

1. We show that all functions that are representable in the theory  $\mathcal{Q}$  are computable.

Recall that an  $n$ -ary function  $f : \mathbb{N}^n \rightarrow \mathbb{N}$  is representable in the theory  $\mathcal{Q}$  if there is an  $(n + 1)$ -ary predicate  $R_f$ , such that for all  $x_1, \dots, x_n \in \mathbb{N}$ :

- $f(x_1, \dots, x_n) = y$  implies  $\models_{\mathcal{Q}} R_f(\underline{x_1}, \dots, \underline{x_n}, \underline{y})$ .
- $f(x_1, \dots, x_n) \neq y$  implies  $\models_{\mathcal{Q}} \sim R_f(\underline{x_1}, \dots, \underline{x_n}, \underline{y})$ .

Furthermore, recall that the theory  $\mathcal{Q}$  is finitely axiomatizable; hence, it is possible (i.e., computable) to enumerate all of the derivations that follow from the axioms of  $\mathcal{Q}$ . Therefore, to calculate  $f(x_1, \dots, x_n)$ , it suffices to search for a  $y$  such that there exists a proof of  $R_f(\underline{x_1}, \dots, \underline{x_n}, \underline{y})$ .

2. Let  $B_1(x)$  and  $B_2(x)$  be formulas in the language  $\mathcal{Q}$  with  $x$  as the sole free variable. We show how to construct formulas  $G_1$  and  $G_2$  such that

$$\models_{\mathcal{Q}} G_1 \Leftrightarrow B_1(\ulcorner G_2 \urcorner) \quad \text{and} \quad \models_{\mathcal{Q}} G_2 \Leftrightarrow B_2(\ulcorner G_1 \urcorner)$$

Recall that  $\mathcal{Q}$  is a theory in which *diag* is representable. Hence, by the Diagonal Lemma, for any formula  $B(x)$  with  $x$  as the sole free variable, there is a formula  $G$  such that  $\models_{\mathcal{Q}} G \Leftrightarrow B(\ulcorner G \urcorner)$ .

Note that  $B_1(\ulcorner B_2(x) \urcorner)$  is a formula with  $x$  as the sole free variable. Hence, there is a formula  $G$  such that  $\models_{\mathcal{Q}} G \Leftrightarrow B_1(\ulcorner B_2(\ulcorner G \urcorner) \urcorner)$ . Define  $G_1 = G$  and  $G_2 = B_2(\ulcorner G \urcorner)$ . Then

$$\models_{\mathcal{Q}} G_1 \Leftrightarrow B_1(\ulcorner G_2 \urcorner) \quad \equiv \quad \models_{\mathcal{Q}} G \Leftrightarrow B_1(\ulcorner B_2(\ulcorner G \urcorner) \urcorner)$$

and

$$\models_{\mathcal{Q}} G_2 \Leftrightarrow B_2(\ulcorner G_1 \urcorner) \quad \equiv \quad \models_{\mathcal{Q}} B_2(\ulcorner G \urcorner) \Leftrightarrow B_2(\ulcorner G \urcorner)$$

3. Let  $Prov$  be a provability predicate for the theory  $\mathcal{Q}$  and  $X$  and  $Y$  be formulas in the language of  $\mathcal{Q}$ . Assume  $\models_{\mathcal{Q}} Prov(\ulcorner X \urcorner) \supset Y$  and  $\models_{\mathcal{Q}} Prov(\ulcorner Y \urcorner) \supset X$ . We show that both  $X$  and  $Y$  are theorems in  $\mathcal{Q}$ .

Recall the properties of a provability predicate:

- (P1) for all  $X$ , if  $\models_{\mathcal{Q}} X$ , then  $\models_{\mathcal{Q}} Prov(\ulcorner X \urcorner)$
- (P2) for all  $X$  and  $Y$ ,  $\models_{\mathcal{Q}} Prov(\ulcorner X \supset Y \urcorner) \supset (Prov(\ulcorner X \urcorner) \supset Prov(\ulcorner Y \urcorner))$
- (P3) for all  $X$ ,  $\models_{\mathcal{Q}} Prov(\ulcorner X \urcorner) \supset Prov(\ulcorner Prov(\ulcorner X \urcorner) \urcorner)$

(NOTE: (P1) is *not* equivalent to “for all  $X$ ,  $\models_{\mathcal{Q}} X \supset Prov(\ulcorner X \urcorner)$ ”.)

