

## Game Playing

Initial State is the initial board/position
Successor Function defines the set of legal moves from any position

Terminal Test determines when the game is over
Utility Function gives a numeric outcome for the game


## Simplified Minimax Algorithm

1. Expand the entire tree below the root.
2. Evaluate the terminal nodes as wins for the minimizer or maximizer (i.e. utility).
3. Select an unlabeled node, $n$, all of whose children have been assigned values. If there is no such node, we're done --return the value assigned to the root.
4. If $n$ is a minimizer move, assign it a value that is the minimum of the values of its children. If $n$ is a maximizer move, assign it a value that is the maximum of the values of its children. Return to Step 3.


Improving Minimax: $\alpha-\beta$ Pruning

Idea: Avoid generating the whole search tree
Approach: Analyze which subtrees have no influence on the solution

$$
\alpha-\beta \text { Search }
$$

$\boldsymbol{\alpha}=$ lower bound on Max's outcome; initially set to $-\infty$ $\boldsymbol{\beta}=$ upper bound on Min's outcome ; initially set to $+\infty$

We'll call $\alpha-\beta$ procedure recursively with a narrowing range between $\alpha$ and $\boldsymbol{\beta}$.

Maximizing levels may reset $\alpha$ to a higher value;
Minimizing levels may reset $\boldsymbol{\beta}$ to a lower value.

## Minimax

function MINIMAX-DECISION(game) returns an operator
for each $o p$ in OPERATORS[game]do
VALUE $[o p] \leftarrow$ MINIMAX-VALUE(APPLY(op,game),game)
end
return the $o p$ with the highest VALUE[op]
function MINIMAX-VALUE(state,game) returns a utility value
if TERMINAL-TEST[game](state) then
return UTILITY[game](state)
else if MAX is to move in state then
return the highest MINIMAX-VALUE of SUCCESSORS(state) else
return the lowest MINIMAX-VALUE of SUCCESSORS(state)
$\alpha-\beta$ Search Algorithm

1. If terminal state, compute $e(n)$ and return the result.
2. Otherwise, if the level is a minimizing level,

- Until no more children or $\beta \leq \alpha$
- $v_{i} \leftarrow \alpha-\boldsymbol{\beta}$ search on a child
- If $v_{i}<\beta, \beta \leftarrow v_{i}$.
- Return $\min \left(v_{i}\right)$

3. Otherwise, the level is a maximizing level:

- Until no more children or $\alpha \geq \beta$,
$-v_{i} \leftarrow \alpha-\boldsymbol{\beta}$ search on a child.
- If $v_{i}>\alpha$, set $\alpha \leftarrow v_{i}$
- Return $\max \left(v_{i}\right)$


## Search Space Size Reductions

Worst Case: In an ordering where worst options evaluated first, all nodes must be examined.

Best Case: If nodes ordered so that the best options are evaluated first, then what?


## Static Evaluation Functions

Minimax depends on the translation of board quality into single, summarizing number. Difficult. Expensive.

- Add up values of pieces each player has (weighted by importance of piece).
- Isolated pawns are bad.
- How well protected is your king?
- How much maneuverability to you have?
- Do you control the center of the board?
- Strategies change as the game proceeds.


## The Need for Imperfect Decisions

Problem: Minimax assumes the program has time to search to the terminal nodes.

Solution: Cut off search earlier and apply a heuristic evaluation function to the leaves.

Design Issues for Heuristic Minimax

## Evaluation Function:

Need to be carefully crafted and depends on game! What criteria should an evaluation function fulfill?


Design Issues for Heuristics Minimax

Search: search to a constant depth

What are problems with constant search depth?


| Backgammon - Rules |
| :--- |
| - If you roll doubles you take 4 moves |
| (example: roll 5,5 , make moves $5,5,5,5$ ). |
| - Moves can be made by one or two pieces |
| (in the case of doubles by $1,2,3$ or 4 pieces) |
| - And a few other rules that concern bearing off and forced |
| moves. |


| Backgammon - Rules |
| :--- |
| - If you roll doubles you take 4 moves |
| (example: roll 5,5, make moves $5,5,5,5$ ). |
| - Moves can be made by one or two pieces |
| (in the case of doubles by $1,2,3$ or 4 pieces) |
| - And a few other rules that concern bearing off and forced |
| moves. |


| Backgammon - Rules |
| :--- |
| - If you roll doubles you take 4 moves |
| (example: roll 5,5, make moves $5,5,5,5$ ). |
| - Moves can be made by one or two pieces |
| (in the case of doubles by $1,2,3$ or 4 pieces) |
| - And a few other rules that concern bearing off and forced |
| moves. |


| Backgammon - Rules |
| :--- |
| - If you roll doubles you take 4 moves |
| (example: roll 5,5, make moves $5,5,5,5$ ). |
| - Moves can be made by one or two pieces |
| (in the case of doubles by $1,2,3$ or 4 pieces) |
| - And a few other rules that concern bearing off and forced |
| moves. |


| Backgammon - Rules |
| :--- |
| - If you roll doubles you take 4 moves |
| (example: roll 5,5, make moves $5,5,5,5$ ). |
| - Moves can be made by one or two pieces |
| (in the case of doubles by $1,2,3$ or 4 pieces) |
| - And a few other rules that concern bearing off and forced |
| moves. |


| Backgammon - Rules |
| :--- |
| - If you roll doubles you take 4 moves |
| (example: roll 5,5, make moves $5,5,5,5$ ). |
| - Moves can be made by one or two pieces |
| (in the case of doubles by $1,2,3$ or 4 pieces) |
| - And a few other rules that concern bearing off and forced |
| moves. |


| Backgammon - Rules |
| :--- |
| - If you roll doubles you take 4 moves |
| (example: roll 5,5, make moves $5,5,5,5$ ). |
| - Moves can be made by one or two pieces |
| (in the case of doubles by $1,2,3$ or 4 pieces) |
| - And a few other rules that concern bearing off and forced |
| moves. |

