

**CS2110 Spring 2014. Concluding Lecture:  
History, Correctness Issues, Summary**

Final review session: Fri, 9 May, 1:00–3:00. Phillips 101.

Final: 7:00–9:30PM, Monday, 12 May, Barton Hall

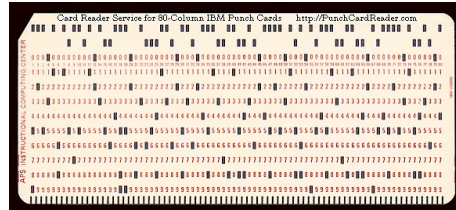
We hope to get you tentative course grades by Wednesday noon, but it may be later. You then visit the CMS and do the assignment to tell us whether you accept the grade or will take the final. There will be a message on the Piazza and the CMS about this when the tentative course grades are available.

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**CS2110 Spring 2014. Concluding Lecture:  
History, Correctness Issues, Summary**

Programming and computers:

Momentous changes since the 1940s –or since even the use of punch cards and attempt at automation ...



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**Punch cards**

**Jacquard loom**

**Loom still  
used in China**

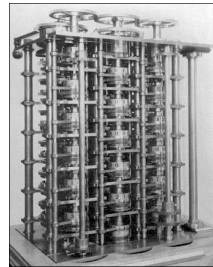


Mechanical loom invented by Joseph Marie Jacquard in 1801. Used the holes punched in pasteboard punch cards to control the weaving of patterns in fabric. Punch card corresponds to one row of the design. Based on earlier invention by French mechanic Falcon in 1728.

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Charles Babbage designed a “difference engine” in 1822

Compute mathematical tables for log, sin, cos, other trigonometric functions.



No electricity



The mathematicians doing the calculations were called **computers**.

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Oxford English Dictionary, 1971

**Computer:** one who computes; a calculator, rekenor. spec. a person employed to make calculations in an observatory, in surveying. etc.

1664: Sir T. Browne. The calendars of these computers.

1704. T. Swift. A very skillful computer.

1744. Walpole. Told by some nice computers of national glory.

1855. Brewster Newton. To pay the expenses of a computer for reducing his observations.

The mathematicians doing the calculations were called **computers**.

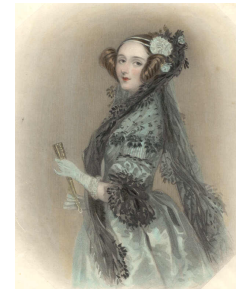
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Charles Babbage planned to use cards to store programs in his Analytical engine. (First designs of real computers, middle 1800s until his death in 1871.)

First programmer was Ada Lovelace, daughter of poet Lord Byron.

Privately schooled in math. One tutor was Augustus De Morgan.

The Right Honourable **Augusta Ada, Countess of Lovelace.**



**Herman Hollerith.**

His tabulating machines used in compiling the 1890 Census. Hollerith's patents were acquired by the Computing-Tabulating-Recording Co. Later became **IBM**.

The operator places each card in the reader, pulls down a lever, and removes the card after each punched hole is counted.



Hollerith 1890 Census Tabulator

**Computers, calculating the US census**



**1935-38. Konrad Zuse - Z1 Computer** **History of computers**

**1935-39.** John Atanasoff and Berry (grad student). Iowa State

**1944.** Howard Aiken & Grace Hopper **Harvard Mark I Computer**

**1946.** John Presper Eckert & John W. Mauchly **ENIAC 1 Computer** 20,000 vacuum tubes later ...

**1947-48** The Transistor, at Bell-labs.

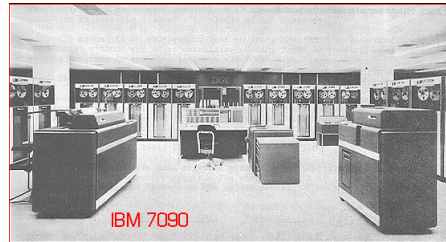
**1953. IBM. the IBM 701.**



How did Gries get into Computer Science?

1959. Took his only computer course. Senior, Queens College.

1960. Mathematician-programmer at the US Naval Weapons Lab in Dahlgren, Virginia.



IBM 7090

1960. Mathematician-programmer at the US Naval Weapons Lab in Dahlgren, Virginia.

Programmed in Fortran and IBM 7090 assembly language

```
if (SEX == 'M') MALES= MALES + 1;
else FEMALES= FEMALES + 1;
```

```
CLI SEX,'M'   Male?
BNO IS_FEM   If not, branch around
L 7,MALES   Load MALES into register 7;
LA 7,1(,7)   add 1;
ST 7,MALES   and store the result
B GO_ON     Finished with this portion
IS_FEM L 7,FEMALES If not male, load FEMALES into register 7;
LA 7,1(,7)   add 1;
ST 7,FEMALES and store
GO_ON EQU *
```

1960: Big Year for Programming Languages

**LISP** (List Processor): McCarthy, MIT (moved to Stanford). First functional programming language. No assignment statement. Write everything as recursive functions.