Recall that  $\mathcal{Q}$  is a theory that can represent the computable functions. Hence, by Löb’s Theorem, we have:

- (L) for all  $X$ , if  $\models_{\mathcal{Q}} Prov(\ulcorner X \urcorner) \supset X$ , then  $\models_{\mathcal{Q}} X$

Finally, recall that  $\mathcal{Q}$  is a consistent theory; hence, the following *modus ponens* style theorems are true:

- (MP1) for all  $X$  and  $Y$ , if  $\models_{\mathcal{Q}} X \supset Y$  and  $\models_{\mathcal{Q}} X$ , then  $\models_{\mathcal{Q}} Y$
- (MP2) for all  $X$  and  $Y$  and  $Z$ , if  $\models_{\mathcal{Q}} X \supset Y$  and  $\models_{\mathcal{Q}} Y \supset Z$ , then  $\models_{\mathcal{Q}} X \supset Z$

We reason as follows:

- (A1)  $\models_{\mathcal{Q}} Prov(\ulcorner X \urcorner) \supset Y$ , by assumption.
- (A2)  $\models_{\mathcal{Q}} Prov(\ulcorner Y \urcorner) \supset X$ , by assumption.
- (B1)  $\models_{\mathcal{Q}} Prov(\ulcorner Prov(\ulcorner X \urcorner) \supset Y \urcorner)$ , by (A1) and (P1).
- (B2)  $\models_{\mathcal{Q}} Prov(\ulcorner Prov(\ulcorner Y \urcorner) \supset X \urcorner)$ , by (A2) and (P1).
- (C1)  $\models_{\mathcal{Q}} Prov(\ulcorner Prov(\ulcorner X \urcorner) \urcorner) \supset Prov(\ulcorner Y \urcorner)$ , by (P2), (B1), and (MP1).
- (C2)  $\models_{\mathcal{Q}} Prov(\ulcorner Prov(\ulcorner Y \urcorner) \urcorner) \supset Prov(\ulcorner X \urcorner)$ , by (P2), (B2), and (MP1).
- (D1)  $\models_{\mathcal{Q}} Prov(\ulcorner Prov(\ulcorner X \urcorner) \urcorner) \supset X$ , by (C1), (A2), and (MP2).
- (D2)  $\models_{\mathcal{Q}} Prov(\ulcorner Prov(\ulcorner Y \urcorner) \urcorner) \supset Y$ , by (C2), (A1), and (MP2).
- (E1)  $\models_{\mathcal{Q}} Prov(\ulcorner X \urcorner) \supset X$ , by (P3), (D1), and (MP2).
- (E2)  $\models_{\mathcal{Q}} Prov(\ulcorner Y \urcorner) \supset Y$ , by (P3), (D2), and (MP2).
- (F1)  $\models_{\mathcal{Q}} X$ , by (E1) and (L).
- (F2)  $\models_{\mathcal{Q}} Y$ , by (E2) and (L).

Thus, both  $X$  and  $Y$  are theorems in  $\mathcal{Q}$ .

4. A pair of natural numbers  $\langle a, b \rangle$  is called a *stamps pair* if for all  $n \geq a + b$  there are natural numbers  $i$  and  $j$  such that  $n = i * a + j * b$ .

