

# PRIORITY QUEUES AND HEAPS

Lecture 18  
CS2110 Fall 2013

# Readings and Homework

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- Read Chapter 26 to learn about heaps
- Salespeople often make matrices that show all the great features of their product that the competitor's product lacks. Try this for a heap versus a BST. First, try and sell someone on a BST: List some desirable properties of a BST that a heap lacks. Now be the heap salesperson: List some good things about heaps that a BST lacks. Can you think of situations where you would favor one over the other?



***With ZipUltra heaps, you've got it made in the shade my friend!***

# The Bag Interface

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## □ A Bag:

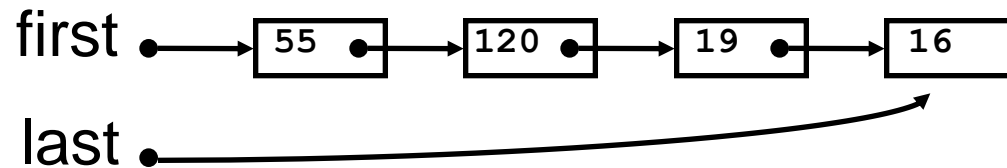
```
interface Bag<E> {  
    void insert(E obj);  
    E extract(); //extract some element  
    boolean isEmpty();  
}
```

Examples: Stack, Queue, PriorityQueue

# Stacks and Queues as Lists

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- Stack (LIFO) implemented as list
  - **insert()**, **extract()** from front of list
- Queue (FIFO) implemented as list
  - **insert()** on back of list, **extract()** from front of list
- All **Bag** operations are  $O(1)$



# Priority Queue

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- A **Bag** in which data items are **Comparable**
- *lesser* elements (as determined by **compareTo ( )**) have *higher* priority
- **extract ( )** returns the element with the highest priority = least in the **compareTo ( )** ordering
- break ties arbitrarily

# Priority Queue Examples

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- Scheduling jobs to run on a computer
  - default priority = arrival time
  - priority can be changed by operator
- Scheduling events to be processed by an event handler
  - priority = time of occurrence
- Airline check-in
  - first class, business class, coach
  - FIFO within each class