## Backgammon - Rules

- Goal: move all of your pieces off the board before your opponent does.
- Black moves counterclockwise toward 0 .
- White moves clockwise toward 25.
- A piece can move to any position except one where there are two or more of the opponent's pieces.
- If it moves to a position with one opponent piece, that piece is captured and has to start it's journey from the beginning.

Game Tree for Backgammon


## Expectiminimax

$\operatorname{Expectiminimax}(\mathrm{n})=$
Utility(n)
for $n$, a terminal state
$\max _{s \in \operatorname{Succ}(n)} \operatorname{expectiminimax}(\boldsymbol{s})$ for $\mathbf{n}$, a Max node
$\min _{s \in \operatorname{Succ}(n)} \operatorname{expectiminimax}(\boldsymbol{s}) \quad$ for $\mathbf{n}$, a Min node
$\Sigma_{\boldsymbol{s} \in \operatorname{Succ}(\boldsymbol{n})} P(\boldsymbol{s}) * \operatorname{expectiminimax}(\boldsymbol{s})$ for $\mathbf{n}$, a chance node


## State of the Art in Checkers

- 1952: Samuel developed a checkers program that learned its own evaluation function through self play.
- 1990: Chinook (J. Schaeffer) wins the U.S. Open. At the world championship, Marion Tinsley beat Chinook.
- 2005: Schaeffer et al. solved checkers for "White Doctor" opening (draw) (about 50 other openings).


## History of Chess in AI

| 500 | Legal chess |
| :--- | :--- |
| 1200 | Occasional player |
| 2000 | World-ranked |
| 2900 | Gary Kasparov |

Early 1950's Shannon and Turing both had programs that (barely) played legal chess ( 500 rank ).

1950's Alex Bernstein's system, $(500+\varepsilon)$
1957 Herb Simon claims that a computer chess program would be world chess champion in 10 years...yeah, right.

## State of the Art in Backgammon

- 1980: $B K G$ using two-ply (depth 2) search and lots of luck defeated the human world champion.
- 1992: Tesauro combines Samuel's learning method with neural networks to develop a new evaluation function (search depth 2-3), resulting in a program ranked among the top 3 players in the world.


## State of the Art in Go

Large branching factor makes regular search methods inappropriate.

Best computer Go programs ranked only "weak amateur".

Employ pattern recognition techniques and limited search.
$\$ 2,000,000$ prize available for first computer program to defeat a top level player.

1966 McCarthy arranges computer chess match, Stanford vs. Russia. Long, drawn-out match. Russia wins.

1967 Richard Greenblatt, MIT. First of the modern chess programs, MacHack (1100 rating).

1968 McCarthy, Michie, Papert bet Levy (rated 2325) that a computer program would beat him within 10 years.

1970 ACM started running chess tournaments. Chess 3.0-6 (rated 1400).
1973 By 1973...Slate: "It had become too painful even to look at Chess 3.6 any more, let alone work on it."

1973 Chess 4.0: smart plausible-move generator rather than speeding up the search. Improved rapidly when put on faster machines.

```
1976 Chess 4.5: ranking of 2070.
1977 Chess 4.5 vs.~Levy.Levy wins.
1980's Programs depend on search speed rather than knowledge (2300 range).
1993 DEEP THOUGHT: Sophisticated special-purpose computer; }\alpha-\boldsymbol{\beta
    search; searches 10-ply; singular extensions; rated about 2600.
1995 DEEP BLUE: searches 14-ply; iterative deepening }\boldsymbol{\alpha}-\boldsymbol{\beta}\mathrm{ search;
    considers 100-200 billion positions per move; regularly reaches depth 14;
    considers 100-200 billion positions per move; regularly reaches depth 
    opening book of 4000 positions; end-game database for 5-6 pieces.
1997 DEEP BLUE: first match won against world-champion (Kasparov). }200
    IBM declines re-match. FRITZ played world champion Vladimir Kramnik. }
    games. Ended in a draw.
```

Informed Search: use heuristic function guide to goal Greedy best-first search
A* search / provably optimal
Search space up to approximately $10^{25}$

## Local search

Greedy / Hillclimbing
Simulated annealing
Tabu search
Genetic Algorithms / Genetic Programming
search space $10^{100}$ to $10^{1000}$
Aversarial Search / Game Playing
minimax Up to $\sim 10^{10}$ nodes, 6-7 ply in chess.
alpha-beta pruning Up to $\sim 10^{20}$ nodes, 14 ply in chess. provably optimal

## Search and AI

## Why such a central role?

- Basically, because lots of tasks in AI are intractable. Search is "only" way to handle them.
- Many applications of search, in e.g., Learning / Reasoning / Planning / NLU / Vision
- Good thing: much recent progress ( $10^{30}$ quite feasible; sometimes up to $10^{1000}$ ).

Qualitative difference from only a few years ago!