(a) We prove in  $\lambda$ -PRL that  $\langle 3, 5 \rangle$  is a stamps pair.

```

 $\vdash \forall n:\text{int}. 8 \leq n \Rightarrow \exists i,j:\text{int}. (0 \leq i) \wedge (0 \leq j) \wedge (n = i * 3 + j * 5)$ 
|
| by induction 1 8 10 3
| \
| 1.  $n:\text{int}$ 
| 2.  $n < 8$ 
| 3.  $8 \leq n+1 \Rightarrow \exists i,j:\text{int}. (0 \leq i) \wedge (0 \leq j) \wedge (n+1 = i * 3 + j * 5)$ 
|  $\vdash 8 \leq n \Rightarrow \exists i,j:\text{int}. (0 \leq i) \wedge (0 \leq j) \wedge (n = i * 3 + j * 5)$ 
| |
| | by impR
| | ...
| 4.  $8 \leq n$ 
|  $\vdash \exists i,j:\text{int}. (0 \leq i) \wedge (0 \leq j) \wedge (n = i * 3 + j * 5)$ 
|
| by arith  $\checkmark$ 
| \
|  $\vdash 8 \leq 8 \Rightarrow \exists i,j:\text{int}. (0 \leq i) \wedge (0 \leq j) \wedge (8 = i * 3 + j * 5)$ 
| |
| | by impR
| 1.  $8 \leq 8$ 
|  $\vdash \exists i,j:\text{int}. (0 \leq i) \wedge (0 \leq j) \wedge (8 = i * 3 + j * 5)$ 
| |
| | by exR 1 1
| | ...
|  $\vdash (0 \leq 1) \wedge (0 \leq 1) \wedge (8 = 1 * 3 + 1 * 5)$ 
|
| by arith  $\checkmark$ 
| \
|  $\vdash 8 \leq 9 \Rightarrow \exists i,j:\text{int}. (0 \leq i) \wedge (0 \leq j) \wedge (9 = i * 3 + j * 5)$ 
| |
| | by impR
| 1.  $8 \leq 9$ 
|  $\vdash \exists i,j:\text{int}. (0 \leq i) \wedge (0 \leq j) \wedge (9 = i * 3 + j * 5)$ 
| |
| | by exR 3 0
| | ...
|  $\vdash (0 \leq 3) \wedge (0 \leq 0) \wedge (9 = 3 * 3 + 0 * 5)$ 
|
| by arith  $\checkmark$ 
| \
|  $\vdash 8 \leq 10 \Rightarrow \exists i,j:\text{int}. (0 \leq i) \wedge (0 \leq j) \wedge (10 = i * 3 + j * 5)$ 
| |
| | by impR
| 1.  $8 \leq 10$ 
|  $\vdash \exists i,j:\text{int}. (0 \leq i) \wedge (0 \leq j) \wedge (10 = i * 3 + j * 5)$ 
| |
| | by exR 0 2
| | ...
|  $\vdash (0 \leq 0) \wedge (0 \leq 2) \wedge (10 = 0 * 3 + 2 * 5)$ 
|
| by arith  $\checkmark$ 
| \
| 1.  $n:\text{int}$ 
| 2.  $10 < n$ 
| 3.  $8 \leq n-3 \Rightarrow \exists i,j:\text{int}. (0 \leq i) \wedge (0 \leq j) \wedge (n-3 = i * 3 + j * 5)$ 
|  $\vdash 8 \leq n \Rightarrow \exists i,j:\text{int}. (0 \leq i) \wedge (0 \leq j) \wedge (n = i * 3 + j * 5)$ 
|
| by impL 3
| \
| | ...
| |  $\vdash 8 \leq n-3$ 
| |
| | by arith  $\checkmark$ 
| | \
| | | ...
| | 3.  $8 \leq n-3$ 
| | 4.  $\exists i,j:\text{int}. (0 \leq i) \wedge (0 \leq j) \wedge (n-3 = i * 3 + j * 5)$ 
| |  $\vdash 8 \leq n \Rightarrow \exists i,j:\text{int}. (0 \leq i) \wedge (0 \leq j) \wedge (n = i * 3 + j * 5)$ 
| |
| | by exR 4

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```

|
...
4. i:int
5. j:int
6. (0 ≤ i) ∧ (0 ≤ j) ∧ (n-3 = i * 3 + j * 5)
⊢ 8 ≤ n ⇒ ∃i,j:int. (0 ≤ i) ∧ (0 ≤ j) ∧ (n = i * 3 + j * 5)
|
...
6. 0 ≤ i
7. 0 ≤ j
8. n-3 = i * 3 + j * 5
⊢ 8 ≤ n ⇒ ∃i,j:int. (0 ≤ i) ∧ (0 ≤ j) ∧ (n = i * 3 + j * 5)
|
...
9. 8 ≤ n
⊢ ∃i,j:int. (0 ≤ i) ∧ (0 ≤ j) ∧ (n = i * 3 + j * 5)
|
...
⊢ (0 ≤ i+1) ∧ (0 ≤ j) ∧ (n = (i+1) * 3 + j * 5)
by arith √

```

(b) We describe the algorithm implicitly contained in the proof of (a) in  $\lambda$ -PRL notation.

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stampsi(n:int):int ≡  n => 0
                    8 => 1
                    9 => 3
                   10 => 0
                    n => stampsi(n-3) + 1

stampsj(n:int):int ≡  n => 0
                    8 => 1
                    9 => 0
                   10 => 2
                    n => stampsj(n-3)

```

(c) If  $\langle a, b \rangle$  is a stamps pair (such that  $a \leq b$ ), then

- $a = 1$ , or
- $a = 2$  and  $b = 1 \bmod 2$ , or
- $a = 3$  and  $b = 4$ , or
- $a = 3$  and  $b = 5$

5. We prove in  $\lambda$ -PRL that every non-empty list of integers has a maximal element.

$$\forall l:\text{list}. (\sim(l=[])) \Rightarrow \exists i:\text{int}. (i \in l \wedge \forall j:\text{int}. (j \in l \Rightarrow j \leq i))$$

We give a  $\lambda$ -PRL definition of the predicate  $i \in l$ :

$$i \in l \equiv \exists k:\text{int}. (\text{index}(k,l)=i \wedge 0 \leq k \wedge k < \text{length}(l))$$

We prove the above statement  $\lambda$ -PRL's refinement logic:

