

## 4410 Lecture 8: monitors

- (finish semaphore impl)
- Monitor design pattern
- (quiz)

## Using semaphores for thread control (per-thread sema)

To sleep current thread:

allocate new semaphore  $s$

put  $s$  in a data structure

$P(s)$

↪ decrement

To wake up sleeping thread:

get  $s$  from data structure

$V(s)$

← increment

delete  $s$

To handle a syscall (monolithic kernel)

① save state as if executing a normal function call

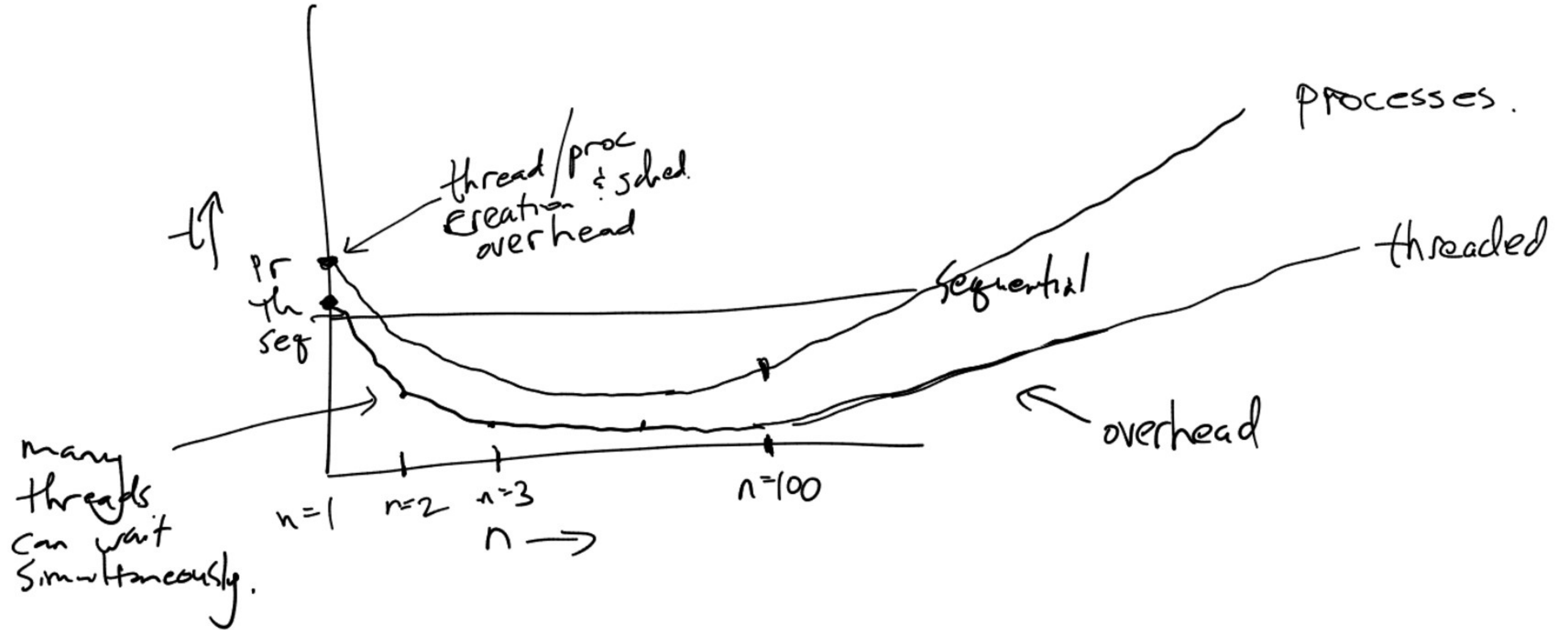
② perform the operation (e.g. communicate w/ device)

- branch (normal jump) to driver code to interact w/ device.

