

CS4810 Dexter Kozen

Decision problems ($\{0,1\}$ -valued functions)

- all instances
- "yes" instances

Possible inputs: strings over a finite alphabet

- Alphabet - any finite set

Ex: $\{0,1\}$ $\{a,b\}$ $\{0,1,2,\dots,9\}$

- Σ elements are symbols or letters a, b, c, \dots

Strings are finite sequences from Σ .

Ex: $\Sigma = \{a,b\}$ $aabab$ a string of length 5

$$|aabab| = 5$$

- ϵ = null string or empty string $|\epsilon| = 0$

- a^n = string of a 's of length n ($n \in \mathbb{N}$)

$$|aaaaa| = 5 = |a^5|$$

$$(\mathbb{N} = \{0,1,2,\dots\})$$

$$a^0 \triangleq \epsilon \quad a^{n+1} \triangleq a^n a \quad \underline{\text{def}}$$

- Set of all strings over Σ is denoted Σ^*

$$\text{Ex: } \{a,b\}^* = \{\epsilon, a, b, aa, ab, ba, bb, aaa, \dots\}$$

$$\{a\}^* = \{\epsilon, a, aa, \dots\} = \{a^n \mid n \geq 0\}$$

$$\emptyset^* = \{\epsilon\} \quad \text{by convention}$$

$$|\varepsilon| = 0 \quad \text{~~xxxxxx~~ } |xa| = |x| + 1$$

$$|\cdot|: \Sigma^* \rightarrow \mathbb{N}$$

$$\{a, b\} = \{b, a\} \quad ab \neq ba$$

$$\{a, a, b\} = \{a, b\} \quad aab \neq ab$$

$$\text{--- } \{a, a, b\} \neq \{a, b\} = \{b, a\}$$

$$\text{Concatenation } \cdot: \Sigma^* \times \Sigma^* \rightarrow \Sigma^*$$

$$xy \in \Sigma^* \text{ if } x \in \Sigma^* \text{ \& } y \in \Sigma^*$$

$$\text{--- } (xy)z = x(yz)$$

$$\text{--- } |a^n| = n$$

$$\text{--- } \varepsilon x = x\varepsilon = x$$

$$\text{--- } a^m a^n = a^{m+n}$$

$$\text{--- } |xy| = |x| + |y|$$

$$\text{--- } |\varepsilon| = 0$$

$(M, \cdot, 1)$ where \cdot is associative

1 is a 2-sided identity

is called a monoid (semigroup w/ ident)

$(\Sigma^*, \cdot, \varepsilon)$ is a monoid

$(\mathbb{N}, +, 0)$ " " "

$|\cdot|: \Sigma^* \rightarrow \mathbb{N}$ is a monoid homomorphism

- operations on sets $A, B \subseteq \Sigma^*$

$$A \cup B \quad A \cap B \quad \sim A = \{x \in \Sigma^* \mid x \notin A\} \\ = \Sigma^* - A = \Sigma^* \setminus A$$

- $A \subseteq B \stackrel{\Delta}{\iff} \forall x \ x \in A \Rightarrow x \in B$

- $A = B \stackrel{\Delta}{\iff} \forall x \ x \in A \iff x \in B$

- $AB \stackrel{\Delta}{=} \{xy \mid x \in A \ \& \ y \in B\}$

set concatenation $\cdot : 2^{\Sigma^*} \times 2^{\Sigma^*} \rightarrow 2^{\Sigma^*}$

$$\mathcal{P}(\Sigma^*) \times \mathcal{P}(\Sigma^*) \rightarrow \mathcal{P}(\Sigma^*)$$

Ex: $\{a, ab\} \cdot \{b, ba\} = \{ab, aba, abb, abba\}$

- $\{\epsilon\}$ is an identity $\{\epsilon\} \cdot A = A \cdot \{\epsilon\} = A$

- $(AB)C = A(BC)$

- $(2^{\Sigma^*}, \cdot, \{\epsilon\})$ forms a monoid

- powers: $A^n = \underbrace{A \cdot A \cdot \dots \cdot A}_n$

$$A^0 \stackrel{\Delta}{=} \{\epsilon\} \quad A^{n+1} \stackrel{\Delta}{=} A^n \cdot A$$