```

 $\vdash \forall l:\text{list}. (\sim(l=[])) \Rightarrow \exists i:\text{int}. (i \in l \wedge \forall j:\text{int}. (j \in l \Rightarrow j \leq i))$ 
|
| \
| |  $\vdash \sim([]=[]) \Rightarrow \exists i:\text{int}. (i \in [] \wedge \forall j:\text{int}. (j \in [] \Rightarrow j \leq i))$ 
| |
| | 1.  $\sim([]=[])$ 
| |  $\vdash \exists i:\text{int}. (i \in [] \wedge \forall j:\text{int}. (j \in [] \Rightarrow j \leq i))$ 
| |
| | ...
| |  $\vdash []=[]$ 
| |
| | \
| | 1.  $x:\text{int}$ 
| |  $\vdash \sim([x]=[]) \Rightarrow \exists i:\text{int}. (i \in [x] \wedge \forall j:\text{int}. (j \in [x] \Rightarrow j \leq i))$ 
| |
| | ...
| | 2.  $\sim([x]=[])$ 
| |  $\vdash \exists i:\text{int}. (i \in [x] \wedge \forall j:\text{int}. (j \in [x] \Rightarrow j \leq i))$ 
| |
| | ...
| |  $\vdash x \in [x] \wedge \forall j:\text{int}. (j \in [x] \Rightarrow j \leq x)$ 
| |
| | \
| | ...
| |  $\vdash x \in [x]$ 
| |
| | ...
| |  $\vdash \exists k:\text{int}. (\text{index}(k,[x])=i \wedge 0 \leq k \wedge k < \text{length}([x]))$ 
| |
| | ...
| |  $\vdash \text{index}(0,[x])=x \wedge 0 \leq 0 \wedge 0 < \text{length}([x])$ 
| |
| | ...
| |  $\vdash x=x \wedge 0 \leq 0 \wedge 0 < 1$ 
| |
| | \
| | ...
| |  $\vdash \forall j:\text{int}. (j \in [x] \Rightarrow j \leq x)$ 
| |
| | ...
| | 3.  $j:\text{int}$ 
| |  $\vdash j \in [x] \Rightarrow j \leq x$ 
| |
| | ...
| | 4.  $j \in [x]$ 
| |  $\vdash j \leq x$ 
| |
| | ...
| | 4.  $\exists k:\text{int}. (\text{index}(k,[x])=j \wedge 0 \leq k \wedge k < \text{length}([x]))$ 
| |  $\vdash j \leq x$ 
| |

```

by induction x::l  
 by impR  
 by notL 1  
 by equality  $\surd$   
 by impR  
 by exR x  
 by andR  
 by unfold  $\in$  in 0  
 by exR 0  
 by simplify in 0  
 by arith  $\surd$   
 by allR  
 by impR  
 by unfold  $\in$  in 4  
 by exL 4

