Lecture 22

## Designing Sequence Algorithms

## Announcements for This Lecture

## Assignments

## Prelim 2

- A5 graded by weekend
- We just starting on it
- Should be working on A6
- Due week from Today
- Work on a method a day
- Should start stenography no later than Sunday
- Friday extension?
- A7 due after class ends
- High scores again
- Mean: 83, Median: 86
- 150/404 scored 90+
- Historical mean: 76
- For-loop, not recursion hard
- But good grade distribution
- A: 90+
- B: Mid-low 70s to high 80s
- C: 50 to mid-low 70s


## Horizontal Notation for Sequences



Example of an assertion about an sequence b. It asserts that:

1. $\mathrm{b}[0 . . \mathrm{k}-1]$ is sorted (i.e. its values are in ascending order)
2. Everything in $\mathrm{b}[0 . . \mathrm{k}-1]$ is $\leq$ everything in $\mathrm{b}[\mathrm{k} . . \operatorname{len}(\mathrm{b})-1]$


Given index $h$ of the first element of a segment and

$(\mathrm{h}+1)-\mathrm{h}=1$

## Developing Algorithms on Sequences

- Specify the algorithm by giving its precondition and postcondition as pictures.
- Draw the invariant by drawing another picture that "generalizes" the precondition and postcondition
- The invariant is true at the beginning and at the end
- The four loop design questions (memorize them)

1. How does loop start (how to make the invariant true)?
2. How does it stop (is the postcondition true)?
3. How does repetend make progress toward termination?
4. How does repetend keep the invariant true?

## Generalizing Pre- and Postconditions

- Dutch national flag: tri-color
- Sequence of 0..n-1 of red, white, blue "pixels"
- Arrange to put reds first, then whites, then blues



## Generalizing Pre- and Postconditions

- Finding the minimum of a sequence.

- Put negative values before nonnegative ones.

| 0 n |  |  |  | n |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| pre: b | ? | ? |  | and $\mathrm{n}>=0$ | (values in $0 . . \mathrm{n}$ are unknown) |
| 0 |  | k | n |  |  |
| post: b |  |  | $>=0$ |  |  |
| 0 |  |  | n | pre: $\mathrm{k}=0$, |  |
| inv: b | $<0$ | $?$ | $>=0$ | pre: $\mathrm{k}=\mathrm{n}$ | (values in k..j-1 are unknown) |

## Partition Algorithm

- Given a sequence $\mathrm{b}[\mathrm{h} . \mathrm{k}]$ with some value x in $\mathrm{b}[\mathrm{h}]$ :

- Swap elements of $\mathrm{b}[\mathrm{h} . \mathrm{k}]$ and store in j to truthify post:

change:
into
or

- x is called the pivot value
- x is not a program variable
- denotes value initially in b[h]


## Partition Algorithm

- Given a sequence $\mathrm{b}[\mathrm{h} . \mathrm{k}]$ with some value x in $\mathrm{b}[\mathrm{h}]$ :

- Swap elements of $b[h . . k]$ and store in $j$ to truthify post:

|  | h | i | i+1 | k |
| :---: | :---: | :---: | :---: | :---: |
| post: b | <= x | x | >= x |  |



- Agrees with precondition when $\mathrm{i}=\mathrm{h}, \mathrm{j}=\mathrm{k}+1$
- Agrees with postcondition when $\mathrm{j}=\mathrm{i}+1$


## Partition Algorithm Implementation

def partition(b, h, k):
"""Partition list b[h..k] around a pivot $\mathrm{x}=\mathrm{b}[\mathrm{h}]$ """
$\mathrm{i}=\mathrm{h} ; \mathrm{j}=\mathrm{k}+\mathrm{l} ; \mathrm{x}=\mathrm{b}[\mathrm{h}]$
\# invariant: $\mathrm{b}[\mathrm{h} . \mathrm{i}-\mathrm{l}]<\mathrm{x}, \mathrm{b}[\mathrm{i}]=\mathrm{x}, \mathrm{b}[\mathrm{j} . \mathrm{k}]>=\mathrm{x}$
while $\mathrm{i}<\mathrm{j}-\mathrm{l}$ :
if $b[i+1]>=x$ :
\# Move to end of block.
_swap(b,i+l,j-l)
$j=j-1$
else: \# b[i+1] < x
_swap(b,i,i+l)
$\mathrm{i}=\mathrm{i}+\mathrm{l}$
\# post: b[h.i-l $]<\mathrm{x}, \mathrm{b}[\mathrm{i}]$ is x , and $\mathrm{b}[\mathrm{i}+\mathrm{l} . \mathrm{k}]>=\mathrm{x}$
return i

## Partition Algorithm Implementation

def partition(b, h, k):
"""Partition list b[h..k] around a pivot $x=b[h]$ """
$\mathrm{i}=\mathrm{h} ; \mathrm{j}=\mathrm{k}+\mathrm{l} ; \mathrm{x}=\mathrm{b}[\mathrm{h}]$
\# invariant: $\mathrm{b}[\mathrm{h} . . \mathrm{i}-1 \mathrm{l}]<\mathrm{x}, \mathrm{b}[\mathrm{i}]=\mathrm{x}, \mathrm{b}[\mathrm{j} . \mathrm{k}]>=\mathrm{x}$
while $\mathrm{i}<\mathrm{j}-\mathrm{l}$ :
if $b[i+1]>=x$ :
\# Move to end of block.
_swap(b,i+l,j-1)
$j=j-1$
else: \#b[i+1] < x
_swap(b,i,i+1)

$$
\mathrm{i}=\mathrm{i}+\mathrm{l}
$$

\# post: $\mathrm{b}[\mathrm{h} . \mathrm{i} \mathrm{i}-\mathrm{l}]<\mathrm{x}, \mathrm{b}[\mathrm{i}]$ is x , and $\mathrm{b}[\mathrm{i}+\mathrm{l} . . \mathrm{k}]>=\mathrm{x}$
return i


| h | i |  |  |  | $\mathrm{i}+1$ | j |  | k |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 1 | 2 | 1 | 3 | 5 | 0 | 6 | 3 | 8 |


| h | i |  |  |  | j |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| l |  |  |  |  |  |  |  |  |
| 1 | 2 | 1 | 3 | 0 | 5 | 6 | 3 | 8 |