# java.util.PriorityQueue<E>

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```
boolean add(E e) {...} //insert an element (insert)
void clear() {...} //remove all elements
E peek() {...} //return min element without removing
                //(null if empty)
E poll() {...} //remove min element (extract)
                //(null if empty)
int size() {...}
```

# Priority Queues as Lists

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- Maintain as **unordered list**
  - **insert ()** puts new element at front –  $O(1)$
  - **extract ()** must search the list –  $O(n)$
- Maintain as **ordered list**
  - **insert ()** must search the list –  $O(n)$
  - **extract ()** gets element at front –  $O(1)$
- In either case,  $O(n^2)$  to process  $n$  elements

Can we do better?



# Important Special Case

9

- Fixed number of priority levels  $0, \dots, p - 1$
- FIFO within each level
- Example: airline check-in
  
- **insert ()** – insert in appropriate queue –  $O(1)$
- **extract ()** – must find a nonempty queue –  $O(p)$

# Heaps

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- A *heap* is a concrete data structure that can be used to implement priority queues
- Gives better complexity than either ordered or unordered list implementation:
  - **insert ()** :  $O(\log n)$
  - **extract ()** :  $O(\log n)$
- $O(n \log n)$  to process  $n$  elements
- Do not confuse with *heap memory*, where the Java virtual machine allocates space for objects – different usage of the word *heap*

# Heaps

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- Binary tree with data at each node
- Satisfies the *Heap Order Invariant*:

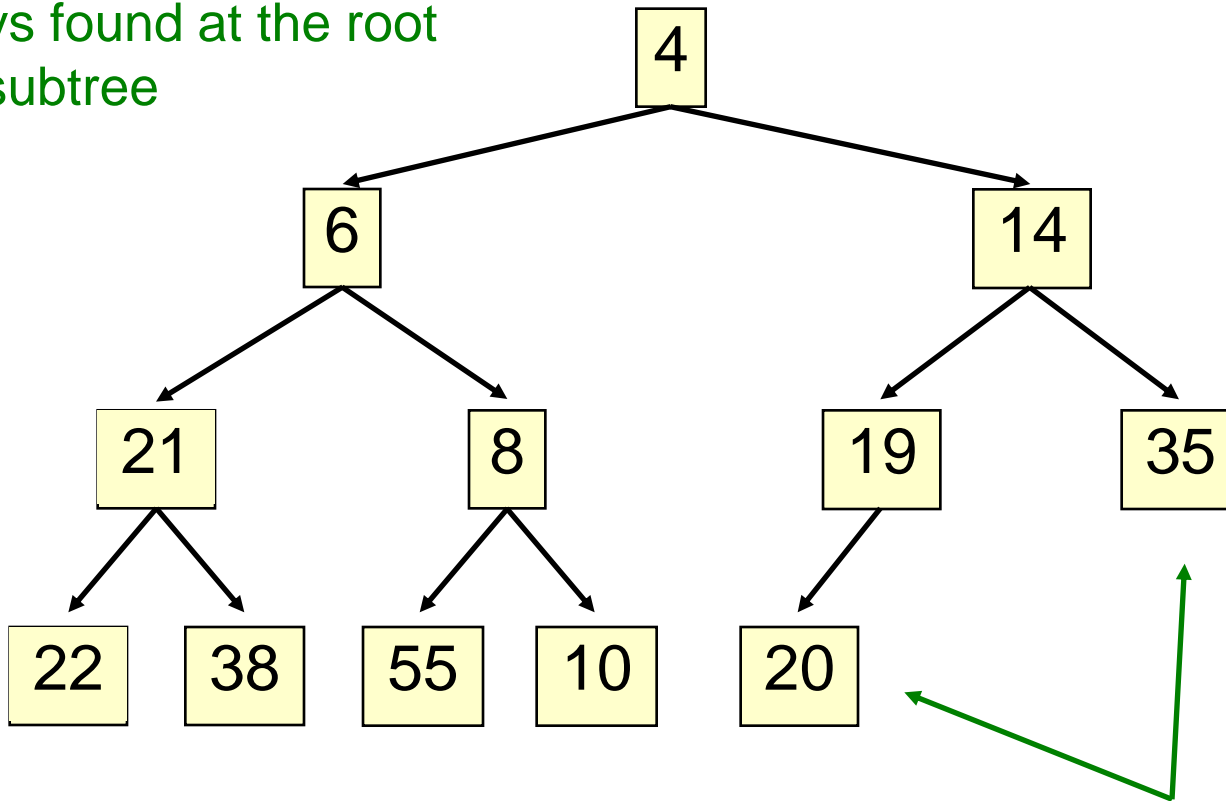
The least (highest priority) element of any subtree is found at the root of that subtree

