L20. More on 2D Arrays

Operations Subscripting Functions* & 2D Arrays

*Will see two new things. Boolean-valued functions Functions that have a function as a parameter.

Two Applications

A commercial setting that involves cost arrays, inventory arrays, and purchase orders.

A setting that requires the visualization of a function of two variables f(x,y) via contour plotting.

A Cost/Inventory Setting

A company has 3 factories that make 5 different products.

The cost of making a product varies from factory to factory.

The inventory varies from factory to factory.

Problems

A customer submits a purchase order that is to be filled by a single factory.

- How much would it cost a factory to fill the order?
- 2. Does a factory have enough inventory to fill the order?
- 3. Among the factories that can fill the order, who can do it most cheaply?

Cost Array

	10	36	22	15	62
C:	12	35	20	12	66
	13	37	21	16	59

The value of C(i,j) is what it costs factory i to make product j.

Inventory Array

	38	5	99	34	42
Inv:	82	19	83	12	42
	51	29	21	56	87

The value of Inv(i,j) is the inventory in factory i of product j.

Purchase Order

The value of PO(j) is the number product j's that the customer wants

How Much Does it Cost for Each Factory to Process a Purchase order?

	-					
		10	36	22	15	62
	C:	12	35	20	12	66
		13	37	21	16	59
	PO:	1	0	12	29	5
or						

For factory 1:

1*10 + 0*36 + 12*22 + 29* 15 + 5*62

	10	36	22	15	62
C:	12	35	20	12	66
	13	37	21	16	59

s = 0;
For
factory 1:
 s = s + C(1,j) * PO(j)
end

	10	36	22	15	62
C:	12	35	20	12	66
	13	37	21	16	59

PO: 1 0 12 29 5

For factory 1:

s =	0;			
for	j=1:5			
S	= s +	C(1,j)	*	PO(j)
end				

	10	36	22	15	62
2:	12	35	20	12	66
	13	37	21	16	59

For factory 1:

s =	0;			
for	j=1:5			
S	= s +	C(1,j)	*	PO(j)
end				

	10	36	22	15	62
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For factory 1:

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S	= s +	C(1,j)	*	PO(j)
end				

	10	36	22	15	62
C:	12	35	20	12	66
	13	37	21	16	59

PO: 1 0 12 29 5

s = 0;
For
factory 1:
 s = s + C(1,j) * PO(j)
end

PO: 1 0 12 29 5
s = 0;
for
for j=1:5

$$s = s + C(2,j)*PO(j)$$

end

	s = 0;
For	for j=1:5
factory i:	s = s + C(i,j)*PO(j)
	end

Encapsulate...

function TheBill = iCost(i,C,PO)

- % The cost when factory i fills
- % the purchase order

```
nProd = length(PO)
```

```
TheBill = 0;
```

```
for j=1:nProd
```

TheBill = TheBill + C(i,j)*PO(j);

end

Finding the Cheapest

	10	36	22	15	62	1019
C:	12	35	20	12	66	930
	13	37	21	16	59	1040
P0:	1	0	12	29	5	

As computed by iCost

Finding Cheapest: Initialization

С:	10 12	36 35	22 20	15 12	62 66	1019 930	
	13	37	21	16	59	1040	Can we do
P0:	1	0	12	29	5		better?
i	Best	t:	0		mir	nBill:	inf

A Note on "inf"

A special value that can be regarded as + infinity.

- x = 10/0 assigns inf to x
- y = 1+x assigns inf to y
- z = 1/x assigns zero to z
- w < inf is always true if w is numeric

Improvement at i = 1



Improvement at i = 2



No Improvement at i = 3



Finding the Cheapest

```
iBest = 0; minBill = inf;
for i=1:nFact
    iBill = iCost(i,C,PO);
    if iBill < minBill
% Found an Improvement
        iBest = i; minBill = iBill;
    end
end
```

Inventory Considerations

What if a factory lacks the inventory to fill the purchase order?

Such a factory should be excluded from the find-the-cheapest computation.

Who Can Fill the Order?

	38	5	99	34	42	Yes
Inv:	82	19	83	12	42	No
	51	29	21	56	87	Yes
PO:	1	0	12	29	5	

Because 12 < 29

Wanted: A True/False Function



B is "true" if factory i can fill the order. B is "false" if factory i cannot fill the order.

Boolean Operations in Matlab

SO FAR we have indicated that expressions like

 $a \leq x \& \& x \leq b$

abs(y) > 10

are either TRUE or FALSE.

The 0-1 Secret

In reality, expressions like

a <= x && x <= b

abs(y) > 10

render the value "1" if they are TRUE and "0" if they are FALSE.

