

Machine Learning Theory (CS 6783)

Lecture 6 : VC dimension recap and continued

1 Recap

1. For any set $V \subset \mathbb{R}^n$:

$$\frac{1}{n} \mathbb{E}_{\epsilon} \left[\sup_{\mathbf{v} \in V} \sum_{t=1}^n \epsilon_t \mathbf{v}[t] \right] \leq \frac{1}{n} \sqrt{2 \left(\sup_{\mathbf{v} \in V} \sum_{t=1}^n \mathbf{v}^2[t] \right) \log |V|}$$

2. We defined growth function as $\Pi(\mathcal{F}, n) = \max_{x_1, \dots, x_n} |\mathcal{F}|_{x_1, \dots, x_n}|$
3. VC dimension : size of largest set of input instances we can shatter

$$\text{VC}(\mathcal{F}) = \max\{d : \Pi(\mathcal{F}, d) = 2^d\}$$

4. VC/Sauer/Shelah Lemma : $\Pi(\mathcal{F}, n) \leq \sum_{i=0}^{\text{VC}(\mathcal{F})} \binom{n}{i}$
5. \mathcal{F} is learnable if and only if $\text{VC}(\mathcal{F}) < \infty$

$$(a) \quad \mathcal{V}_n^{\text{stat}}(\mathcal{F}) \leq \sqrt{\frac{\text{VC}(\mathcal{F}) \log n}{n}}$$

$$(b) \quad \text{VC}(\mathcal{F}) = \infty \quad \implies \quad \mathcal{V}_n^{\text{stat}}(\mathcal{F}) \geq 1/4$$

2 VC dimension examples and utility lemmas

1. To show $\text{VC}(\mathcal{F}) \geq d$ show that you can at least pick d points x_1, \dots, x_d that can be shattered.
2. To show that $\text{VC}(\mathcal{F}) \leq d$ show that no configuration of $d + 1$ points can be shattered.

Examples : Intervals, axis aligned rectangle, lines, convex polygons

Claim 1. *VC dimension of half-spaces in \mathbb{R}^d is $d + 1$*

Proof. We consider half-spaces that map vector in \mathbb{R}^d to $\{\pm 1\}$. That is

$$\mathcal{F} = \{\mathbf{x} \mapsto \text{sign}(\mathbf{f}^\top \mathbf{x} + f_0) : \mathbf{f} \in \mathbb{R}^d, f_0 \in \mathbb{R}\}$$

We prove the statement as follows.

1. $\text{VC}(\mathcal{F}) \geq d + 1$:

We can shatter the points $\mathbf{e}_1, \dots, \mathbf{e}_d, \mathbf{0}$. To see this, note that given any $y_1, \dots, y_{d+1} \in \{\pm 1\}^{d+1}$, if we consider $f \in \mathcal{F}$ given by $f_0 = y_{d+1}$ and for all $i \in [d]$, $\mathbf{f}[i] = y_i - y_{d+1}$. Hence note that, $f(\mathbf{0}) = \text{sign}(\mathbf{f}^\top \mathbf{0} + f_0) = \text{sign}(y_{d+1}) = y_{d+1}$. Also, for any $i \in [d]$, $f(\mathbf{e}_i) = \text{sign}(\mathbf{f}^\top \mathbf{e}_i + f_0) = \text{sign}(y_i - y_{d+1} + y_{d+1}) = y_i$.

2. $\text{VC}(\mathcal{F}) < d + 2$:

By Radon theorem, any set of $d + 2$ points in \mathbb{R}^d can be partitioned into two disjoint subsets whose convex hulls have a non-empty intersection. Label one of these partitions +1 and other -1. No half-space can successfully label points in the intersection.

□

Claim 2. For any binary hypothesis class \mathcal{F} ,

$$\text{VC}(\mathcal{F}) \leq \log_2 |\mathcal{F}|$$

Proof. Note that for any d , $\Pi(\mathcal{F}, d) \leq |\mathcal{F}|$. From the definition of VC dimension, we have, $\text{VC}(\mathcal{F}) = \max\{d : \Pi(\mathcal{F}, d) = 2^d\}$. Hence $2^{\text{VC}(\mathcal{F})} \leq |\mathcal{F}|$ □

Claim 3. Consider any fixed function $g : \{\pm 1\}^k \mapsto \{\pm 1\}$. For every $i \in [k]$, let \mathcal{F}_i be a some class of binary hypotheses mapping from input space \mathcal{X} . Let $\mathcal{G} = \{x \mapsto g(f_1(x), \dots, f_k(x)) : f_i \in \mathcal{F}_i\}$. Then

$$\text{VC}(\mathcal{G}) \leq 2k \text{VC}(\mathcal{F}) \log(5k)$$

Proof.

$$\Pi(\mathcal{G}, n) \leq \prod_{i=1}^k \Pi(\mathcal{F}_i, n) \leq (en)^{\sum_{i=1}^k \text{VC}(\mathcal{F}_i)}$$

By definition of VC dimension, we have that $\text{VC}(\mathcal{G}) = \max\{d : \Pi(\mathcal{G}, d) = 2^d\}$. Hence,

$$2^{\text{VC}(\mathcal{G})} \leq (e \text{VC}(\mathcal{G}))^{\sum_{i=1}^k \text{VC}(\mathcal{F}_i)}$$

Hence

$$\text{VC}(\mathcal{G}) \leq \log_2 (e \text{VC}(\mathcal{G})) \sum_{i=1}^k \text{VC}(\mathcal{F}_i)$$

Hence, we can conclude that $\text{VC}(\mathcal{G}) \leq \sum_{i=1}^k \text{VC}(\mathcal{F}_i) \log \left(\sum_{i=1}^k \text{VC}(\mathcal{F}_i) \cdot \log \left(\sum_{i=1}^k \text{VC}(\mathcal{F}_i) \right) \right)$ □