- Size of the heap is “fixed” at  $n$ . (But can usually double  $n$  if heap fills up)

# Heaps

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Least element in any subtree  
is always found at the root  
of that subtree



Note:  $19, 20 < 35$ : we can often find smaller elements deeper in the tree!

# Examples of Heaps

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- Ages of people in family tree
  - parent is always older than children, but you can have an uncle who is younger than you
- Salaries of employees of a company
  - bosses generally make more than subordinates, but a VP in one subdivision may make less than a Project Supervisor in a different subdivision

# Balanced Heaps

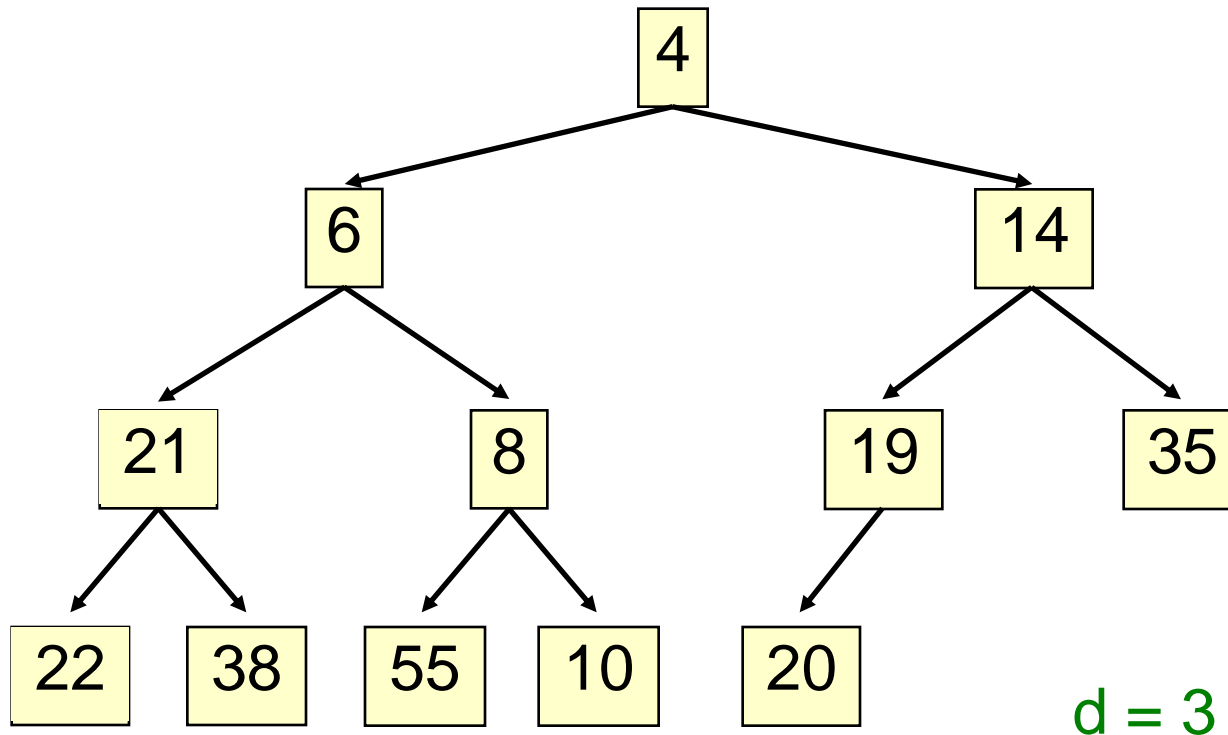
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These add two restrictions:

1. Any node of depth  $< d - 1$  has exactly 2 children, where  $d$  is the height of the tree
  - implies that any two maximal paths (path from a root to a leaf) are of length  $d$  or  $d - 1$ , and the tree has at least  $2^d$  nodes
- All maximal paths of length  $d$  are to the left of those of length  $d - 1$

# Example of a Balanced Heap

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# Store in an ArrayList or Vector

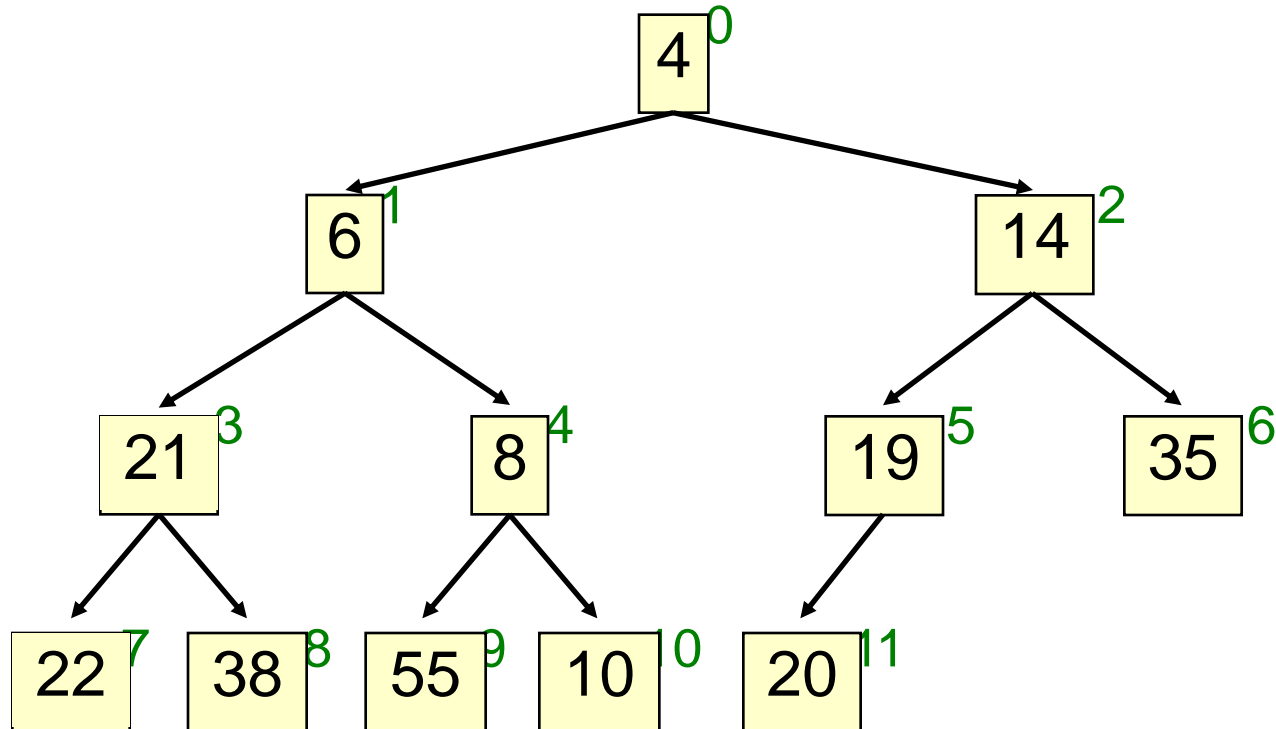
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- Elements of the heap are stored in the array in order, going across each level from left to right, top to bottom
- The children of the node at array index  $n$  are found at  $2n + 1$  and  $2n + 2$
- The parent of node  $n$  is found at  $(n - 1)/2$



# Store in an ArrayList or Vector

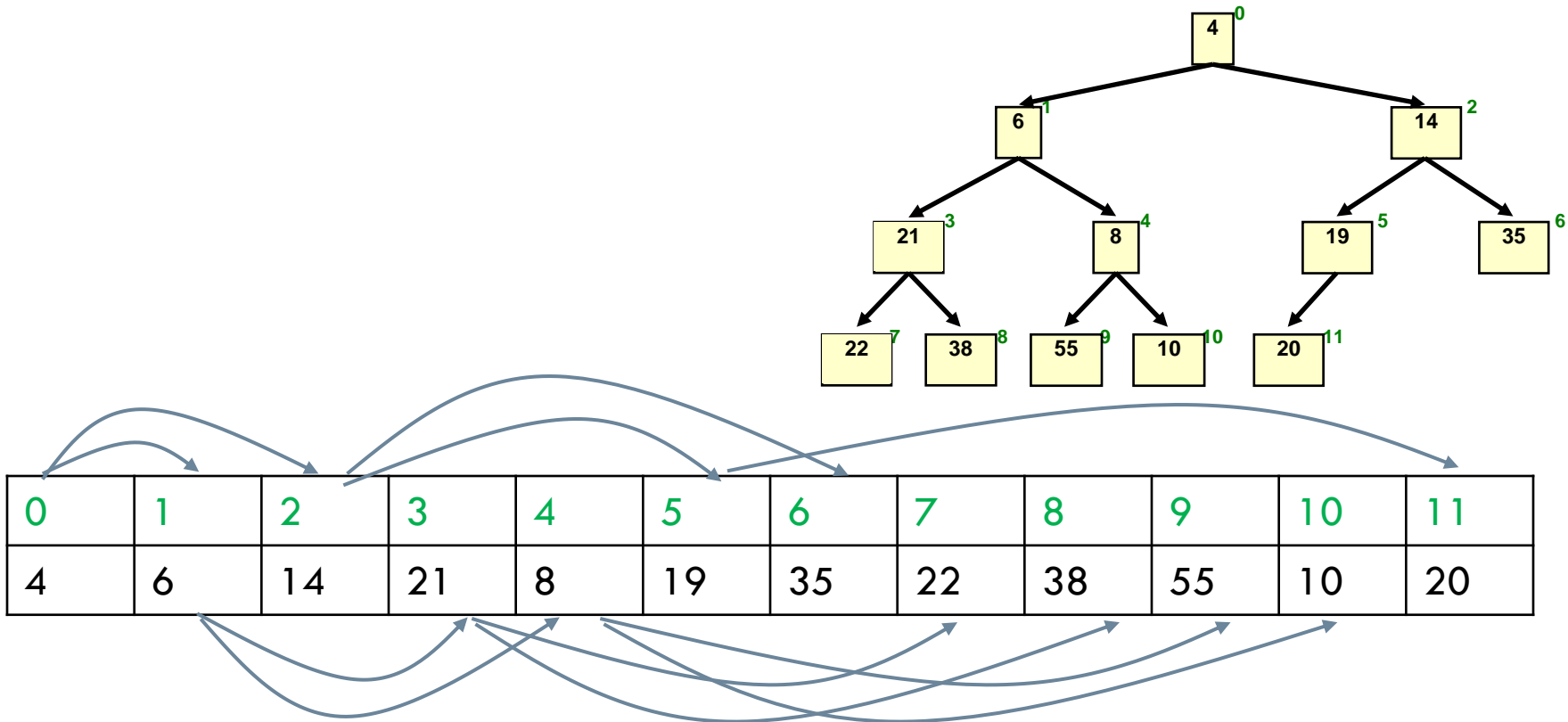
17



children of node  $n$  are found at  $2n + 1$  and  $2n + 2$

