

Computational Requirements of Scientific Applications

Computational Science Applications

Simulation of physical phenomena

- fluid flow over aircraft (Boeing 777 designed by simulation)
- fatigue fracture in aircraft bodies
- bone growth
- evolution of galaxies

Two main approaches

- **continuous methods**: fields and partial differential equations (pde's) (eg. Navier-Stokes equations, Maxwell's equations, elasticity equations..)
- **discrete methods**: particles and forces between them (eg. Gravitational/Coulomb forces)

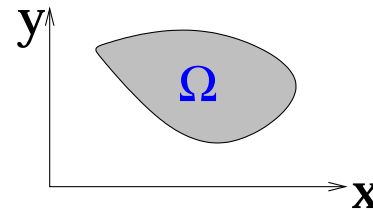
We will focus on pde's in this lecture.

Modeling physical phenomena using pde's

PDE : $L u = f$

eg: $\left(\frac{\delta^2}{\delta x^2} + \frac{\delta^2}{\delta y^2} \right) u = 0$

Domain: Ω



Boundary conditions: on $\delta\Omega$

$u(x,y) = x + y \quad | \quad (x,y) \text{ on } \delta\Omega$

General technique: find an approximate solution that is a linear combination of known functions

$$u^*(x,y) = \sum_i c_i \Phi_i(x,y)$$

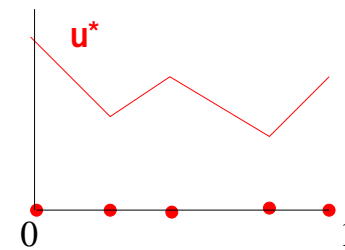
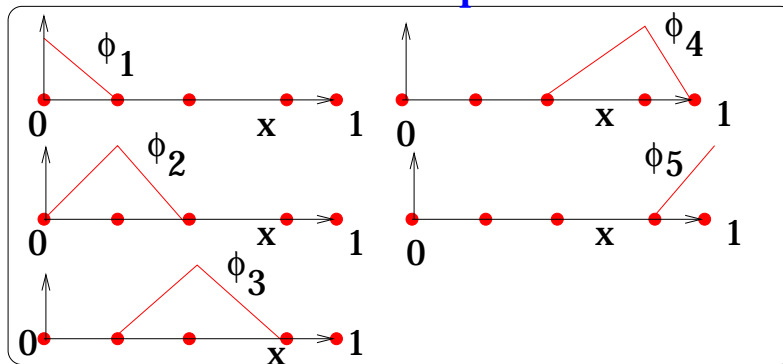
Question: How do we choose the known functions?

How do we find the best choice of c's, given the functions?

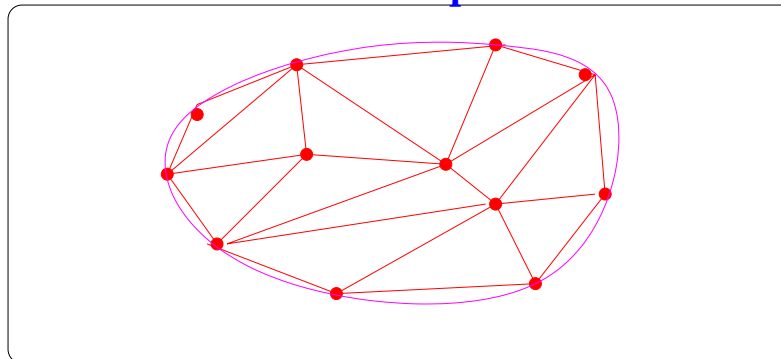
Choice of known functions:

- periodic boundary conditions: can use sines and cosines
- finite element method : generate a mesh that discretizes the domain
use low degree piecewise polynomials on mesh

1-D example



2-D example

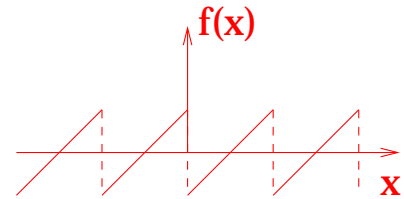


Mesh generation

Finding the best choices of the coefficients:

Analogy with Fourier series:

$$f(x) = a_0 + \sum_i a_i \cos(ix) + \sum_i b_i \sin(ix)$$



How do you find 'best' choices for a's and b's?

