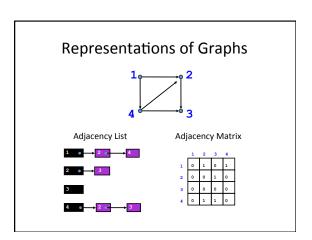


# **Graph Algorithms**

- Search
  - Depth-first search
  - Breadth-first search
- Shortest paths
  - Dijkstra's algorithm
- Minimum spanning trees
  - Prim's algorithm
  - Kruskal's algorithm



### Adjacency Matrix or Adjacency List?

- Definitions:
  - n =number of vertices
  - -m = number of edges
  - -d(u) = degree of u = number of edges leaving u
- Adjacency Matrix
  - Uses space O(n²)
  - Can iterate over all edges in time  $O(n^2)$
  - Can answer "Is there an edge from u to v?" in O(1) time
  - Better for dense graphs (lots of edges)
- Adjacency List
  - Uses space O(m + n)
  - Can iterate over all edges in time O(m + n)
  - Can answer "Is there an edge from u to v?" in O(d(u)) time
  - Better for sparse graphs (fewer edges)

# **Depth-First Search**

Given a graph and one of its nodes u
 (say node 1 below)



## Depth-First Search

- Given a graph and one of its nodes u
   (say node 1 below)
- We want to "visit" each node reachable from u
   (nodes 1, 0, 2, 3, 5)



There are many paths to some nodes.

How do we visit all nodes efficiently, without doing extra work?

# Depth-First Search

### boolean[] visited;

- Node u is visited means: visited[u] is true
- To visit u means to: set visited[u] to true
- Node v is REACHABLE from node u if there is a path (u, ..., v) in which all nodes of the path are unvisited.



Suppose all nodes are unvisited.

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Nodes REACHABLE from node 1: {1, 0, 2, 3, 5}

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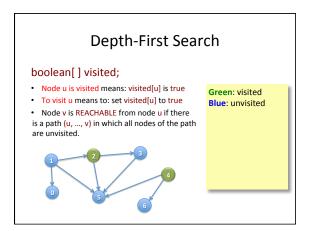
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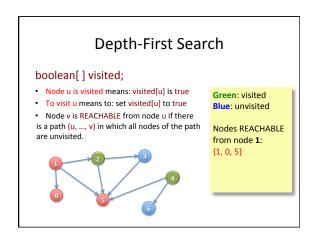


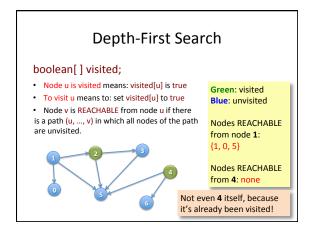
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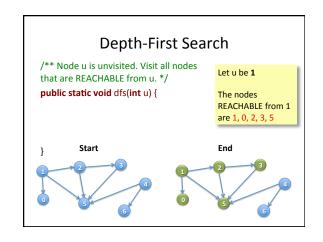
Nodes REACHABLE from node 1: {1, 0, 2, 3, 5}

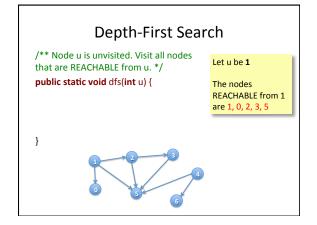
Nodes REACHABLE from 4: {4, 5, 6}

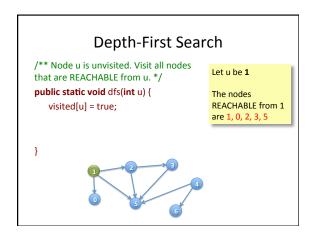


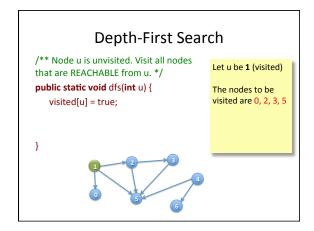


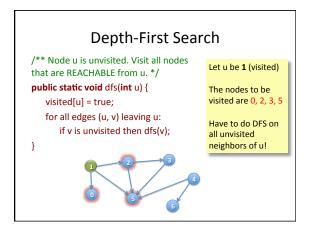


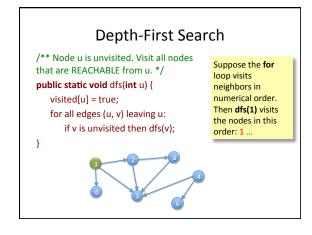


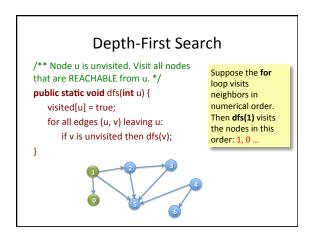


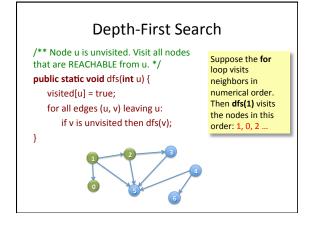


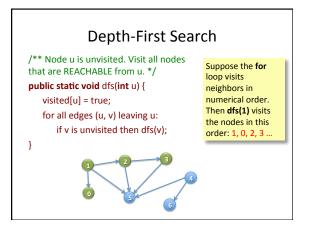












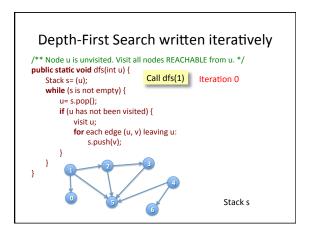
# /\*\* Node u is unvisited. Visit all nodes that are REACHABLE from u. \*/ public static void dfs(int u) { visited[u] = true; for all edges (u, v) leaving u: if v is unvisited then dfs(v); }