```

| |
| | ...
| | 4. k:int
| | 5. index(k,[x])=j  $\wedge$  0 $\leq$ k  $\wedge$  k<length([x])
| |  $\vdash$  j $\leq$ x
| | |
| | | by simplify in 5
| | | ...
| | | 5. index(k,[x])=j  $\wedge$  0 $\leq$ k  $\wedge$  k<1
| | |  $\vdash$  j $\leq$ x
| | | |
| | | | by andL 5
| | | | ...
| | | | 5. index(k,[x])=j
| | | | 6. 0 $\leq$ k
| | | | 7. k<1
| | | |  $\vdash$  j $\leq$ x
| | | | |
| | | | | by cut k=0
| | | | | \
| | | | | | ...
| | | | | |  $\vdash$  k=0
| | | | | | |
| | | | | | | by arith  $\checkmark$ 
| | | | | | | \
| | | | | | | | ...
| | | | | | | | 8. k=0
| | | | | | | |  $\vdash$  j $\leq$ x
| | | | | | | | |
| | | | | | | | | by cut index(k,[x])=j
| | | | | | | | | \
| | | | | | | | | | ...
| | | | | | | | | |  $\vdash$  index(k,[x])=j
| | | | | | | | | | |
| | | | | | | | | | | by hyp 5  $\checkmark$ 
| | | | | | | | | | | \
| | | | | | | | | | | | ...
| | | | | | | | | | | | 9. index(k,[x])=j
| | | | | | | | | | | |  $\vdash$  j $\leq$ x
| | | | | | | | | | | | |
| | | | | | | | | | | | | by subst 8 9
| | | | | | | | | | | | | \
| | | | | | | | | | | | | | ...
| | | | | | | | | | | | | | 9. index(0,[x])=j
| | | | | | | | | | | | | |  $\vdash$  j $\leq$ x
| | | | | | | | | | | | | | |
| | | | | | | | | | | | | | | by simplify 9
| | | | | | | | | | | | | | | \
| | | | | | | | | | | | | | | | ...
| | | | | | | | | | | | | | | | 9. x=j
| | | | | | | | | | | | | | | |  $\vdash$  j $\leq$ x
| | | | | | | | | | | | | | | | |
| | | | | | | | | | | | | | | | | by arith  $\checkmark$ 
| | \
| | 1. x:int
| | 2. l:list
| | 3.  $\sim(1=[]) \Rightarrow \exists i:int. (i \in 1 \wedge \forall j:int. (j \in 1 \Rightarrow j \leq i))$ 
| |  $\vdash \sim(x::l=[]) \Rightarrow \exists i:int. (i \in x::l \wedge \forall j:int. (j \in x::l \Rightarrow j \leq i))$ 
| | |
| | | by cut (1=[])  $\vee$   $\sim(1=[])$ 
| | | \
| | | | ...
| | | |  $\vdash (1=[]) \vee \sim(1=[])$ 
| | | | |
| | | | | by equality  $\checkmark$ 
| | | | | \
| | | | | | ...
| | | | | | 4. (1=[])  $\vee$   $\sim(1=[])$ 
| | | | | |  $\vdash \sim(x::l=[]) \Rightarrow \exists i:int. (i \in x::l \wedge \forall j:int. (j \in x::l \Rightarrow j \leq i))$ 
| | | | | | |
| | | | | | | by orL 4
| | | | | | | \
| | | | | | | | ...
| | | | | | | | 4. l=[]
| | | | | | | |  $\vdash \sim(x::l=[]) \Rightarrow \exists i:int. (i \in x::l \wedge \forall j:int. (j \in x::l \Rightarrow j \leq i))$ 
| | | | | | | | |
| | | | | | | | | by subst 4 0
| | | | | | | | | \
| | | | | | | | | | ...
| | | | | | | | | |  $\vdash \exists i:int. (i \in x::[] \wedge \forall j:int. (j \in x::[] \Rightarrow j \leq i))$ 
| | | | | | | | | | |
| | | | | | | | | | | by simplify 0

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| |
| ...
| ⊢ ∃i:int. (i∈[x] ∧ ∀j:int. (j∈[x] ⇒ j≤i))
|
| \
| ...
| 4. ∼(l=[])
| ⊢ ∼(x::l=[]) ⇒ ∃i:int. (i∈x::l ∧ ∀j:int. (j∈x::l ⇒ j≤i))
| |
| | by impR
| ...
| 5. ∼(x::l=[])
| ⊢ ∃i:int. (i∈x::l ∧ ∀j:int. (j∈x::l ⇒ j≤i))
| |
| | by impL 3
| | \
| | | ...
| | | ⊢ ∼(l=[])
| | |
| | | \
| | | ...
| | | 3. ∼(l=[])
| | | 4. ∼(x::l=[])
| | | 5. ∼(l=[])
| | | 6. ∃i:int. (i∈l ∧ ∀j:int. (j∈l ⇒ j≤i))
| | | ⊢ ∃i:int. (i∈x::l ∧ ∀j:int. (j∈x::l ⇒ j≤i))
| | | |
| | | | by exL 6
| | | ...
| | | 6. i:int
| | | 7. i∈l ∧ ∀j:int. (j∈l ⇒ j≤i)
| | | ⊢ ∃i:int. (i∈x::l ∧ ∀j:int. (j∈x::l ⇒ j≤i))
| | | |
| | | | by andL 7
| | | ...
| | | 7. i∈l
| | | 8. ∀j:int. (j∈l ⇒ j≤i)
| | | ⊢ ∃i:int. (i∈x::l ∧ ∀j:int. (j∈x::l ⇒ j≤i))
| | | |
| | | | by unfold ∈ 7
| | | ...
| | | 7. ∃k:int. (index(k,l)=i ∧ 0≤k ∧ k<length(x::l))
| | | 8. ∀j:int. (j∈l ⇒ j≤i)
| | | ⊢ ∃i:int. (i∈x::l ∧ ∀j:int. (j∈x::l ⇒ j≤i))
| | | |
| | | | by exL 7
| | | ...
| | | 7. k:int
| | | 8. index(k,l)=i ∧ 0≤k ∧ k<length(x::l)
| | | 9. ∀j:int. (j∈l ⇒ j≤i)
| | | ⊢ ∃i:int. (i∈x::l ∧ ∀j:int. (j∈x::l ⇒ j≤i))
| | | |
| | | | by andL 8
| | | ...
| | | 8. index(k,l)=i
| | | 9. 0≤k
| | | 10. k<length(x::l)
| | | 11. ∀j:int. (j∈l ⇒ j≤i)
| | | ⊢ ∃i:int. (i∈x::l ∧ ∀j:int. (j∈x::l ⇒ j≤i))
| | | |
| | | | by simplify 10
| | | ...
| | | 8. index(k,l)=i
| | | 9. 0≤k
| | | 10. k<length(l)+1
| | | 11. ∀j:int. (j∈l ⇒ j≤i)
| | | ⊢ ∃i:int. (i∈x::l ∧ ∀j:int. (j∈x::l ⇒ j≤i))
| | | |
| | | | by cut (x ≤ i) ∨ (x > i)
| | | \
| | | | ...
| | | | ⊢ (x ≤ i) ∨ (x > i)
| | | |
| | | | by arith √