**COBOL** (Common Business-Oriented Language). Became most widely used language for business, data processing.

**ALGOL** (Algorithmic Language). Developed by an international team over a 3-year period. McCarthy was on it, John Backus was on it (developed Fortran in mid 1950's). Gries's soon-to-be PhD supervisor, Fritz Bauer of Munich, led the team.

1959. Took his only computer course. Senior, Queens College.  
 1960. Mathematician-programmer at the US Naval Weapons Lab in Dahlgren, Virginia.  
 1962. Back to grad school, in Math, at University of Illinois  
 Graduate Assistantship: Help two Germans write the ALCOR-Illinois 7090 Compiler.  
 John Backus, FORTRAN, mid 1950's: 30 people years  
 This compiler: 6 ~people-years  
 Today, CS compiler writing course: 2 students, one semester  
 1963-66 Dr. rer. nat. in Math in Munich Institute of Technology  
 1966-69 Asst. Professor, Stanford CS  
 1969- Cornell!

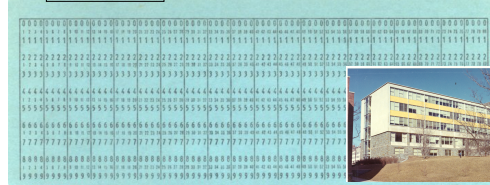
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**Late 1960s**



IBM 360

**Mainframes**



Write programs on IBM "punch cards. Deck of cards making up a program trucked to Langmuir labs by the airport 2-3 times a day; get them back, with output, 3-4 hours later



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**About 1973. BIG STEP FORWARD**

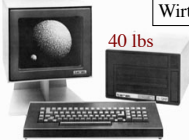
1. Write program on punch cards.
2. Wait in line (20 min) to put cards in card reader in Upson basement
3. Output comes back in 5 minutes

**About 1973. BIG STEP FORWARD**

Switched to using the programming language Pascal, developed by Niklaus Wirth at Stanford.

**About 1979. Teraks**

Prof. Tim Teitelbaum sees opportunity. He and grad student Tom Reps develop "Cornell Program Synthesizer". Year later, Cornell uses Teraks in its prog course.



40 lbs

November 1981, Terak with 56K RAM, one floppy drive: \$8,935.

Want 10MB hard drive? \$8,000 more

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**1983-84**

Switched to Macintosh in labs

**Late 1980s**

Put fifth floor addition on Upson. We made the case that our labs were in our office and therefore we need bigger offices.

**1980s**

CS began getting computers on their desks.

**Nowadays**

Everybody has a computer in their office.

**2014**

Moved into Gates Hall!

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**Programming languages. Dates approximate**

Year	Major languages	Teach at Cornell
1956's	Fortran	
1960	Algol, LISP, COBOL	
1965	PL/I	PL/C (1969)
1970	C	
1972	Pascal	
1980's	Smalltalk (object-oriented)	Pascal (1980's)
1980's (late)	C++	
1996	Java	C and C++
1998		Java / Matlab

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**The NATO Software Engineering Conferences**

homepages.cs.ncl.ac.uk/brian.randell/NATO/

7-11 Oct 1968, Garmisch, Germany  
 27-31 Oct 1969, Rome, Italy

Download Proceedings, which have transcripts of discussions. See photographs.

**Software crisis:**


Academic and industrial people. Admitted for first time that they did not know how to develop software efficiently and effectively.



**Software Engineering, 1968**



Next 10-15 years: intense period of research on software engineering, language design, proving programs correct, etc. 19



Software Engineering, 1968 20

**During 1970s, 1980s, intense research on**  
 How to prove programs correct,  
 How to make it practical,  
 Methodology for developing algorithms

The way we understand recursive methods is based on that methodology.  
 Our understanding of and development of loops is based on that methodology.

Throughout, we try to give you thought habits and strategies to help you solve programming problems for effectively, e.g.  
**Write good method specs.**  
**Keep methods short.**  
**Use method calls to eliminate nested loops.**  
**Put local variable declarations near first use.**

Mark Twain: Nothing needs changing so much as the habits of **others**.

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The way we understand recursive methods is based on that methodology.  
 Our understanding of and development of loops is based on that methodology.

Simplicity is key:  
 Learn not only to simplify, learn not to complicate.

Separate concerns, and focus on one at a time.

Develop and test incrementally.

Throughout, we try to give you thought habits to help you solve programming problems for effectively

Don't solve a problem until you know what the problem is (give precise and thorough specs).

Learn to read a program at different levels of abstraction.