- $A^* \stackrel{\Delta}{=} \bigcup_{n \geq 0} A^n = A^0 \cup A^1 \cup A^2 \cup \dots$

$$= \{x_1 x_2 \dots x_n \mid n \geq 0 \ \& \ x_i \in A, 1 \leq i \leq n\}$$

Kleene asterate

$$A \cup B = \{x \mid x \in A \text{ or } x \in B\}$$

$$A \cap B = \{x \mid x \in A \text{ and } x \in B\}$$

$$\{a, ab\} \cup \{ab, aab\} = \{a, ab, aab\}$$

$$\{a, ab\} \cap \{ab, aab\} = \{ab\}$$

$$\sim A = \{x \in \Sigma^* \mid x \notin A\} = \Sigma^* - A = \Sigma^* \setminus A$$

$$AB = \{xy \mid x \in A \text{ and } y \in B\}$$

$$\{a, ab\}\{b, ba\} = \{ab, aba, abb, abba\}$$

$$A^0 = \{\varepsilon\}$$

$$A^{n+1} = AA^n$$

$$\{ab, aab\}^0 = \{\varepsilon\}$$

$$\{ab, aab\}^1 = \{ab, aab\}$$

$$\{ab, aab\}^2 = \{abab, abaab, aabab, aabaab\}$$

$$\{ab, aab\}^3 = \{ababab, ababaab, abaabab, aababab, \\ abaabaab, aababaab, aabaabab, aabaabaab\}$$

$$\{a, b\}^n = \{x \in \{a, b\}^* \mid |x| = n\}$$

$$A^* = \bigcup_{n \geq 0} A^n = A^0 \cup A^1 \cup A^2 \cup A^3 \cup \dots$$

$$= \{x_1 x_2 \cdots x_n \mid n \geq 0 \text{ and } x_i \in A, 1 \leq i \leq n\}$$

$$A^+ = AA^* = \bigcup_{n \geq 1} A^n$$

Associativity:

$$(A \cup B) \cup C = A \cup (B \cup C) \quad (A \cap B) \cap C = A \cap (B \cap C) \quad (AB)C = A(BC)$$

Commutativity:

$$A \cup B = B \cup A$$

$$A \cap B = B \cap A$$

Identities:

$$A \cup \emptyset = \emptyset \cup A = A$$

$$\{\varepsilon\}A = A\{\varepsilon\} = A$$

Annihilation:

$$A\emptyset = \emptyset A = \emptyset$$

Distributivity:

$$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$$

$$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$$

$$A(B \cup C) = AB \cup AC$$

$$(A \cup B)C = AC \cup BC$$

$$A\left(\bigcup_{i \in I} B_i\right) = \bigcup_{i \in I} AB_i$$

$$\left(\bigcup_{i \in I} B_i\right)A = \bigcup_{i \in I} B_iA$$

De Morgan laws:

$$\sim(A \cup B) = \sim A \cap \sim B$$

$$\sim(A \cap B) = \sim A \cup \sim B$$

Some properties of *:

$$A^*A^* = A^*$$

$$(A^*)^* = A^*$$

$$A^* = \{\varepsilon\} \cup AA^* = \{\varepsilon\} \cup A^*A$$

$$\emptyset^* = \{\varepsilon\}$$

Deterministic finite automata (DFAs)

(singular: automaton, plural: automata)

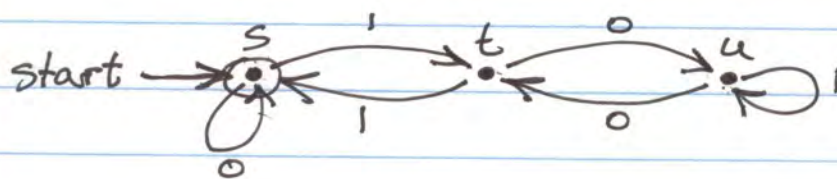
A DFA is a 5-tuple $M = (Q, \Sigma, \delta, s, F)$

~~where~~ where:

- Q is a finite set (the states)
- Σ is a finite set (the input alphabet)
- s is the start state, $s \in Q$
- F is the set of accept states or final states, $F \subseteq Q$
- $\delta: Q \times \Sigma \rightarrow Q$ is the transition function
 $\delta(q, a)$ is the state to move to next when in state q , seeing input symbol a .

To specify,

- write down components
- draw a diagram



$110101 \in \Sigma^*$

$1001 \in \Sigma^*$

$\Sigma = \{0, 1\}$ start state is s , denoted by \rightarrow

$F = \{s\}$, denoted by a circle

- table:

		0	1
s	s	t	s
t	t	u	s
u	u	t	u