$$\begin{aligned} \int_{-\pi}^{+\pi} f(x) \cos(kx) dx &= \int_{-\pi}^{+\pi} (a_0 + \sum_i a_i \cos(ix) + \sum_i b_i \sin(ix)) \cos(kx) dx \\ &= \int_{-\pi}^{+\pi} a_k \cos(kx) \cos(kx) dx \\ &= a_k \pi \end{aligned}$$

Key idea:

- residual $f(x) - a_0 + \sum_i a_i \cos(ix) + \sum_i b_i \sin(ix)$
- weight residual by known function and integrate to find corresponding coefficient

Weighted Residual Technique:

$$\text{Residual: } (\mathbf{L} \mathbf{u}^* - \mathbf{f}) = (\mathbf{L} (\sum_i^N \mathbf{c}_i \phi_i) - \mathbf{f})$$

$$\text{Weighted Residual} = (\mathbf{L} (\sum_i^N \mathbf{c}_i \phi_i) - \mathbf{f}) \phi_k$$

$$\text{Equation for } k^{\text{th}} \text{ unknown: } \int_{\Omega} \phi_k^* (\mathbf{L} (\sum_i^N \mathbf{c}_i \phi_i) - \mathbf{f}) dV = 0 \Rightarrow$$

If the differential equation is linear:

$$\mathbf{c}_1 \int_{\Omega} \phi_k^* \mathbf{L} \phi_1 dV + \dots + \mathbf{c}_N \int_{\Omega} \phi_k^* \mathbf{L} \phi_N dV = \int_{\Omega} \phi_k^* \mathbf{f} dV$$

$k = 1, 2, \dots, N$

This system can be written as

$$\mathbf{K} \mathbf{c} = \mathbf{b} \text{ where}$$

$$\mathbf{K}(i,j) = \int_{\Omega} \phi_i^* \mathbf{L} \phi_j dV \quad \mathbf{b}(i) = \int_{\Omega} \phi_i^* \mathbf{f} dV$$

Key insight: Calculus problem of solving pde is converted to linear algebra problem of solving $\mathbf{K} \mathbf{c} = \mathbf{b}$ where \mathbf{K} is sparse

Solving system of linear algebraic equations:

- $K c = b$
- Orders of magnitude for realistic problems
 - large (~ 10 million unknowns) (roughly equal to number of mesh points)
 - sparse (~ 100 non-zero entries per row)
(roughly equal to connectivity of a point)
 - same K, many b's in some problems
- Algorithms:
 - iterative methods (Jacobi, conjugate gradient, GMRES)

start with an initial approximation to solution
and keep refining it till you get close enough
 - factorization methods (LU, Cholesky, QR)

factorize K into LU where L is lower triangular and U is upper triangular
 $LUc = b$
Solve for c by solving two triangular systems

Jacobi: a (slow) iterative solver

Example:

$$4x + 2y = 8$$

$$3x + 4y = 11$$

Iterative system:

$$x_{n+1} = (8 - 2y_n)/4$$

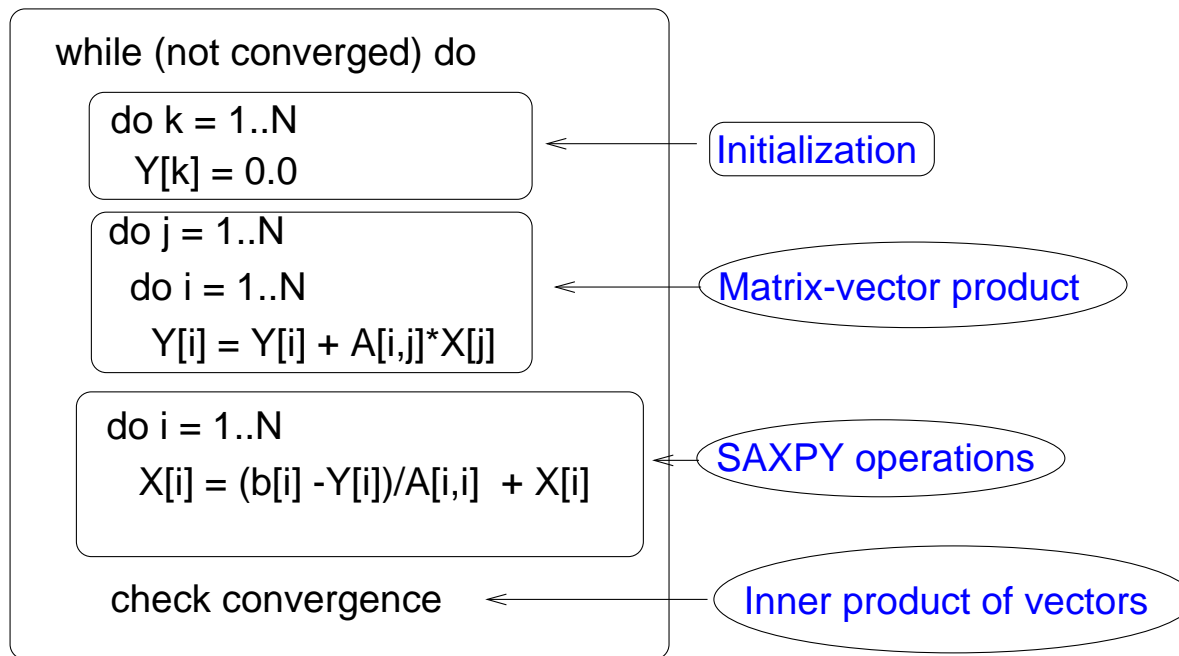
$$y_{n+1} = (11 - 3x_n)/4$$

n	1	2	3	4	5	6	7	8	
-	-	-	-	-	-	-	-	-	
x	0	2	0.625	1.375	0.8594	1.1406	0.9473	1.0527	...
=									
y	0	2.75	1.250	2.281	1.7188	2.1055	1.8945	2.0396	...

Matrix view of Jacobi Iteration

Iterative method for solving linear systems $Ax = b$

Jacobi method: $M * X_{k+1} = (M - A) * X_k + b$ (M is DIAGONAL(A))



Matrix-vector product: $O(N^2)$ work

SAXPY, Inner product: $O(N)$ work

Most of the time is spent in matrix-vector product.

Lesson for software systems people: optimize MVM

Reality check:

- Jacobi is a very old method of solving linear systems iteratively.
- More modern methods: conjugate gradient (CG), GMRES, etc. converge faster in most cases.
- However, the structure of these algorithms is similar: **MVM is the key operation.**
- Major area of research in numerical analysis: speeding up iterative algorithms further by *preconditioning*.

Tangential Discussion

- Calculus problem $Lu = f \Rightarrow$ linear algebra problem $Kc = b$.
- In some problems, we need to solve for multiple variables at each mesh point (temperature, pressure, velocity etc.)
 \Rightarrow solve many linear equations with same K , different b 's.
- This is viewed as matrix equation $KC = B$ where C and B are matrices.
- Algorithms for solving single system can be used to solve multiple systems as well.
- **Key computation in iterative methods: matrix-matrix multiplication (MMM) rather than matrix-vector multiplication (MVM).**
- Non-linear pde's lead to non-linear algebraic systems which are solved iteratively (Newton's method etc.).
Key computation: MMM or MVM.

Computational Requirements

Let us estimate storage and time requirements.

- Assume 10^6 mesh points (rows/columns of A)
- Assume iterative solver needs 100 iterations to converge
- Assume simulation runs for 1000 time steps.

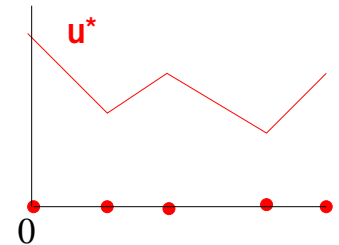
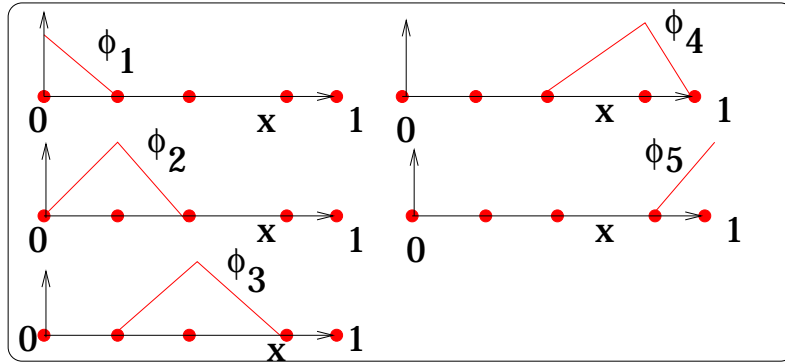
One MVM requires roughly 10^{12} flops

=>

Overall simulation requires 10^{17} flops and 10^{12} bytes of storage!