# Store in an ArrayList or Vector

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children of node  $n$  are found at  $2n + 1$  and  $2n + 2$

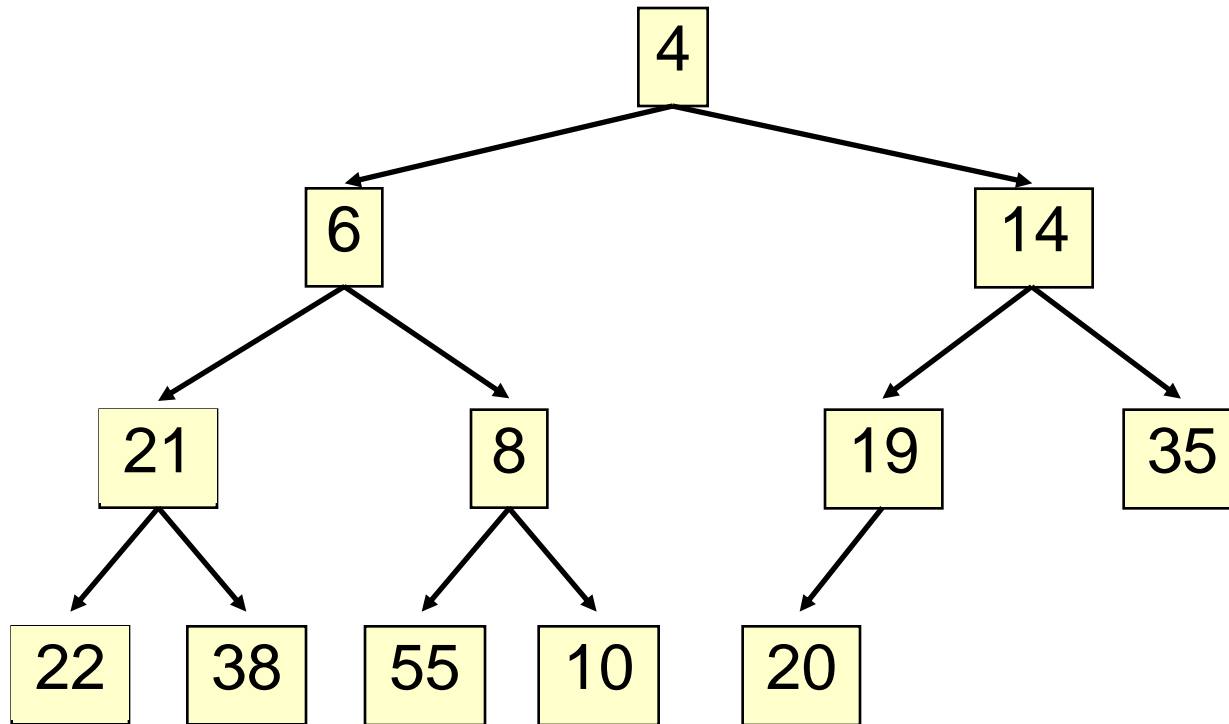
# insert ()

19

- Put the new element at the end of the array
- If this violates heap order because it is smaller than its parent, swap it with its parent
- Continue swapping it up until it finds its rightful place
- The heap invariant is maintained!

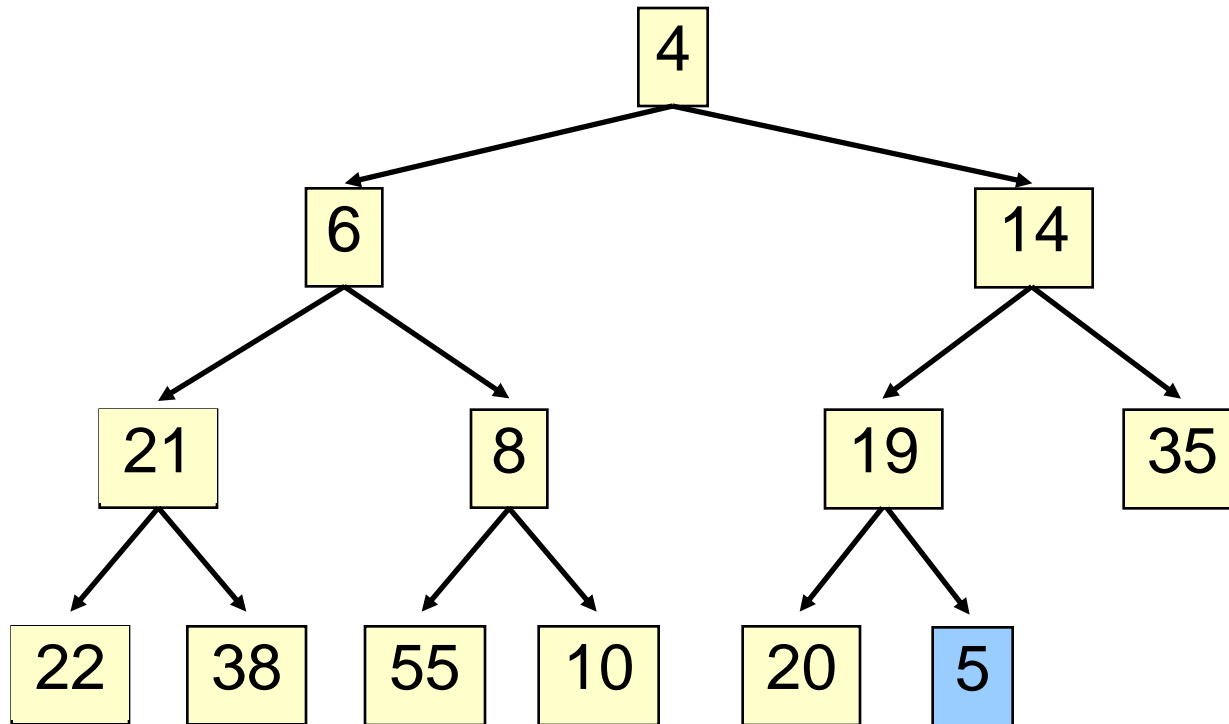
# insert ()

20



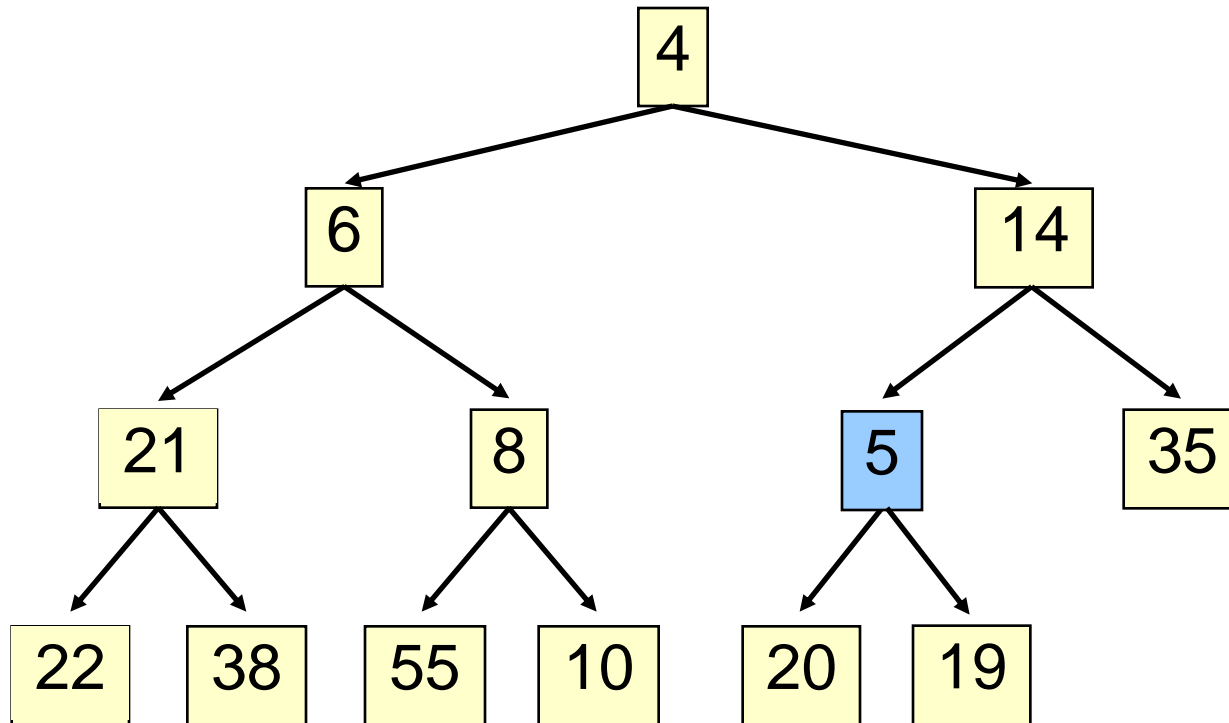
# insert()

21



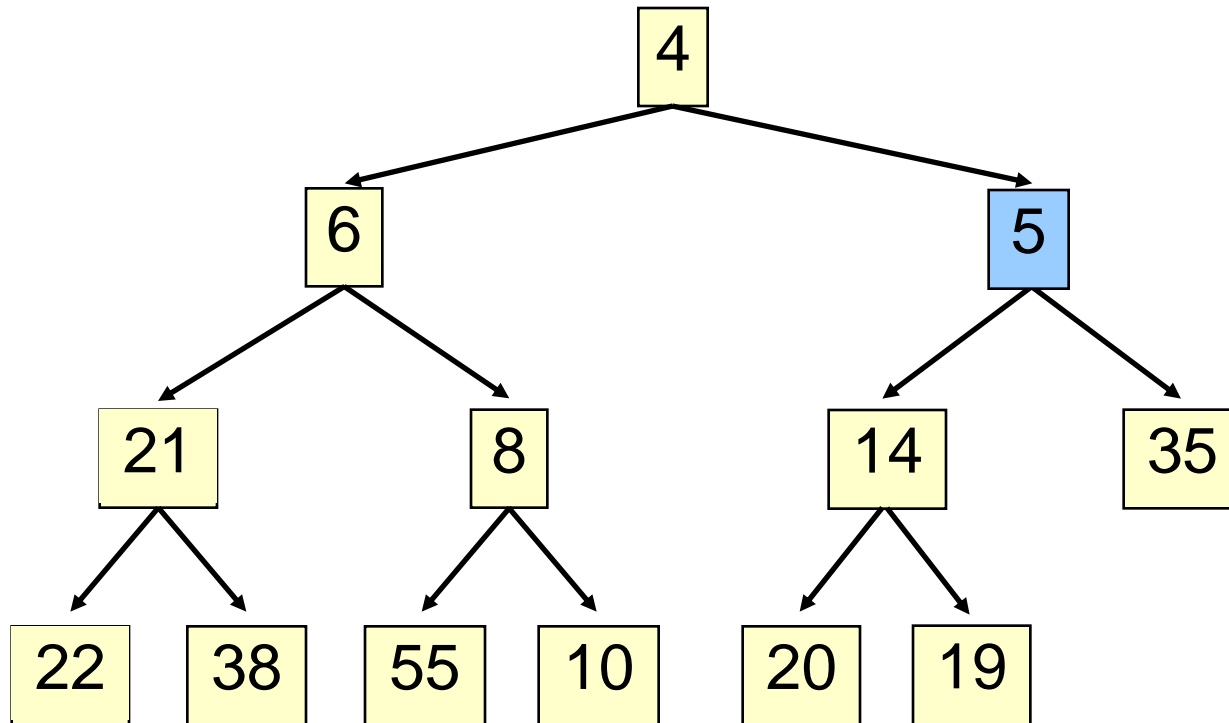
# insert()

22



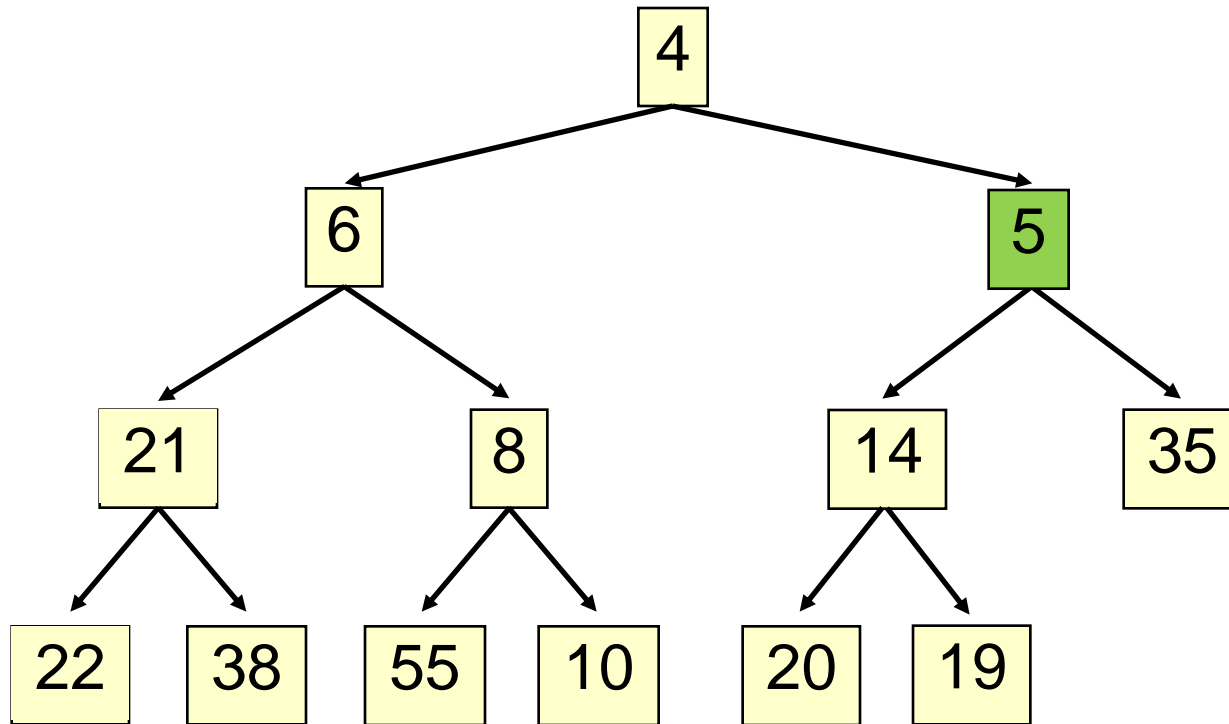
# insert()

23



# insert()

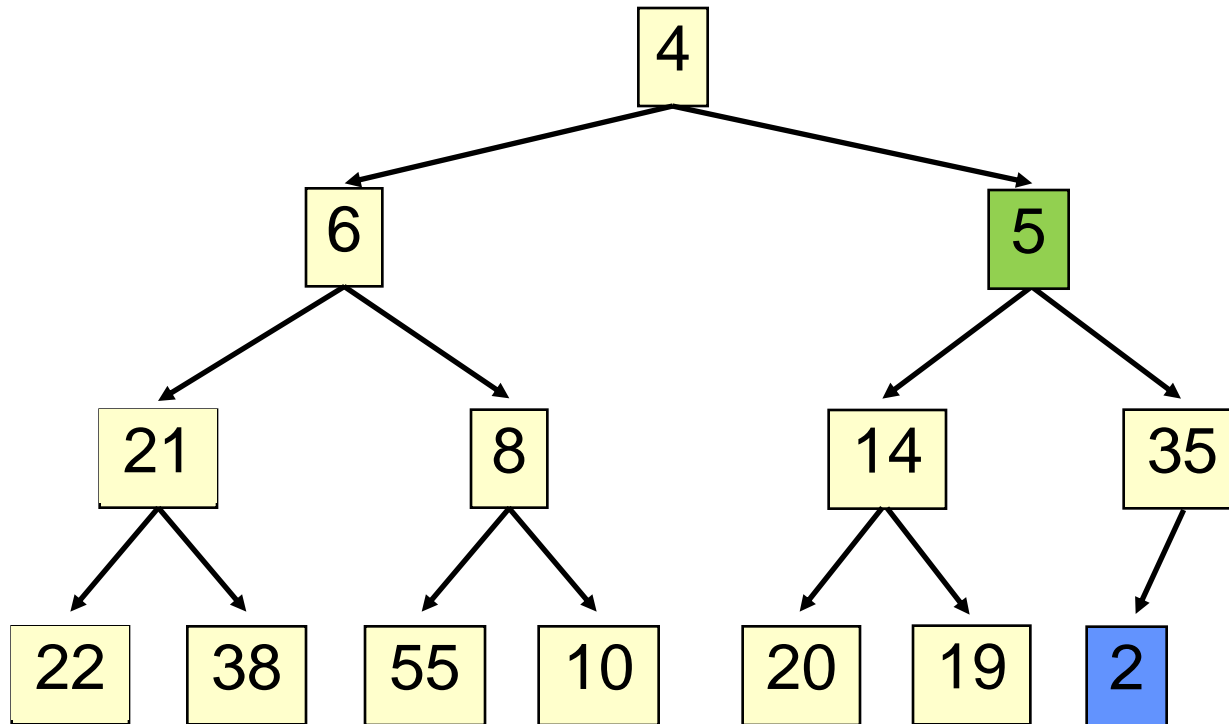
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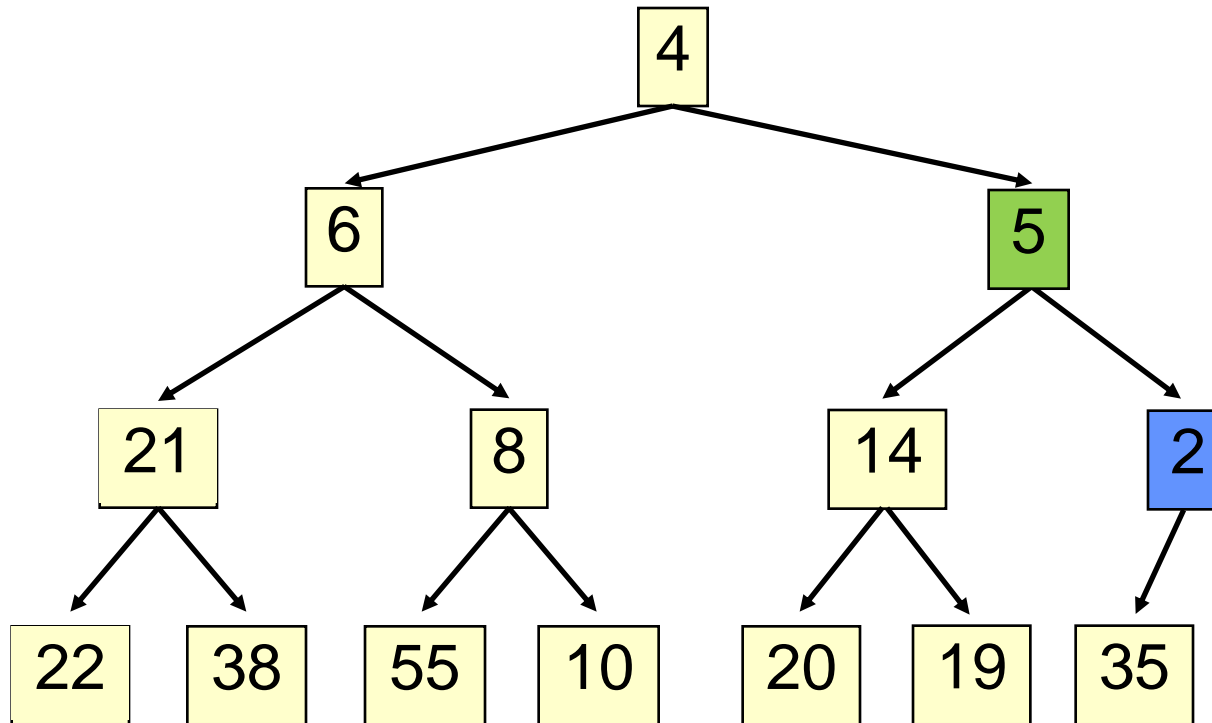
# insert()

25



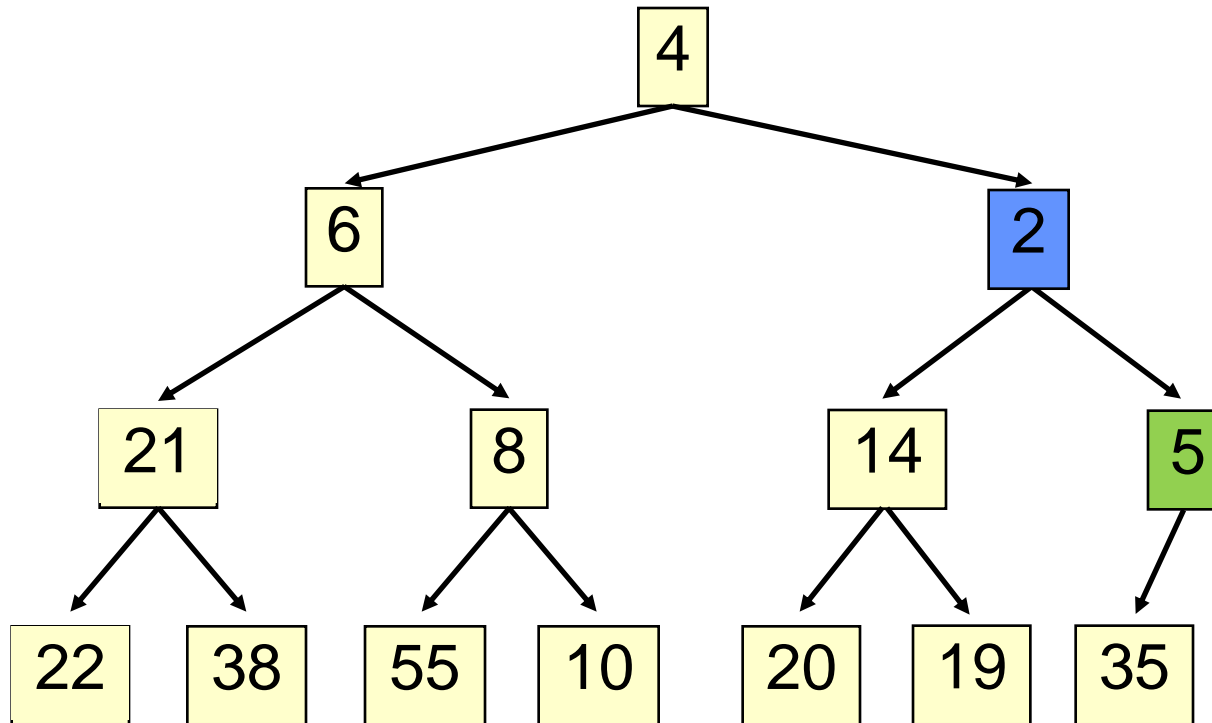
# insert()

26



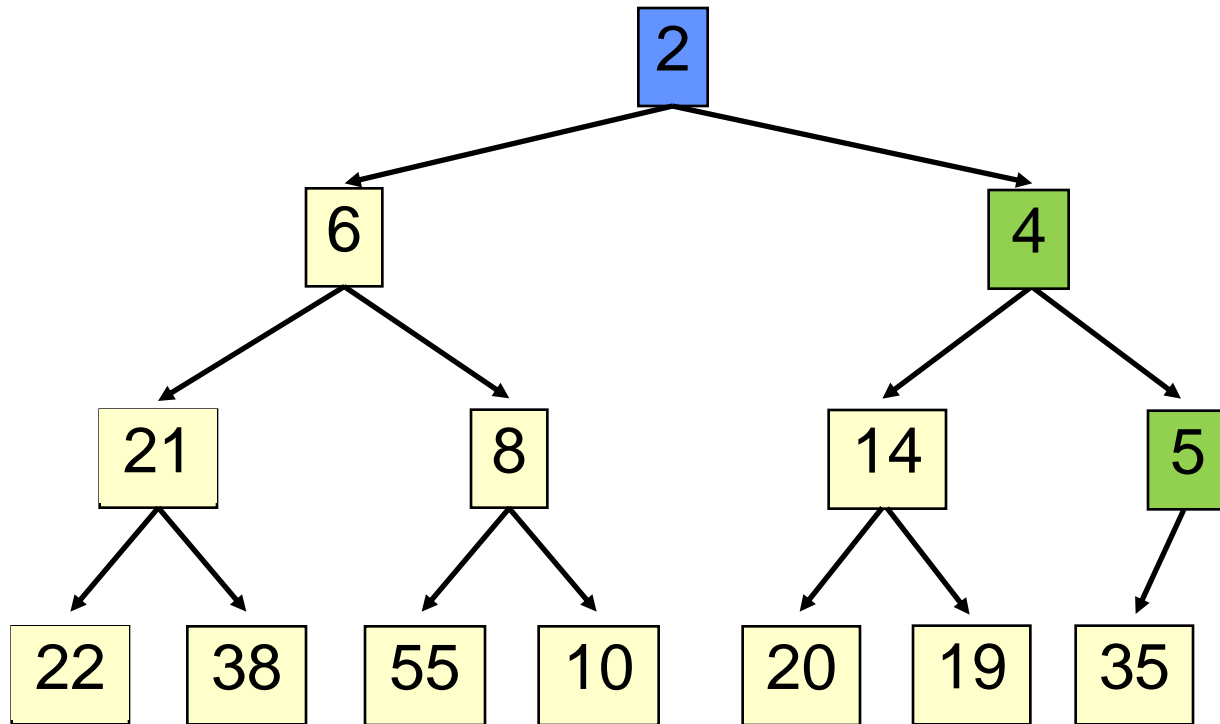
# insert()

27



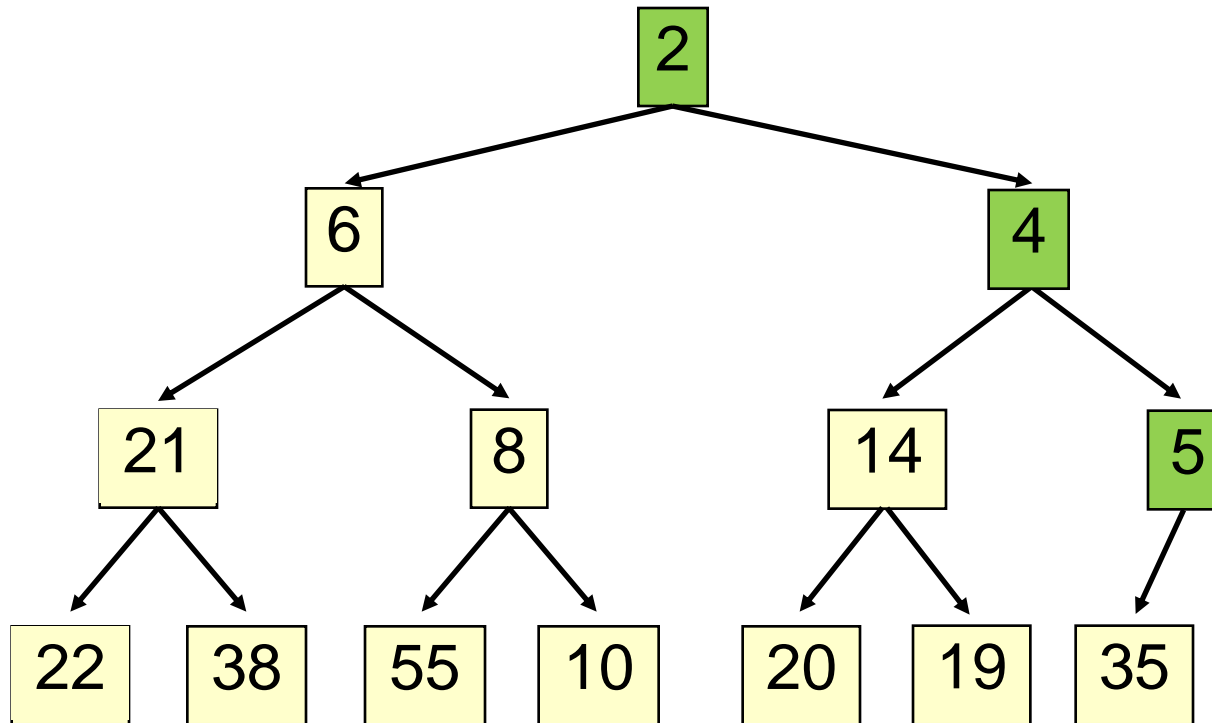
# insert ()

28



# insert()

29



# insert()

30

- Time is  $O(\log n)$ , since the tree is balanced
  - size of tree is exponential as a function of depth
  - depth of tree is logarithmic as a function of size

# insert()

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```
class PriorityQueue<E> extends java.util.Vector<E> {  
  
    public void insert(E obj) {  
        super.add(obj); //add new element to end of array  
        rotateUp(size() - 1);  
    }  
  
    private void rotateUp(int index) {  
        if (index == 0) return;  
        int parent = (index - 1)/2;  
        if (elementAt(parent).compareTo(elementAt(index)) <= 0)  
            return;  
        swap(index, parent);  
        rotateUp(parent);  
    }  
}
```