## Dutch National Flag Variant

- Sequence of integer values
- 'red' = negatives, 'white' $=0$, 'blues' = positive
- Only rearrange part of the list, not all



## Dutch National Flag Algorithm

$\operatorname{def} \operatorname{dnf}(\mathrm{b}, \mathrm{h}, \mathrm{k})$ :
"""Returns: partition points as a tuple (i,j)"""
$\mathrm{t}=\mathrm{h} ; \mathrm{i}=\mathrm{k}+\mathrm{l}, \mathrm{j}=\mathrm{k}$;
\# inv: b[h.t-l] < 0, b[t..i-l] ?, b[i..j] = $0, b[j+1 . . k]>0$
while t < i :
if $b[i-1]<0$ :
$\operatorname{swap}(\mathrm{b}, \mathrm{i}-1, \mathrm{t})$


$$
t=t+l
$$

elif $b[i-1]==0$ :

$$
\mathrm{i}=\mathrm{i}-1
$$

else:


$$
\begin{aligned}
& \operatorname{swap}(b, i-1, j) \\
& i=i-1 ; j=j-1
\end{aligned}
$$

\# post: b[h.i-i-l] < $0, b[i . . j]=0, b[j+1 . . k]>0$ return (i, j)

## Linear Search

- Vague: Find first occurrence of v in $\mathrm{b}[\mathrm{h} . \mathrm{k}-1]$.
- Better: Store an integer in i to truthify result condition post: post: $1 . \mathrm{v}$ is not in $\mathrm{b}[\mathrm{h} . \mathrm{i}-1]$

2. $\mathrm{i}=\mathrm{k} \quad$ OR $\mathrm{v}=\mathrm{b}[\mathrm{i}]$


## Linear Search



## Linear Search

def linear_search(b,c,h):
"""Returns: first occurrence of c in b[h..]"""
\# Store in $i$ the index of the first c in $\mathrm{b}[\mathrm{h}$. .]
$\mathrm{i}=\mathrm{h}$
\# invariant: c is not in b[0..i-1]
while $\mathrm{i}<\operatorname{len}(\mathrm{b})$ and $\mathrm{b}[\mathrm{i}]$ != c :

$$
\mathrm{i}=\mathrm{i}+\mathrm{l}
$$

\# post: b[i] == c and c is not in b[h.i-l]
return i if i < len(b) else -l

## Analyzing the Loop

1. Does the initialization make inv true?
2. Is post true when inv is true and condition is false?
3. Does the repetend make progress?
4. Does the repetend keep inv true?


## Binary Search

- Vague: Look for v in sorted sequence segment b[h..k].
- Better:
- Precondition: b[h.k-1] is sorted (in ascending order).
- Postcondition: b[h.i] <= v and $\mathrm{v}<\mathrm{b}[\mathrm{i}+1 . \mathrm{k}-1]$
- Below, the array is in non-descending order:


Called binary search because each iteration of the loop cuts the array segment still to be processed in half

## Extras Not Covered in Class

## Loaded Dice

- Sequence $p$ of length $n$ represents $n$-sided die
- Contents of $p$ sum to 1
- $\mathrm{p}[\mathrm{k}]$ is probability die rolls the number k

weighted d6, favoring 5, 6
- Goal: Want to "roll the die"
- Generate random number $r$ between 0 and 1
- Pick p[i] such that $\mathrm{p}[\mathrm{i}-1]<\mathrm{r} \leq \mathrm{p}[\mathrm{i}]$

| 0.1 | 0.1 | 0.1 | 0.1 | 0.3 | 0.3 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0.1 | 0.2 | 0.3 | 0.4 | 0.7 | 1.0 |

## Loaded Dice

- Want: Value i such that $\mathrm{p}[\mathrm{i}-1]<\mathrm{r}<=\mathrm{p}[\mathrm{i}]$

- Same as precondition if $\mathrm{i}=0$
- Postcondition is invariant + false loop condition


## Loaded Dice

def $\operatorname{roll}(p)$ :
"""Returns: randint in 0..len(p)-1; i returned with prob. p[i] Precondition: p list of positive floats that sum to 1. """
$\mathrm{r}=$ random.random() \#rin [0,1)
\# Think of interval [ 0,1 ] divided into segments of size p[i]
\# Store into it the segment number in which r falls.
$\mathrm{i}=0 ; \quad$ sum_of $=\mathrm{p}[0]$
\# inv: $\mathrm{r}>=$ sum of $\mathrm{p}[0]$.. $p[i-1]$; $p E n d=$ sum of $p[0]$.. $p[i]$
while $\mathrm{r}>=$ sum_of:
sum_of = sum_of + p[i+1]
$\mathrm{i}=\mathrm{i}+\mathrm{l}$
\# post: sum of p[0] .. p[i-1] <= $\mathrm{p}<$ sum of $\mathrm{p}[0]$.. $p[i]$ return i


## Reversing a Sequence