Can we do better?

1-D case



$$K(i,j) = \int_{\Omega} \phi_i * L(\phi_j) d\Omega$$

Structure of the K matrix for any pde: $K[i,j]$ is 0 if ϕ_i and ϕ_j do not overlap!

For our example, K is

$$\begin{bmatrix} x & x & 0 & 0 & 0 \\ x & x & x & 0 & 0 \\ 0 & x & x & x & 0 \\ 0 & 0 & x & x & x \\ 0 & 0 & 0 & x & x \end{bmatrix}$$

Half the entries are zero!

In 2-D and 3-D, an even larger percentage of matrix is zero!

Typical 3-D numbers: 10^6 rows but only 100-500 non-zeros per row!

Matrix is sparse.

Exploiting sparsity

Store sparse matrices in special formats to avoid storing zeros

=> storage costs are reduced!

Avoid computing with zeros when working with sparse matrices.

=> MFlops needs are reduced!

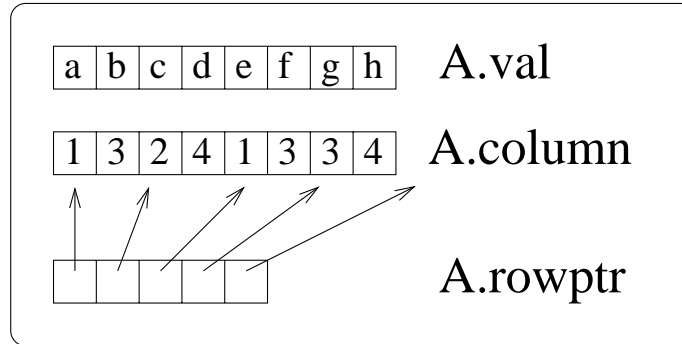
Question: How do we represent sparse matrices and how do we compute with them?

Three Sparse Matrix Representations

	1	2	3	4
1	a		b	
2		c		d
3	e		f	
4			g	h

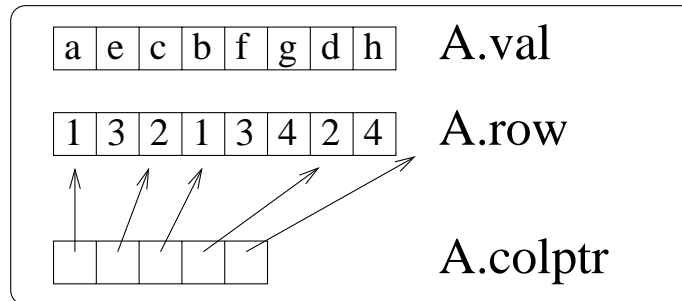
A

CRS



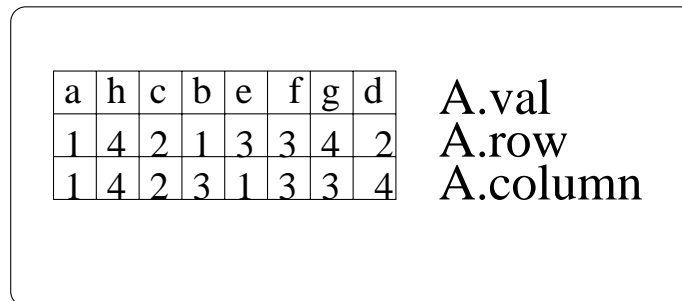
Indexed access to a row

CCS



Indexed access to a column

**Co-ordinate
Storage**



Indexed access to neither
rows nor columns

MVM for CRS

```
for I = 1 to N do
  for JJ = A.rowptr(I) to A.rowptr(I+1) -1 do
    Y(I) = Y(I) + A.val(JJ)*X(A.column(JJ))
  od
od
```

MVM for Co-ordinate storage

```
for P = 1 to NZ do
  Y(A.row(P)) = Y(A.row(P)) + A.val(P)*X(A.column(P))
od
```

Sparse matrix computations introduce subscripts with indirection.