# extract ()

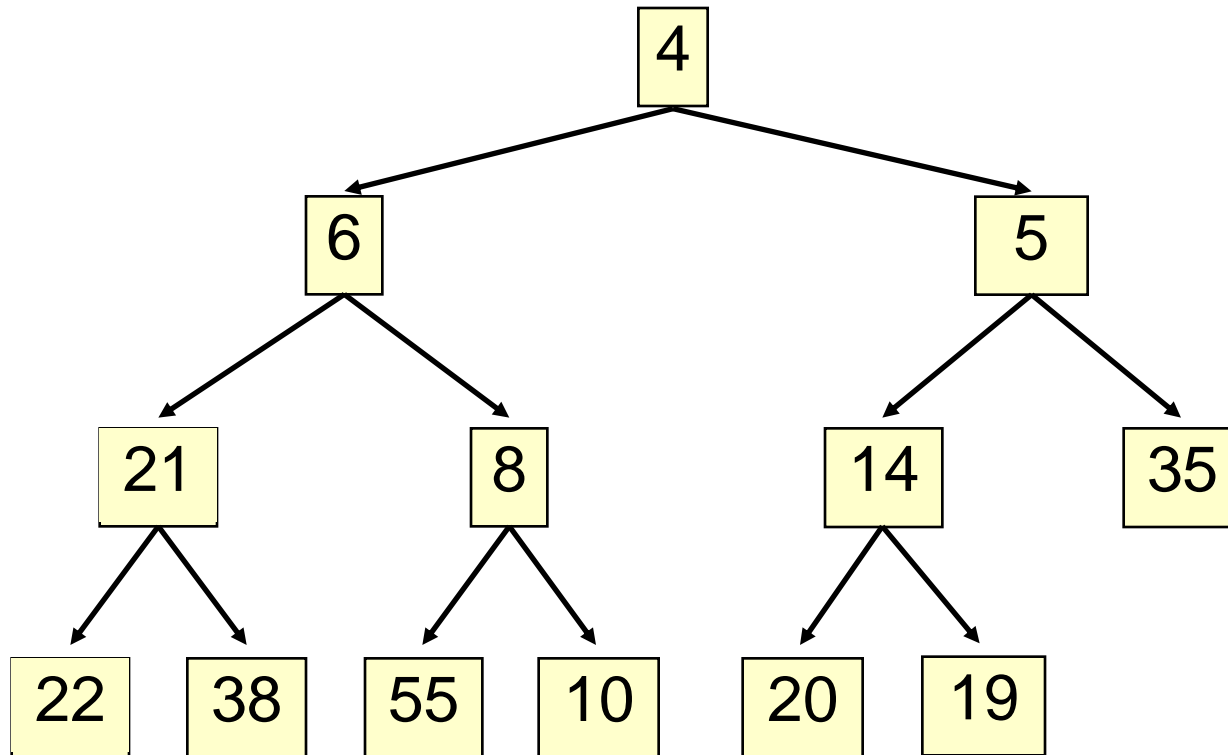
32

- Remove the least element – it is at the root
- This leaves a hole at the root – fill it in with the last element of the array
- If this violates heap order because the root element is too big, swap it down with the smaller of its children
- Continue swapping it down until it finds its rightful place
- The heap invariant is maintained!



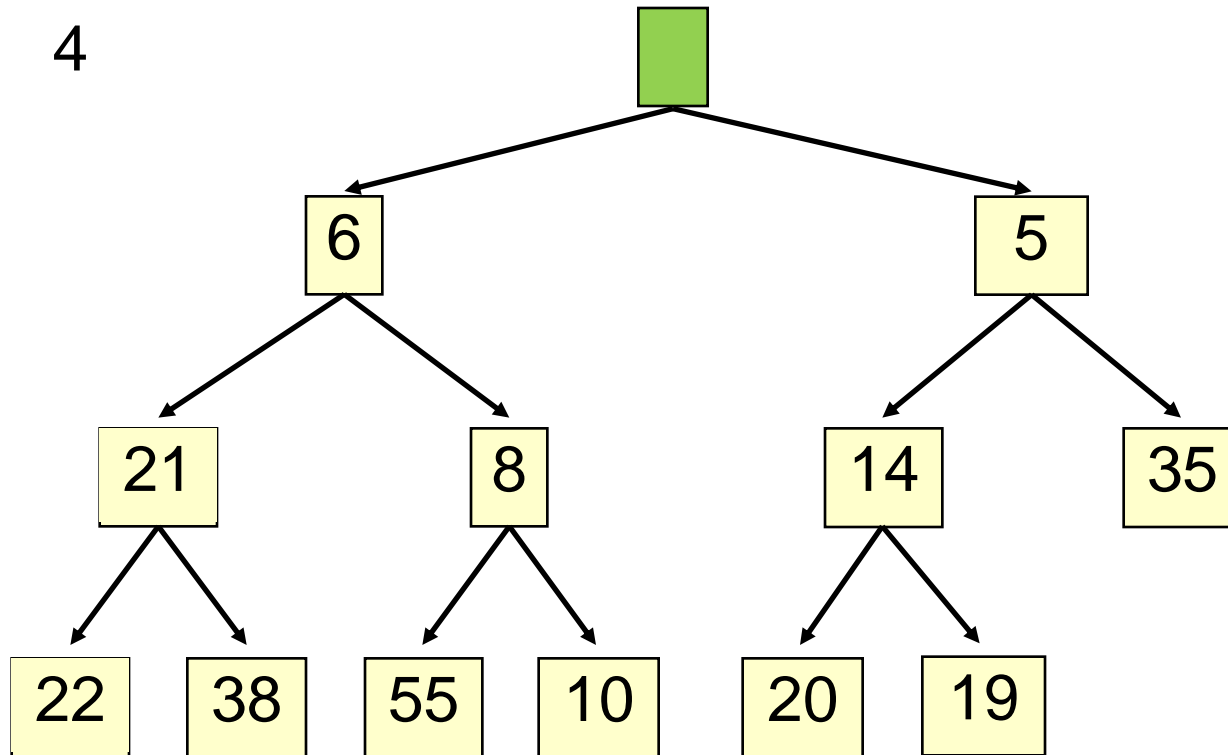
# extract ()

33



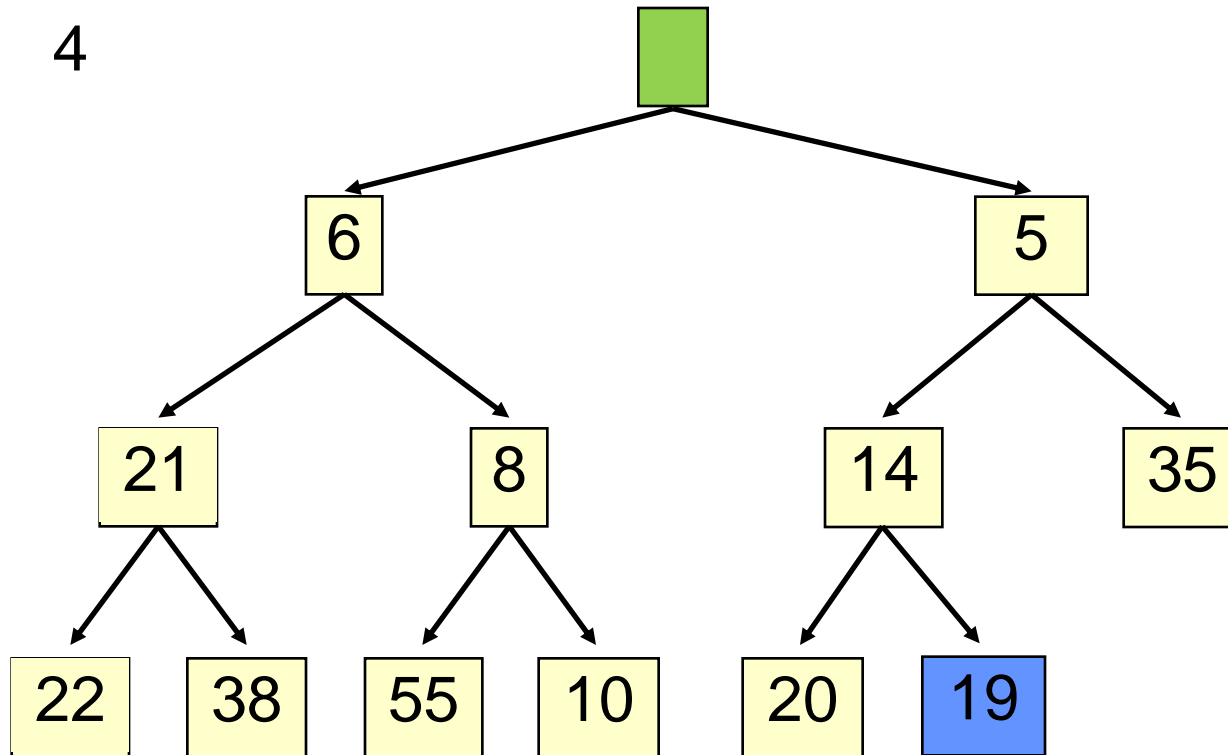
# extract ()

