Foundations of Artificial Intelligence

Constraint Satisfaction

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Moving to a different formalism...

SEND + More Money

Consider state space for cryptarithmetic (e.g. DFS).

Is this (DFS) how humans tackle the problem?

Human problem solving appears more sophisticated! For example, we derive new constraints on the fly. \rightarrow little or no search!

Constraint Satisfaction Problems (CSP)

A powerful representation for (discrete) search problems

A Constraint Satisfaction Problem (CSP) is defined by: **X** is a set of n variables $X_1, X_2, ..., X_n$ each defined by a finite domain $D_1, D_2, ..., D_n$ of possible values.

 ${\bf C}$ is a set of constraints $C_1, C_2, \ldots, C_m.$ Each C_i involves a subset of the variables; specifies the allowable combinations of values for that subset.

A solution is an assignment of values to the variables that satisfies all constraints.

Cryptarithmetic as a CSP

Problem: TWO + TWO = FOUR Variables: $T \in \{0,...,9\}; W \in \{0,...,9\}; 0 \in \{0,...,9\};$ $F \in \{0,...,9\}; U \in \{0,...,9\}; R \in \{0,...,9\};$ $X_1 \in \{0,...,9\}; X_2 \in \{0,...,9\}; X_3 \in \{0,...,9\};$ Constraints: $0 + 0 = R + 10 * X_1$ $X_1 + W + W = U + 10 * X_2$ $X_2 + T + T = O + 10 * X_3$ $X_3 = F$ each letter has a different digit (F \neq T, F \neq U, etc.);





Types of Constraints

Unary Constraints: Restriction on single variable

Binary Constraints: Restriction on pairs of variables

Higher-Order Constraints: Restriction on more than two variables

Constraint Satisfaction Problems (CSP)

For a given CSP the problem is one of the following:

- 1. find all solutions
- 2. find one solution
 - · just a feasible solution, or
 - · A "reasonably good" feasible solution, or
 - \cdot the optimal solution given an objective function
- 3. determine if a solution exists

How to View a CSP as a Search Problem?

Initial State - state in which all the variables are unassigned.

Successor function - assign a value to a variable from a set of possible values.

Goal test - check if all the variables are assigned and all the constraints are satisfied.

Path cost - assumes constant cost for each step

Branching Factor

Approach 1- any unassigned variable at a given state can be assigned a value by an operator: branching factor as high as sum of size of all domains.

Approach 2 - since order of variable assignment not relevant, consider as the successors of a node just the different values of a *single* unassigned variable: max branching factor = max size of domain.

Maximum Depth of Search Tree n the number of variables \rightarrow all solutions at depth n.

Prefer DFS or BFS?

CSP – Goal Decomposed into Constraints

Backtracking Search: a DFS that

- · chooses values for variables one at a time
- checks for *consistency* with the constraints.

Decisions during search:

- · Which variable to choose next for assignment.
- · Which value to choose next for the variable.

Forward Checking

- **Idea:** Reduce domain of unassigned variables based on assigned variables.
- Each time variable is instantiated, delete from domains of the uninstantiated variables all of those values that conflict with current variable assignment.
- Identify dead ends without having to try them via backtracking.

General Purpose Heuristics

Variable and value ordering:

Minimum remaining values (MRV): choose the variable with the fewest possible values.

Degree heuristic: assign a value to the variable that is involved in the largest number of constraints on other unassigned variables.

Least-constraining value heuristic: choose a value that rules out the smallest number of values in variables connected to the current variable by constraints.

Comparison of CSP Algorithms

Problem	BT	BT+MRV	BT+FC	BT+FC+MRV
USA	(>1,000K)	(>1,000K)	2K	60
N-queens	(>40,000K)	13,500K	(>40,000K)	817K
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Constraint Propagation (Arc Consistency)

Arc Consistency - state is arc-consistent, if every variable has some value that is consistent with each of its constraints (consider pairs of variables)

- Init: Q is queue with all (directed) arcs (X_i, X_i) in CSP
 - WHILE Q is not empty
 - (X_i, X_j) = remove_first(Q) FOREACH $x \in dom(X_i)$

 - *IF no $y \in dom(Xj)$ satisfies constraint (X_i, X_i)
 - •THEN remove x from $dom(X_i)$
 - IF *dom*(X_i) changed

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*THEN add all arcs $(X_k, X_i) \notin Q$ to Q

Constraint Propagation (K-Consistency)

- K-Consistency generalizes arc-consistency (2-consistency).
- Consistency of groups of K variables.



Remarks

- · Infinite discrete domains and continuous domains
- · Exploiting special problem structure

• Dramatic recent progress in Constraint Satisfaction. Methods can now handle problems with 10,000 to 100,000 variables, and up to 1,000,000 constraints.