Computational Requirements with sparse matrices

- Assume 10^6 mesh points (rows/columns of A).
- Assume roughly 100 non-zeros per row.
- Assume iterative solver needs 100 iterations to converge.
- Assume simulation runs for 1000 time steps.

One MVM requires roughly 10^8 flops

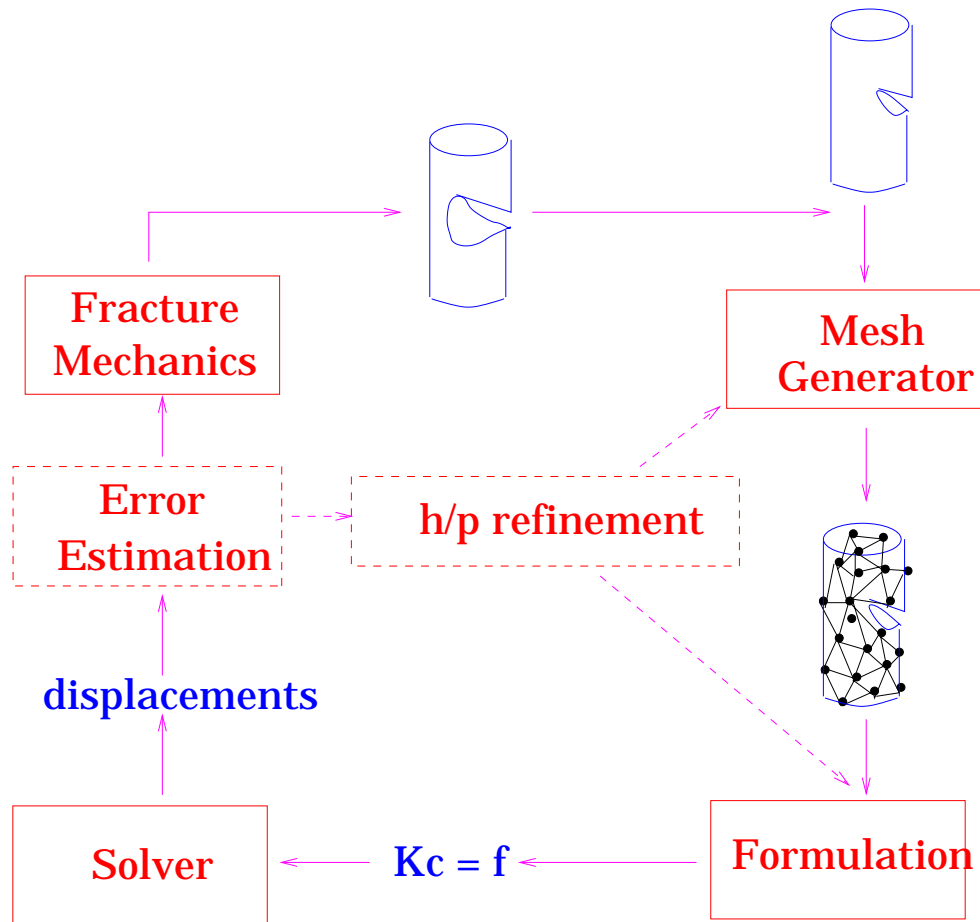
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Overall simulation requires 10^{13} flops and 10^8 bytes of storage!

This is roughly 100 seconds on a 100 Gflop supercomputer.

Doable!

Flow-chart of Adaptive Finite-element Simulation of Fracture



Summary

- Computational science applications: solving pde's or pushing particles
- PDE's are solved using approximate techniques like fe method
- Time-consuming part: solving large linear algebraic systems
- Two approaches: iterative methods and direct (factorization) methods
- Key operations in iterative methods:
Basic Linear Algebra Subroutines (BLAS)
 - Level-1 BLAS: inner-product of vectors, saxpy
 - Level-2 BLAS: matrix-vector product, triangular-solve
 - Level-3 BLAS: matrix-matrix product, triangular-solve with multiple right-hand-sides
- Important to exploit sparsity in matrix
- Exploiting sparsity complicates code.