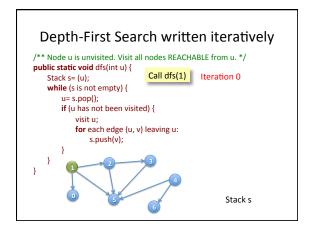
```
/** Node u is unvisited. Visit all nodes that are REACHABLE from u. */
public static void dfs(int u) {
  visited[u] = true;
  for all edges (u, v) leaving u:
      if v is unvisited then dfs(v);
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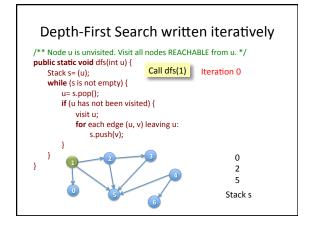
Suppose n nodes are REACHABLE along e edges (in total). What is
  • Worst-case execution?
  • Worst-case space?
```

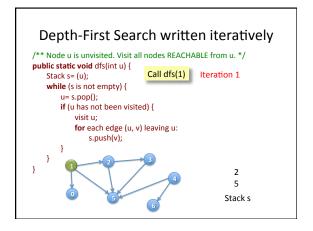
```
Depth-First Search
 /** Node u is unvisited. Visit all nodes
                                                That's all there is to
 that are REACHABLE from u. */
                                                basic DFS. You may
 public static void dfs(int u) {
                                                have to change it to
    visited[u] = true;
                                                fit a particular
                                                situation.
    for all edges (u, v) leaving u:
        if v is unvisited then dfs(v);
                                                If you don't have
                                                this spec and you
                                                do something
Example: Use different way (other than array
visited) to know whether a node has been visited
                                               different, it's
                                                probably wrong.
Example: We really haven't said what data
structures are used to implement the graph
```

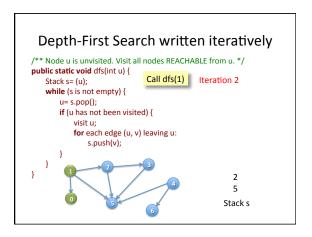
```
Depth-First Search in OO fashion
public class Node {
                                         Each node of the
   boolean visited;
                                         graph is an object
   List<Node> neighbors;
                                         of type Node
   /** This node is unvisited. Visit all nodes
      REACHABLE from this node */
   public void dfs() {
                                         No need for a
                                         parameter. The
      visited= true;
                                         object is the node.
      for (Node n: neighbors) {
         if (!n.visited) n.dfs();
```











### 

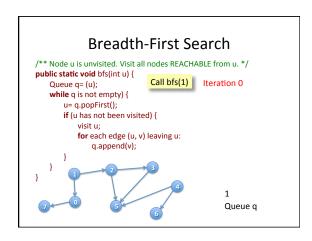
```
Depth-First Search written iteratively

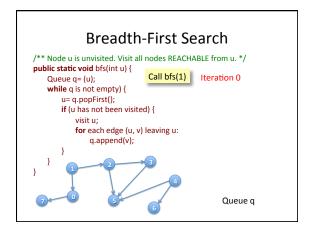
/** Node u is unvisited. Visit all nodes REACHABLE from u. */
public static void dfs(int u) {
    Stack s= (u);
    while (s is not empty) {
        u= s.pop();
    if (u has not been visited) {
        visit u;
        for each edge (u, v) leaving u:
            s.push(v);
    }
}

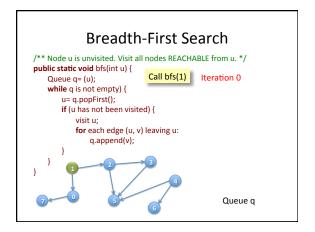
}

Stack s
```

```
Depth-First Search written iteratively
/** Node u is unvisited. Visit all nodes REACHABLE from u. */
public static void dfs(int u) {
                              Call dfs(1) Iteration 2
   Stack s= (u);
   while (s is not empty) {
       u= s.pop();
                                           Yes, 5 is put on the
       if (u has not been visited) {
                                           stack twice, once for
           visit u;
                                           each edge to it. It will
           for each edge (u, v) leaving u:
                                           be visited only once.
               s.push(v);
                                                      3
                                                     5
                                                     5
                                                   Stack s
```







```
## Node u is unvisited. Visit all nodes REACHABLE from u. */

public static void bfs(int u) {

Queue q= (u);

while q is not empty) {

u= q.popFirst();

if (u has not been visited) {

visit u;

for each edge (u, v) leaving u:

q.append(v);

}

0 2

Queue q
```