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|
| \
| ...
| 12. (x ≤ i) ∨ (x > i)
| ⊢ ∃i:int. (i∈x::l ∧ ∀j:int. (j∈x::l ⇒ j≤i))
| | | | | by orL 12
| | \
| | | ...
| | | 12. x ≤ i
| | | ⊢ ∃i:int. (i∈x::l ∧ ∀j:int. (j∈x::l ⇒ j≤i))
| | | | | by exR i
| | | | ...
| | | | ⊢ i∈x::l ∧ ∀j:int. (j∈x::l ⇒ j≤i)
| | | | | by andR
| | | | \
| | | | | ...
| | | | | ⊢ i∈x::l
| | | | | | | by unfold ∈ in 0
| | | | | | ...
| | | | | | ⊢ ∃k:int. (index(k,x::l)=i ∧ 0≤k ∧ k<length(x::l))
| | | | | | | by exR k+1
| | | | | | ...
| | | | | | ⊢ index(k+1,x::l)=i ∧ 0≤k+1 ∧ k+1<length(x::l)
| | | | | | | by def index(k+1,x::l) up
| | | | | | \
| | | | | | | ...
| | | | | | | ⊢ 0 < k+1
| | | | | | | | | by arith √
| | | | | | | \
| | | | | | | | ...
| | | | | | | | 13. 0 < k+1
| | | | | | | | 14. index(k+1,x::l) = index(k+1-1,tl(x::l))
| | | | | | | | ⊢ index(k+1,x::l)=i ∧ 0≤k+1 ∧ k+1<length(x::l)
| | | | | | | | | by subst 14 0
| | | | | | | | | ...
| | | | | | | | | ⊢ index(k+1-1,tl(x::l))=i ∧ 0≤k+1 ∧ k+1<length(x::l)
| | | | | | | | | | by simplify 0
| | | | | | | | | ...
| | | | | | | | | ⊢ index(k,l)=i ∧ 0≤k+1 ∧ k+1<length(l)+1
| | | | | | | | | | by arith √
| | | | | | | \
| | | | | | | | ...
| | | | | | | | ⊢ ∀j:int. (j∈x::l ⇒ j≤i)
| | | | | | | | | by allR
| | | | | | | | ...
| | | | | | | | 13. j: int
| | | | | | | | ⊢ j∈x::l ⇒ j≤i
| | | | | | | | | by impR
| | | | | | | | ...
| | | | | | | | 14. j∈x::l
| | | | | | | | ⊢ j≤i
| | | | | | | | | by unfold ∈ in 14
| | | | | | | | ...
| | | | | | | | 14. ∃k:int. (index(k,x::l)=j ∧ 0≤k ∧ k<length(x::l))
| | | | | | | | ⊢ j≤i
| | | | | | | | | by exL 14
| | | | | | | | ...
| | | | | | | | 14. k':int
| | | | | | | | 15. index(k',x::l)=j ∧ 0≤k' ∧ k'<length(x::l)
| | | | | | | | ⊢ j≤i
| | | | | | | | | by andL 15
| | | | | | | |

```

	...	
	15. $\text{index}(k', x :: 1) = j$	
	16. $0 \leq k'$	
	17. $k' < \text{length}(x :: 1)$	
	$\vdash j \leq i$	
		by simplify 17
	...	
	17. $k' < \text{length}(l) + 1$	
	$\vdash j \leq i$	
		cut $k' = 0 \vee k' > 0$
	\	
	...	
	$\vdash k' = 0 \vee k' > 0$	
		by arith $\sqrt{\quad}$
	\	
	...	
	18. $k' = 0 \vee k' > 0$	
	$\vdash j \leq i$	
		by orL 18
	\	
	...	
	18. $k' = 0$	
	$\vdash j \leq i$	
		by cut $\text{index}(k', x :: 1) = j$
	\	
	...	
	$\vdash \text{index}(k', x :: 1) = j$	
		by hyp 15
	\	
	...	
	19. $\text{index}(k', x :: 1) = j$	
	$\vdash j \leq i$	
		by subst 18 19
	...	
	19. $\text{index}(0, x :: 1) = j$	
	$\vdash j \leq i$	
		by simplify 19
	...	
	19. $x = j$	
	$\vdash j \leq i$	
		by arith $\sqrt{\quad}$
	\	
	...	
	18. $k' > 0$	
	$\vdash j \leq i$	
		by allL 11 j
	...	
	19. $j \in 1 \Rightarrow j \leq i$	
	$\vdash j \leq i$	
		by impL 19
	\	
	...	
	$\vdash j \in 1$	
		by unfold $\in$ in 0
	...	
	$\vdash \exists k : \text{int}. (\text{index}(k, l) = j \wedge 0 \leq k \wedge k < \text{length}(l))$	
		by exR $k' - 1$
	...	
	$\vdash \text{index}(k' - 1, l) = j \wedge 0 \leq k' - 1 \wedge k' - 1 < \text{length}(l)$	
		by def $\text{index}(k', x :: 1)$ up
	\	
	...	
	$\vdash 0 \leq k'$	
		by hyp 18