34



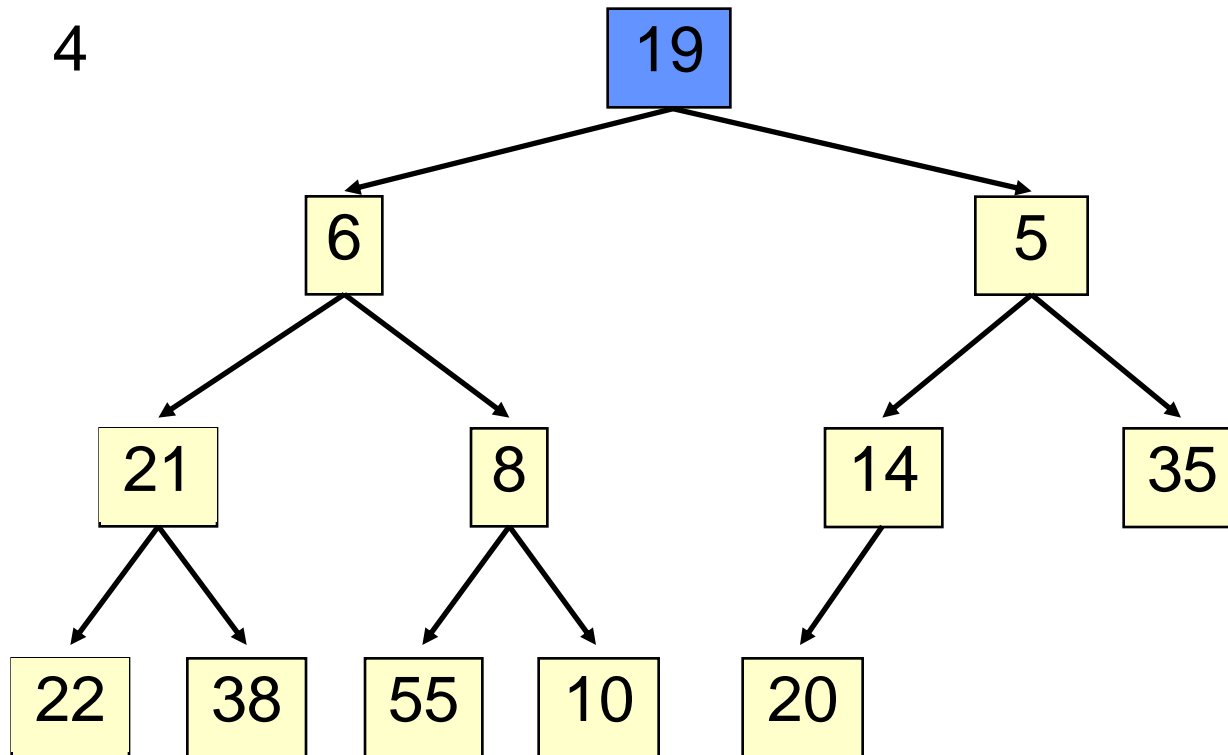
# extract()

35



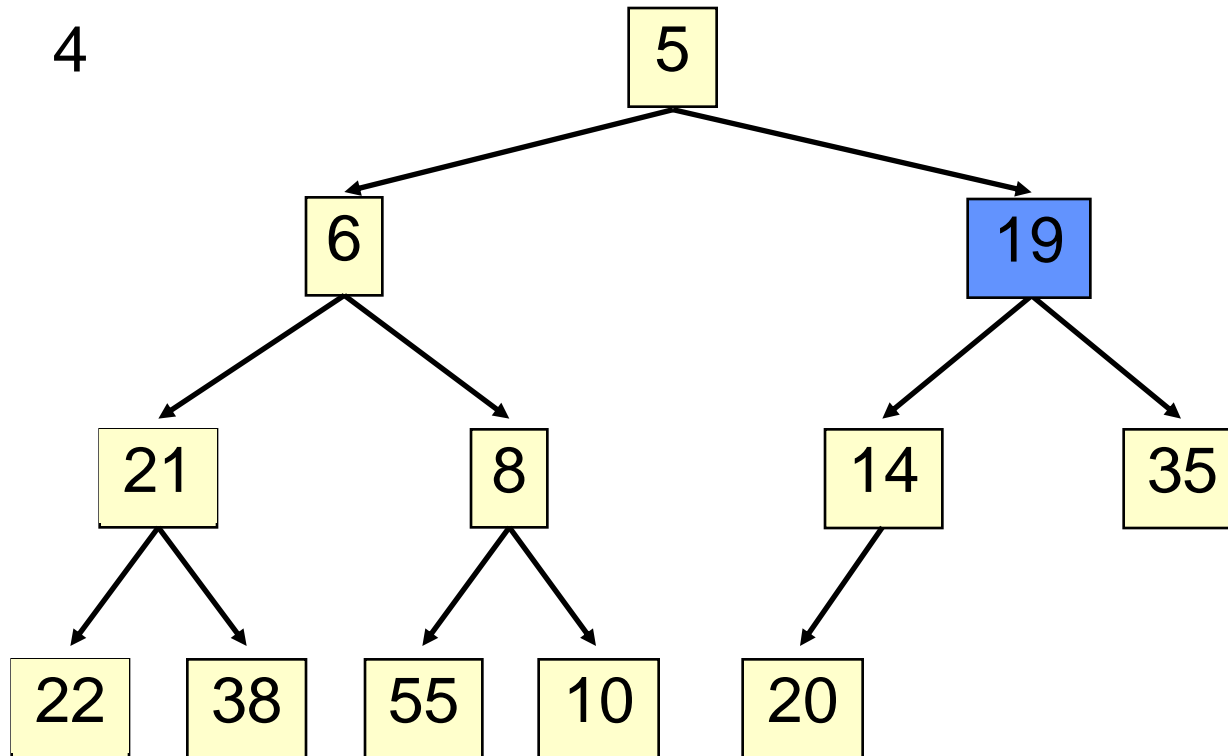
# extract()

36



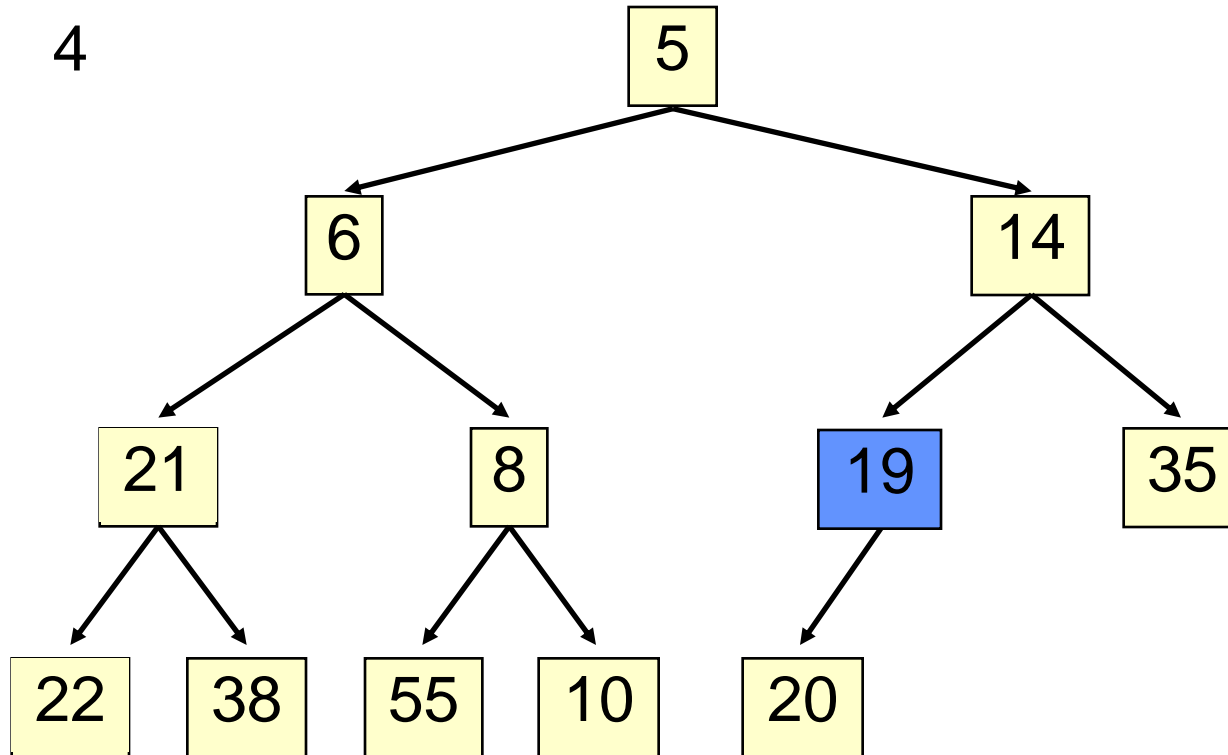
# extract()

37



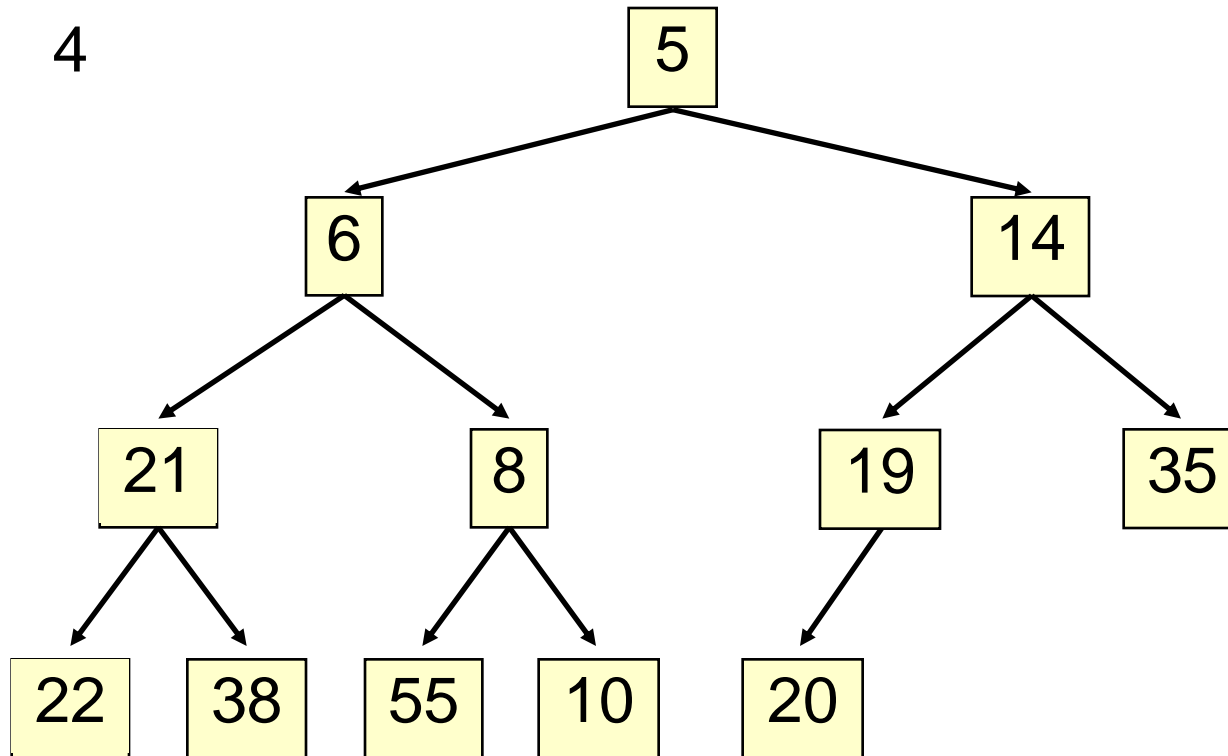
# extract ()