Use methods and method calls so that you don't have nested loops

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**Simplicity and beauty: keys to success**

CS has its field of computational complexity. Mine is computational simplicity,

**David Gries**

Inside every large program is a little program just trying to come out. **Tony Hoare**


Beauty is our Business. **Edsger Dijkstra**

**CS professor's non-dilemma**  
 I do so want students to see beauty and simplicity.  
 A language used just has to be one only with that property.  
 Therefore, and most reasonably, I will not and do not teach C.  
**David Gries**

**Admonition a little Grook**  
 In correctness concerns one must be immersed.  
 To use only testing is simply accursed. 23

**Correctness of programs, the teaching of programming**

simplicity  
 elegance  
 perfection  
 intellectual honesty



Edsger W. Dijkstra Sir Tony Hoare

**Dijkstra:** The competent programmer is fully aware of the limited size of his own skull, so he approaches the programming task in full humility, and among other things, he avoids clever tricks like the plague.

**Hoare:** Two ways to write a program:  
 (1) Make it so simple that there are obviously no errors.  
 (2) Make it so complicated that there are no obvious errors. 24

**Axiomatic Basis for Computer Programming.**  
**Tony Hoare, 1969**

Provide a definition of programming language statements not in terms of how they are executed but in terms of proving them correct.

{precondition P}  
 Statement S  
 {Postcondition Q}

Meaning: If P is true, then execution of S is guaranteed to terminate and with Q true

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**Assignment statement  $x = e$ ;**

{true}	{x+1 >= 0}	{2*x = 82}
x = 5;	x = x + 1;	x = 2*x;
{x = 5}	{x >= 0}	{x = 82}

**Definition of notation:**

$P[x:=e]$  (read P with x replaced by e) stands for a copy of expression P in which each occurrence of x is replaced by e

Example:  $(x >= 0)[x:= x+1] = x+1 >= 0$

**Definition of the assignment statement:**

{P[x:= e]}  
 x = e;  
 {P}

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**Assignment statement  $x = e$ ;**

**Definition of the assignment statement:**

{P[x:= e]}  
 x = e;  
 {P}

{x+1 >= 0}	{2*x = 82}
x = x + 1;	x = 2*x;
{x >= 0}	{x = 82}

{2.0xy + z = (2.0xy + z)/6}	x = x/6
x = 2.0*x*y + z;	
{x = x/6}	2.0xy + z = (2.0xy + z)/6

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**If statement defined as an “inference rule”:**

**Definition of if statement: If**  
 {P && B} ST {Q} and  
 {P && !B} SF {Q}

**Then**  
 {P}  
**if (B) ST**  
**else SF**  
 {Q}

The then-part, ST, must end with Q true  
 The else-part, SF, must end with Q true

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**Hoare’s contribution 1969:**

Axiomatic basis: Definition of a language in terms of how to prove a program correct.

But it is difficult to prove a program correct after the fact. How do we develop a program and its proof hand-in-hand?

Dijkstra showed us how to do that in 1975.

His definition, called “weakest preconditions” is defined in such a way that it allows us to “calculate” a program and its proof of correctness hand-in-hand, with the proof idea leading the way.

Dijkstra: *A Discipline of Programming*. Prentice Hall, 1976. A research monograph

Gries: *The Science of Programming*. Springer Verlag, 1981. Undergraduate text.

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**How to prove concurrent programs correct.**  
**Use the principle of non-interference**

Thread T1  
 {P0}  
 S1;  
 {P1}  
 S2;  
 {P2}  
 ...  
 Sn;  
 {Pn}

Thread T2  
 {Q0}  
 Z1;  
 {Q1}  
 Z2;  
 {Q2}  
 ...  
 Zm;  
 {Qm}

We have a proof that T1 works in isolation and a proof that T2 works in isolation. But what happens when T1 and T2 execute simultaneously, operating on the same variables?

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Thread T1 {P0} S1; {P1} S2; {P2} ... Sn; {Pn}	Thread T2 {Q0} Z1; {Q1} Z2; {Q2} ... Zm; {Qm}	<p><b>How to prove concurrent programs correct.</b></p> <p>Turn what previously seemed to be an exponential problem, looking at all executions, into a problem of size <math>n*m</math>.</p>
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Prove that execution of T1 does not interfere with the **proof** of T2, and vice versa.  
Basic notion: Execution of Si does not falsify an assertion in T2:  
e.g. {Pi && Q1} S2 {Z2}

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Thread T1 {P0} S1; {P1} S2; {P2} ... Sn; {Pn}	Thread T2 {Q0} Z1; {Q1} Z2; {Q2} ... Zm; {Qm}	<p><b>Interference freedom.</b>  Susan Owicki's Cornell thesis, under Gries, in 1975.</p> <p>A lot of progress since then! But still, there are a lot of hard issues to solve in proving concurrent programs correct in a practical manner.</p>
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Prove that execution of T1 does not interfere with the **proof** of T2, and vice versa.  
Basic notion: Execution of Si does not falsify an assertion in T2:  
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