```

|   | |
|   | | \
|   | | ...
|   | | 20.  $0 \leq k'$ 
|   | | 21.  $\text{index}(k', x::l) = \text{index}(k'-1, \text{tl}(x::l))$ 
|   | |  $\vdash \text{index}(k'-1, l) = j \wedge 0 \leq k'-1 \wedge k'-1 < \text{length}(l)$ 
|   | | | by simplify 21
|   | | ...
|   | | 21.  $\text{index}(k', x::l) = \text{index}(k'-1, l)$ 
|   | |  $\vdash \text{index}(k'-1, l) = j \wedge 0 \leq k'-1 \wedge k'-1 < \text{length}(l)$ 
|   | | | by arith  $\checkmark$ 
|   | | \
|   | | 19.  $j \in l$ 
|   | | 20.  $j \leq i$ 
|   | |  $\vdash j \leq i$ 
|   | | | by hyp 20  $\checkmark$ 
|   | | \
|   | | ...
|   | | 12.  $x > i$ 
|   | |  $\vdash \exists i:\text{int}. (i \in x::l \wedge \forall j:\text{int}. (j \in x::l \Rightarrow j \leq i))$ 
|   | | | by exR x
|   | | ...
|   | |  $\vdash x \in x::l \wedge \forall j:\text{int}. (j \in x::l \Rightarrow j \leq x)$ 
|   | | | by andR
|   | | | \
|   | | | ...
|   | | |  $\vdash x \in x::l$ 
|   | | | | by unfold  $\in$  in 0
|   | | | ...
|   | | |  $\vdash \exists k:\text{int}. (\text{index}(k, x::l) = x \wedge 0 \leq k \wedge k < \text{length}(x::l))$ 
|   | | | | by exR 0
|   | | | ...
|   | | |  $\vdash \text{index}(0, x::l) = x \wedge 0 \leq 0 \wedge 0 < \text{length}(x::l)$ 
|   | | | | by simplify 0
|   | | | ...
|   | | |  $\vdash x = x \wedge 0 \leq 0 \wedge 0 < \text{length}(l) + 1$ 
|   | | | | by arith  $\checkmark$ 
|   | | | \
|   | | | ...
|   | | |  $\vdash \forall j:\text{int}. (j \in x::l \Rightarrow j \leq x)$ 
|   | | | | by allR
|   | | | ...
|   | | | 13.  $j:\text{int}$ 
|   | | |  $\vdash j \in x::l \Rightarrow j \leq x$ 
|   | | | | by impR
|   | | | ...
|   | | | 14.  $j \in x::l$ 
|   | | |  $\vdash j \leq x$ 
|   | | | | by unfold  $\in$  in 14
|   | | | ...
|   | | | 14.  $\exists k:\text{int}. (\text{index}(k, x::l) = j \wedge 0 \leq k \wedge k < \text{length}(x::l))$ 
|   | | |  $\vdash j \leq x$ 
|   | | | | by exL 14
|   | | | ...
|   | | | 14.  $k':\text{int}$ 
|   | | | 15.  $\text{index}(k', x::l) = j \wedge 0 \leq k' \wedge k' < \text{length}(x::l)$ 
|   | | |  $\vdash j \leq x$ 
|   | | | | by andL 15
|   | | | ...
|   | | | 15.  $\text{index}(k', x::l) = j$ 
|   | | | 16.  $0 \leq k'$ 
|   | | | 17.  $k' < \text{length}(x::l)$ 
|   | | |  $\vdash j \leq x$ 
|   | | | | by simplify 17

```