38



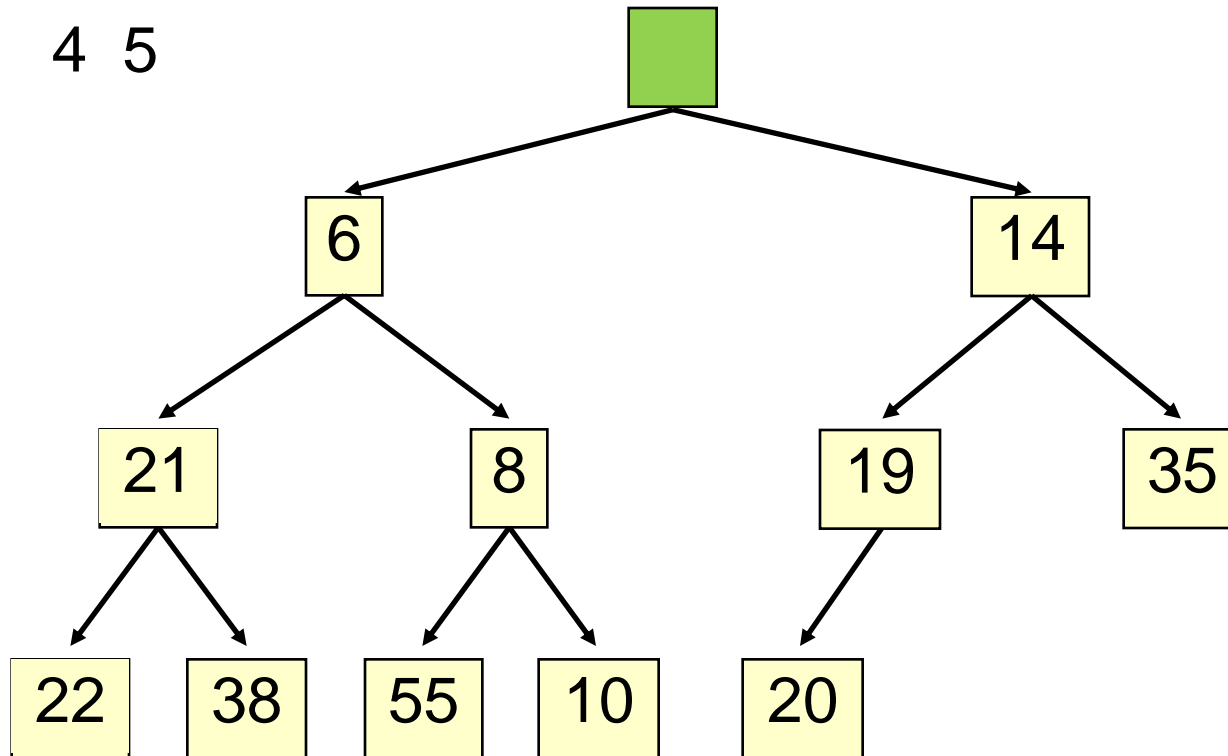
# extract ()

39



# extract ()

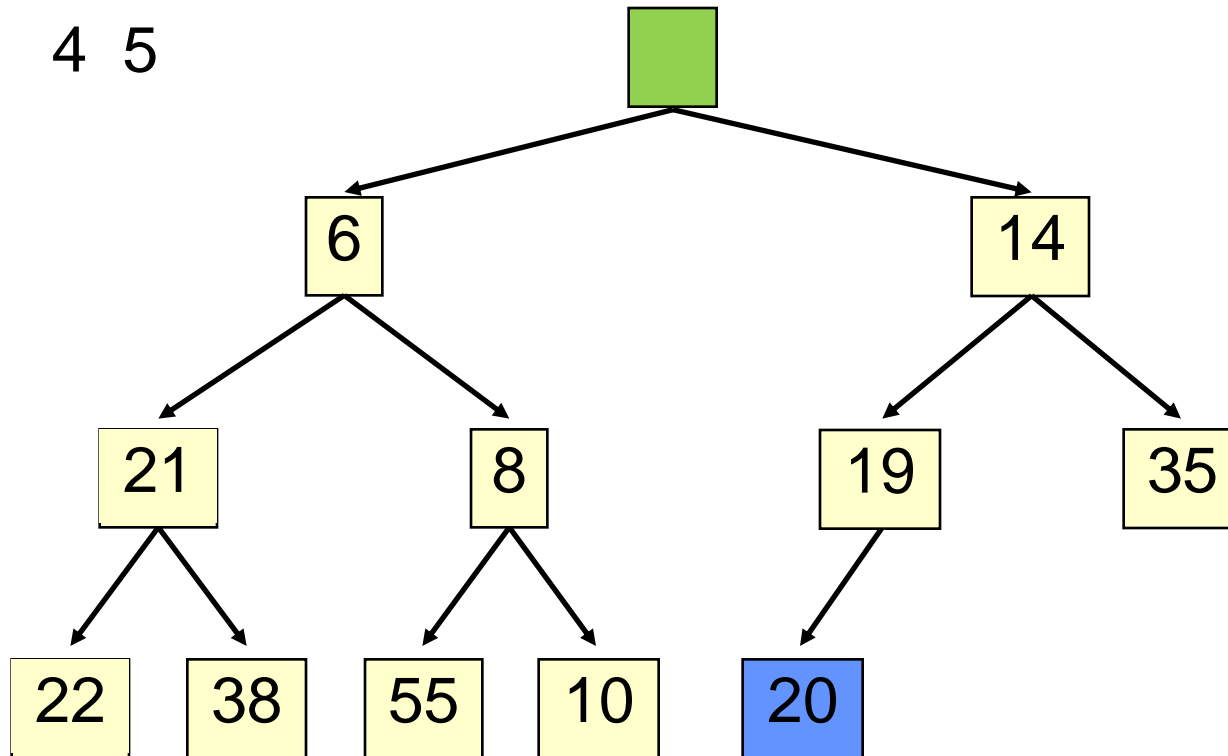
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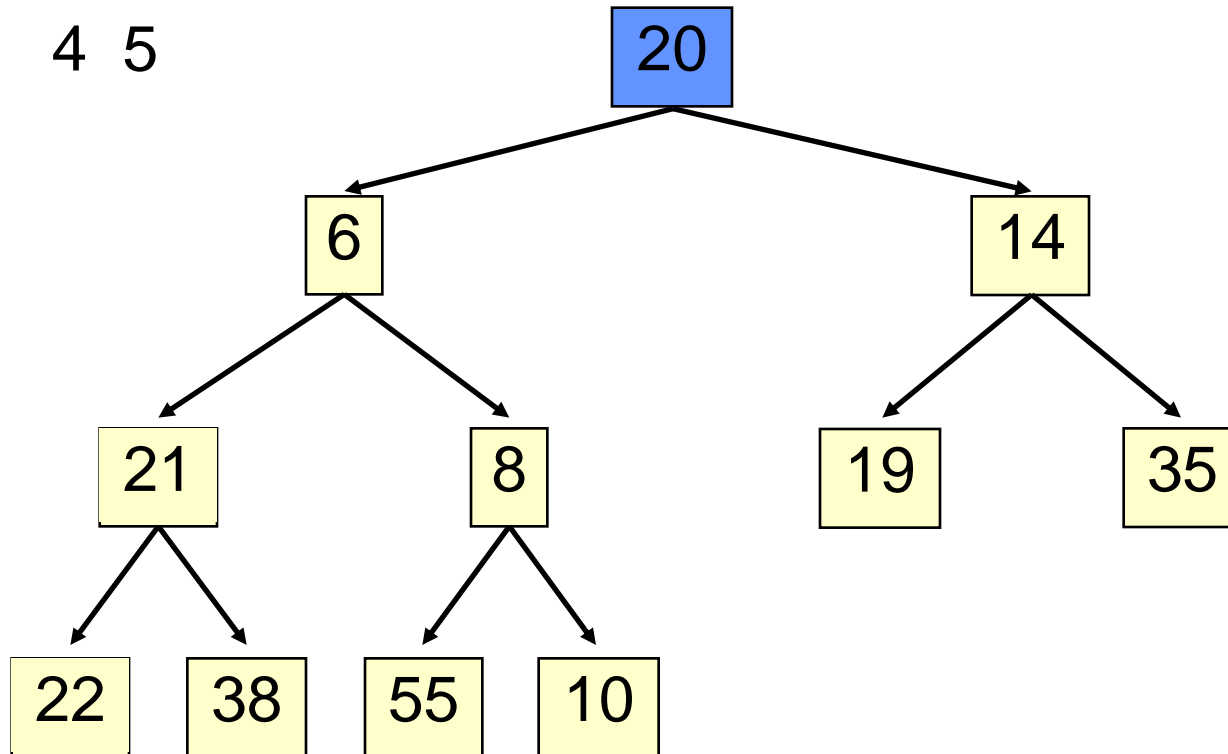
# extract()

41



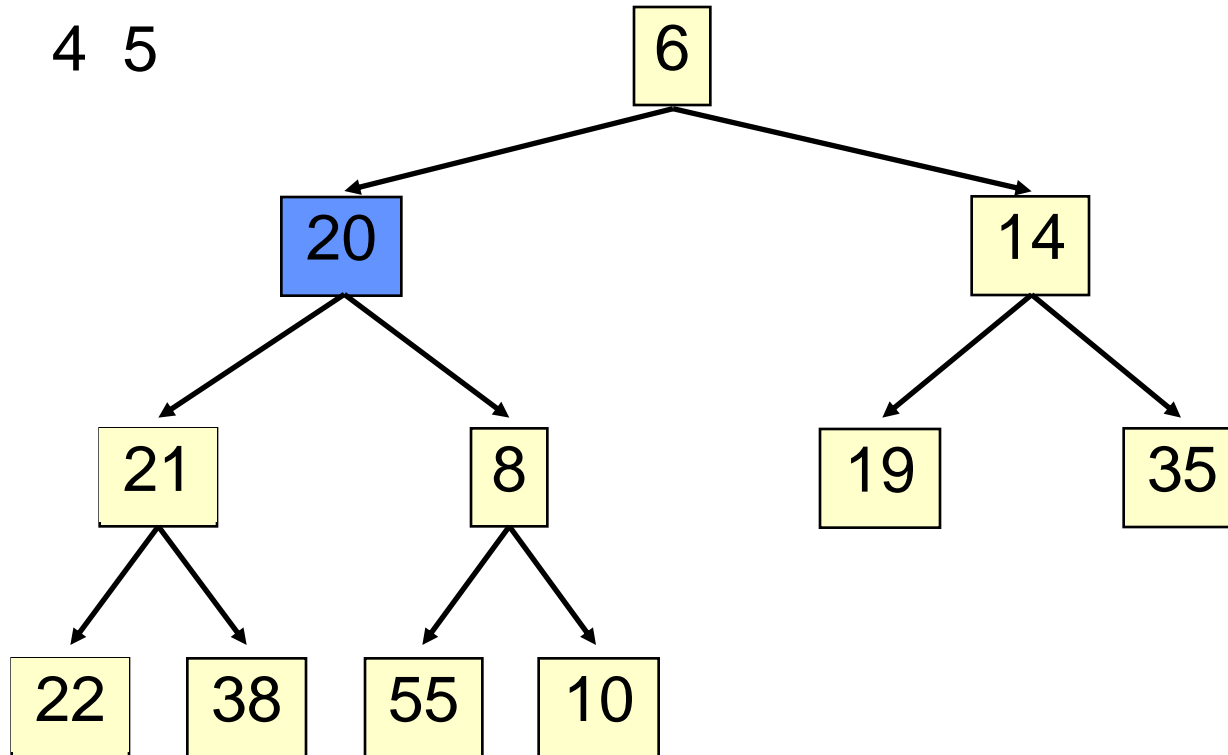
# extract()

