
  - stripping types preserves or improves semantics
  - static code can assume type annotations are true



# C#'s dynamic Type

- `bool Equal(object left, object right) {  
    return left == right;  
}`
- `Equal(0, 0)` returns false

# C#'s dynamic Type

- interface Getter<T> { T get(); }
- class Five : Getter<int>, Getter<string> {  
    int Getter<int>.get()  
        { return 5; }  
    double Getter<string>.get()  
        { return 5.0; }  
}
- void Print(Getter<int> getter) {  
    Console.WriteLine(getter.get());  
}

Crashes if changed  
to dynamic!

# C#'s dynamic Type

- `List<T> Snoc<T>(IEnumerable<T> start, T end) {  
    var elems = ToList(start);  
    elems.add(end);  
    return elems;  
}`
- `Snoc(Singleton("Hello"), 5)` works

Crashes if  
made dynamic!

# Prerequisite Language Properties

- Static Behavioral Subtyping
  - ▣ Using a more precise type for a subexpression improves the typability of the whole expression
- Decidability
  - ▣ Typing must be reliably doable at run time
- Principality
  - ▣ Every execution has a most precise typing