Example

- >> x = 8; y = 7;
- >> B = x<y
- в =
 - 0
- >> B = x>y

1

B =

A Boolean-Valued Function

function	в =	Overlap	(a,b)	,c,d)
----------	-----	---------	-------	-------

- % B is true if intervals [a,b]
- % and [c,d] intersect.
- % Otherwise B is false.
- % Assume a<b and c<d.

abToLeft = b < c;</pre>

abToRight = d < a;

B = ~(abToLeft || abToRight);

Using Overlap

```
S = 0;
for k=1:100
   a = rand; b = a + rand;
   c = rand; d = c + rand;
   if Overlap(a,b,c,d)
     s = s+1;
   end
end
probOverlap = s/100
```

Back to Inventory Problem

	38	5	99	34	42
Inv:	82	19	83	12	42
	51	29	21	56	87

PO: 1	0	12	29	5
-------	---	----	----	---

Initialization

	38	5	99	34	42
Inv:	82	19	83	12	42
	51	29	21	56	87

B: 1

Still True...

	38	5	99	34	42
Inv:	82	19	83	12	42
	51	29	21	56	87

B: 1

PO: 1 0 12 29 5

B = B && (Inv(2,1) >= PO(1))

Still True...

	38	5	99	34	42
Inv:	82	19	83	12	42
	51	29	21	56	87

B: 1

B = B && (Inv(2,2) >= PO(2))

Still True...

	38	5	99	34	42
Inv:	82	19	83	12	42
	51	29	21	56	87

B: 1

PO: 1 0 12 29 5

B = B && (Inv(2,3) >= PO(3))

No Longer True...

 \mathbf{O}

	38	5	99	34	42	
Inv:	82	19	83	12	42	В:
	51	29	21	56	87	

B = B && (Inv(2,4) >= PO(4))

Encapsulate...

```
function B = iCanDo(i, Inv, PO)
% B is true if factory i can fill
% the purchase order. Otherwise, false
  nProd = length(PO);
  j = 1;
 B = 1;
 while j<=nProd && B
     B = B \&\& (Inv(i,j) >= PO(j));
     j = j+1;
  end
```

Back To Finding the Cheapest

```
iBest = 0; minBill = inf;
 for i=1:nFact
    iBill = iCost(i,C,PO);
    if iBill < minBill
%
       Found an Improvement
       iBest = i; minBill = iBill;
    end
end
```

Don't bother with this unless sufficient inventory.

Back To Finding the Cheapest

```
iBest = 0; minBill = inf;
 for i=1:nFact
    if iCanDo(i,Inv,PO)
       iBill = iCost(i,C,PO);
       if iBill < minBill
%
          Found an Improvement
          iBest = i; minBill = iBill;
       end
    end
```

end

```
function [iBest,minBill] = ...
                   Cheapest(C, Inv, PO)
[nFact,nProd] = size(C);
iBest = 0; minBill = inf;
for i=1:nFact
   if iCanDo(i,Inv,PO)
      iBill = iCost(i,C,PO);
      if iBill < minBill
         iBest = i; minBill = iBill;
      end
   end
end
```

Finding the Cheapest



New Problem

Visualizing a function of the form z = f(x,y).

Think of z as an elevation which depends on the coordinates x and y of the location.

Sample Elevation Function

function z = Elev(x,y)

- r1 = $(x-1)^2$ + $3*(y-1.5)^2$; r2 = $2*(x+2)^2$ + $(y-.5)^2$; r3 = $(x-.5)^2$ + $7*y^2$;
- z = 100*exp(-.5*r1) + ... 90*exp(-.3*r2) + ... 80*exp(-.4*r3);

Three Hills at (1,1.5),(-2,.5), (.5,0)

Its Contour Plot



Making a Contour Plot

```
x = linspace(-5, 4, 200);
```

```
y = linspace(-2.5, 6.5, 200);
```

```
A = zeros(200, 200);
```

```
for i=1:200
```

```
for j=1:200
```

```
A(i,j) = Elev(x(j),y(i));
```

end

end

contour(x,y,A,15)

Set up a matrix of function evals

General Set-Up

```
function A = SetUp(f, xVals, yVals)
Nx = length(xVals);
Ny = length(yVals);
A = zeros(Ny,Nx);
for i=1:Ny
  for j=1:Nx
     A(i,j) = f(xVals(j),yVals(i));
  end
end
```

Calling SetUp

- x = linspace(-5, 4, 200);
- y = linspace(-2.5, 6.5, 200);
- F = SetUp(@Elev,x,y);

Not just 'Elev' The @ is required for function parameters.

Generating a Cross Section



Enter endpoints via ginput Sample Elev(x,y) along the line segment

Mouse Input Via ginput

To draw a line segment connecting (a(1),b(1)) and (a(2),b(2)):

[a,b] = ginput(2);
plot(a,b)

[a,b] = ginput(n) puts the mouseclick coords in length-n arrays a and b.

- n = 100;
- t = linspace(0,1,n);
- x = linspace(a(1),a(2),n);
- y = linspace(b(1),b(2),n);

for i=1:n

- % At "time" t(i) we are at (x(i),y(i)).
- % Compute elevation at time t(i).

```
f(i) = Elev(x(i),y(i));
```

end

figure

plot(t,f)