42



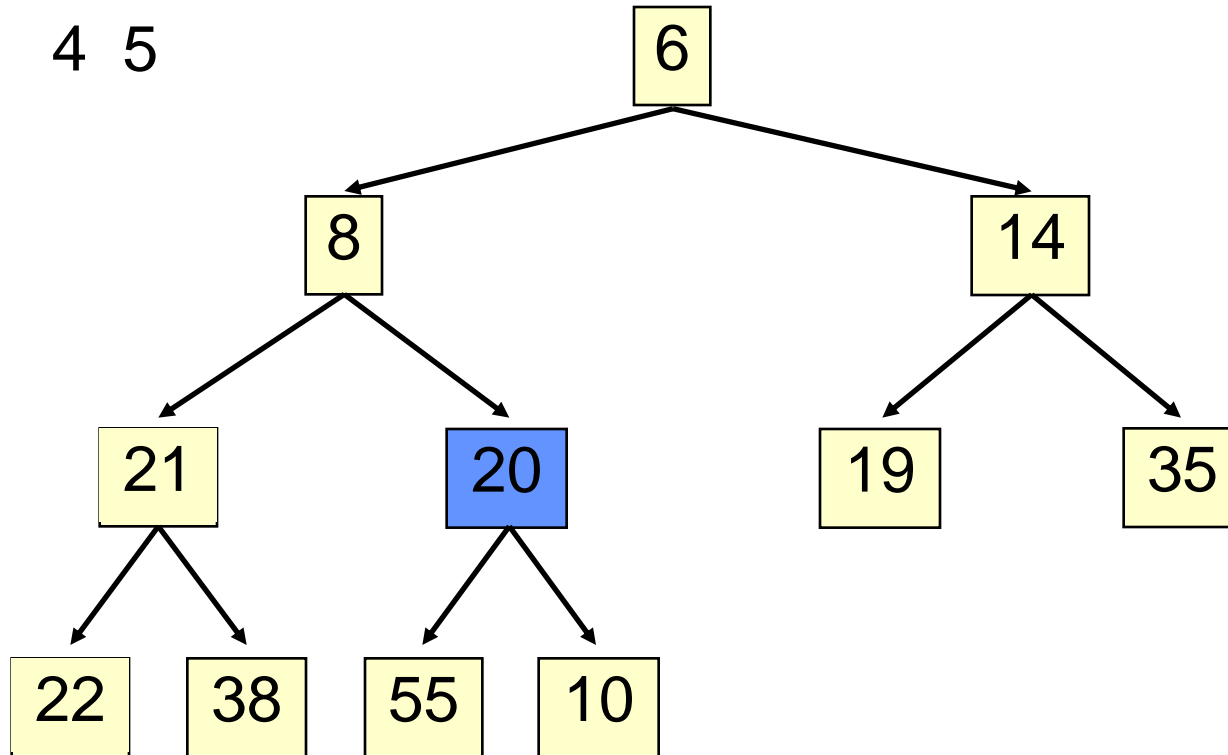
# extract ()

43



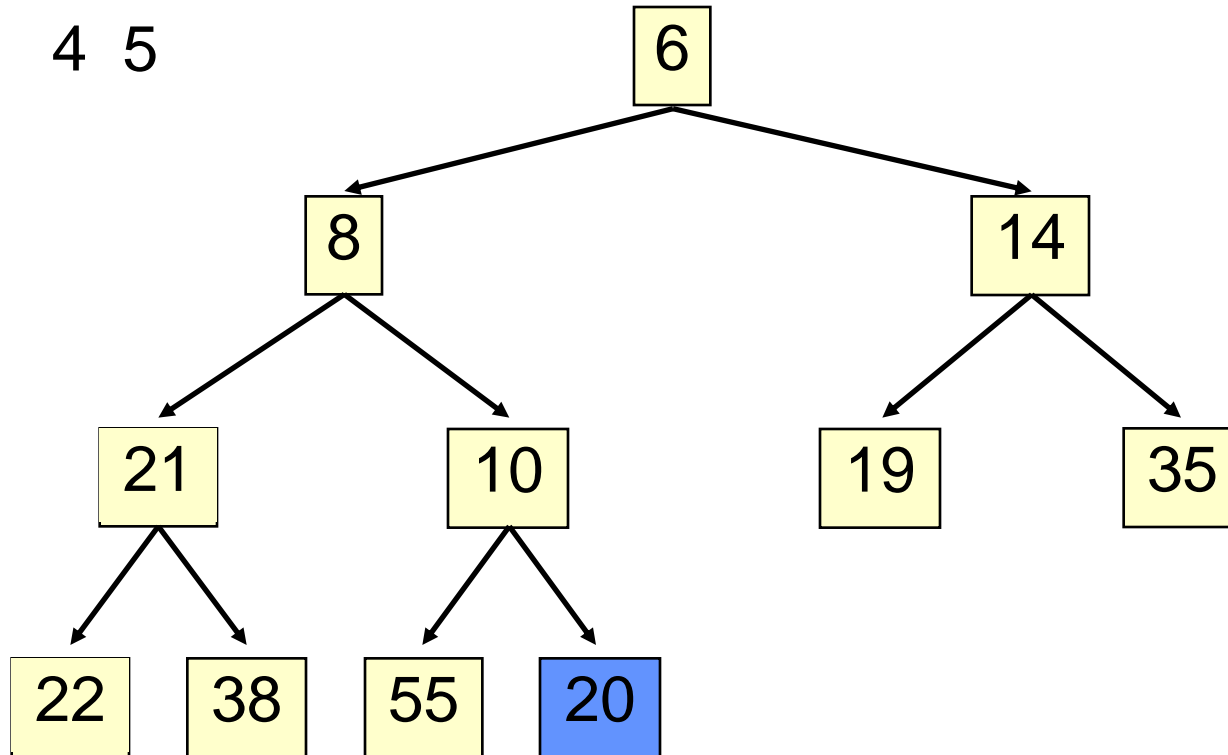
# extract()

44



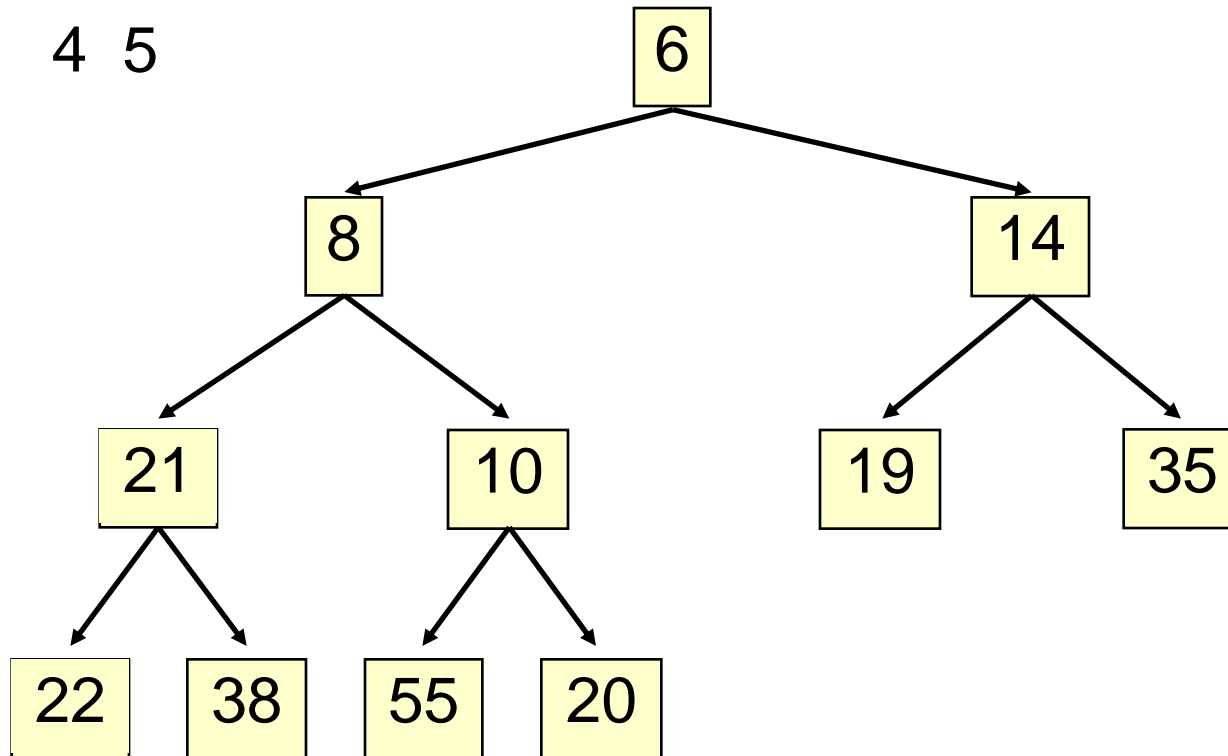
# extract ()

45



# extract ()

46



# extract ()

47

- Time is  $O(\log n)$ , since the tree is balanced

# extract()

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```
public E extract() {
    if (size() == 0) return null;
    E temp = elementAt(0);
    setElementAt(elementAt(size() - 1), 0);
    setSize(size() - 1);
    rotateDown(0);
    return temp;
}

private void rotateDown(int index) {
    int child = 2*(index + 1); //right child
    if (child >= size()
        || elementAt(child - 1).compareTo(elementAt(child)) < 0)
        child -= 1;
    if (child >= size()) return;
    if (elementAt(index).compareTo(elementAt(child)) <= 0)
        return;
    swap(index, child);
    rotateDown(child);
}
```



# HeapSort

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Given a `Comparable []` array of length  $n$ ,

- Put all  $n$  elements into a heap –  $O(n \log n)$
- Repeatedly get the min –  $O(n \log n)$

```
public static void heapSort(Comparable[] a) {  
    PriorityQueue<Comparable> pq = new PriorityQueue<Comparable>(a);  
    for (int i = 0; i < a.length; i++) { a[i] = pq.extract(); }  
}
```

# PQ Application: Simulation

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□ Example: Probabilistic model of bank-customer arrival times and transaction times, how many tellers are needed?

- Assume we have a way to generate random inter-arrival times
- Assume we have a way to generate transaction times
- Can simulate the bank to get some idea of how long customers must wait

## Time-Driven Simulation

- Check at each *tick* to see if any event occurs

## Event-Driven Simulation

- Advance clock to next event, skipping intervening *ticks*
- This uses a PQ!