```

|
...
17.  $k' < \text{length}(l)+1$ 
 $\vdash j \leq x$ 
|
| \ by cut  $k'=0 \vee k'>0$ 
| \
| ...
|  $\vdash k'=0 \vee k'>0$ 
|
| \ by arith  $\sqrt{\quad}$ 
| \
| ...
18.  $k'=0 \vee k'>0$ 
 $\vdash j \leq x$ 
|
| \ by orL 18
| \
| ...
| 18.  $k'=0$ 
|  $\vdash j \leq x$ 
|
| \ by cut  $\text{index}(k', x::l)=j$ 
| \
| | ...
| |  $\vdash \text{index}(k', x::l)=j$ 
| |
| | \ by hyp 15
| | \
| | ...
| | 19.  $\text{index}(k', x::l)=j$ 
| |  $\vdash j \leq x$ 
| |
| | \ by subst 18 19
| | \
| | ...
| | 19.  $\text{index}(0, x::l)=j$ 
| |  $\vdash j \leq x$ 
| |
| | \ by simplify 19
| | \
| | ...
| | 19.  $x=j$ 
| |  $\vdash j \leq x$ 
| |
| | \ by arith  $\sqrt{\quad}$ 
| | \
| | ...
18.  $k'>0$ 
 $\vdash j \leq x$ 
|
| \ by allL 11 j
| \
| ...
19.  $j \in l \Rightarrow j \leq i$ 
 $\vdash j \leq x$ 
|
| \ by impL 19
| \
| | ...
| |  $\vdash j \in l$ 
| |
| | \ by unfold  $\in$  in 0
| | \
| | | ...
| | |  $\vdash \exists k:\text{int}. (\text{index}(k, l)=j \wedge 0 \leq k \wedge k < \text{length}(l))$ 
| | |
| | | \ by def  $\text{index}(k', x::l)=j$  up
| | | \
| | | | ...
| | | |  $\vdash \text{index}(k'-1, l)=j \wedge 0 \leq k'-1 \wedge k'-1 < \text{length}(l)$ 
| | | |
| | | | \ by hyp 18
| | | | \
| | | | ...
| | | | 20.  $0 \leq k'$ 
| | | |
| | | | \
| | | | 21.  $\text{index}(k', x::l) = \text{index}(k'-1, \text{tl}(x::l))$ 
| | | |  $\vdash \text{index}(k'-1, l)=j \wedge 0 \leq k'-1 \wedge k'-1 < \text{length}(l)$ 
| | | |
| | | | \ by simplify 21
| | | |

```

```

| |
| ...
| 20.  $0 \leq k'$ 
| 21.  $\text{index}(k', x :: l) = \text{index}(k'-1, l)$ 
|  $\vdash \text{index}(k'-1, l) = j \wedge 0 \leq k'-1 \wedge k'-1 < \text{length}(l)$ 
|
| by arith  $\checkmark$ 
\
...
19.  $j \in l$ 
20.  $j \leq i$ 
 $\vdash j \leq x$ 

```

by arith  $\checkmark$

We describe the extracted algorithm in  $\lambda$ -PRL notation.

```

listmax(l:list):int  $\equiv$ 
  0 => 0
  1 => hd(l)
  l => if tl(l) = []
      then hd(l)
      else if hd(l)  $\leq$  listmax(tl(l))
          then listmax(tl(l))
          else hd(l)

```