(note: if microkernel, schedule driver process)

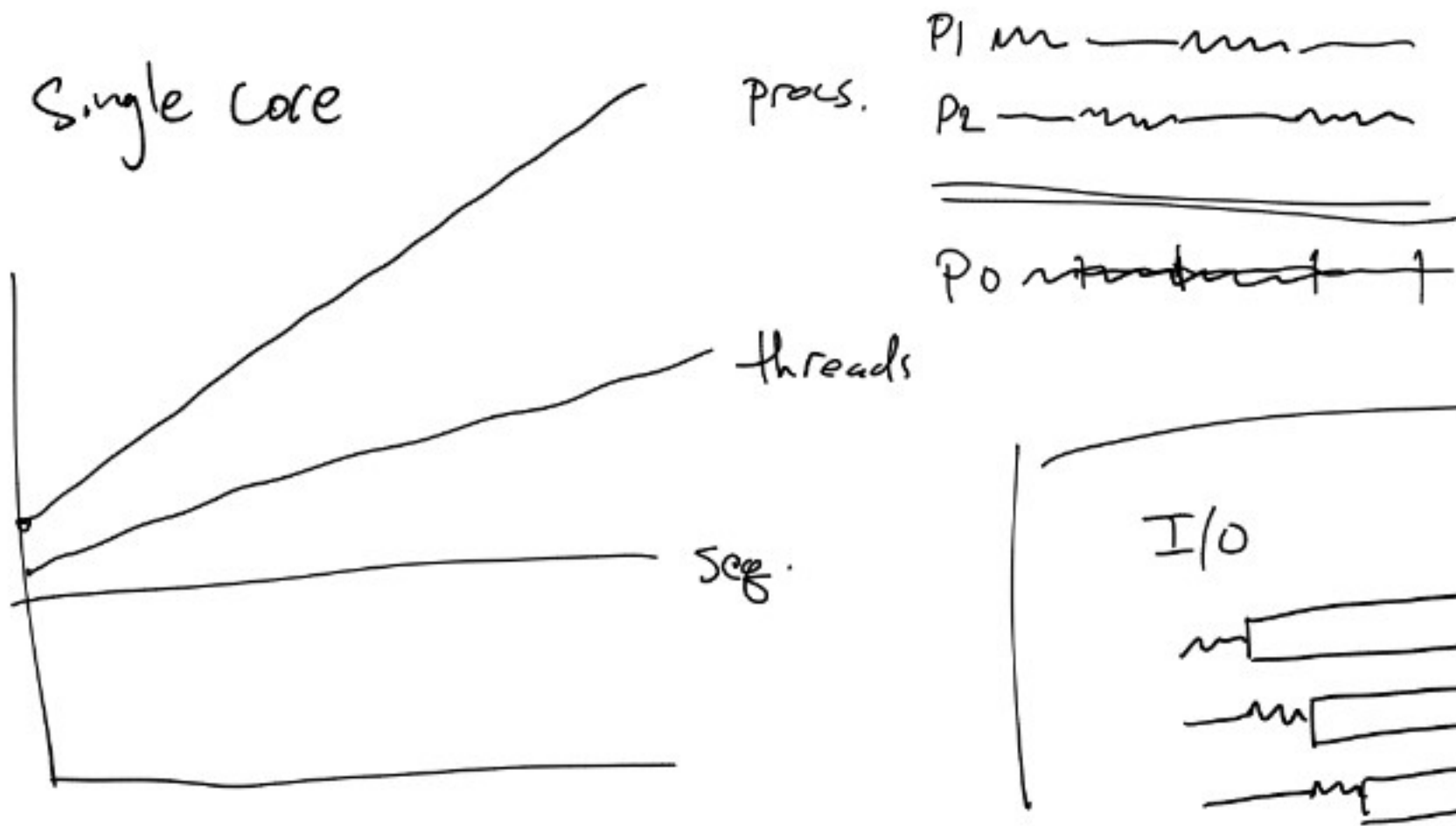
③ return from syscall (restoring state)

# I/O bound



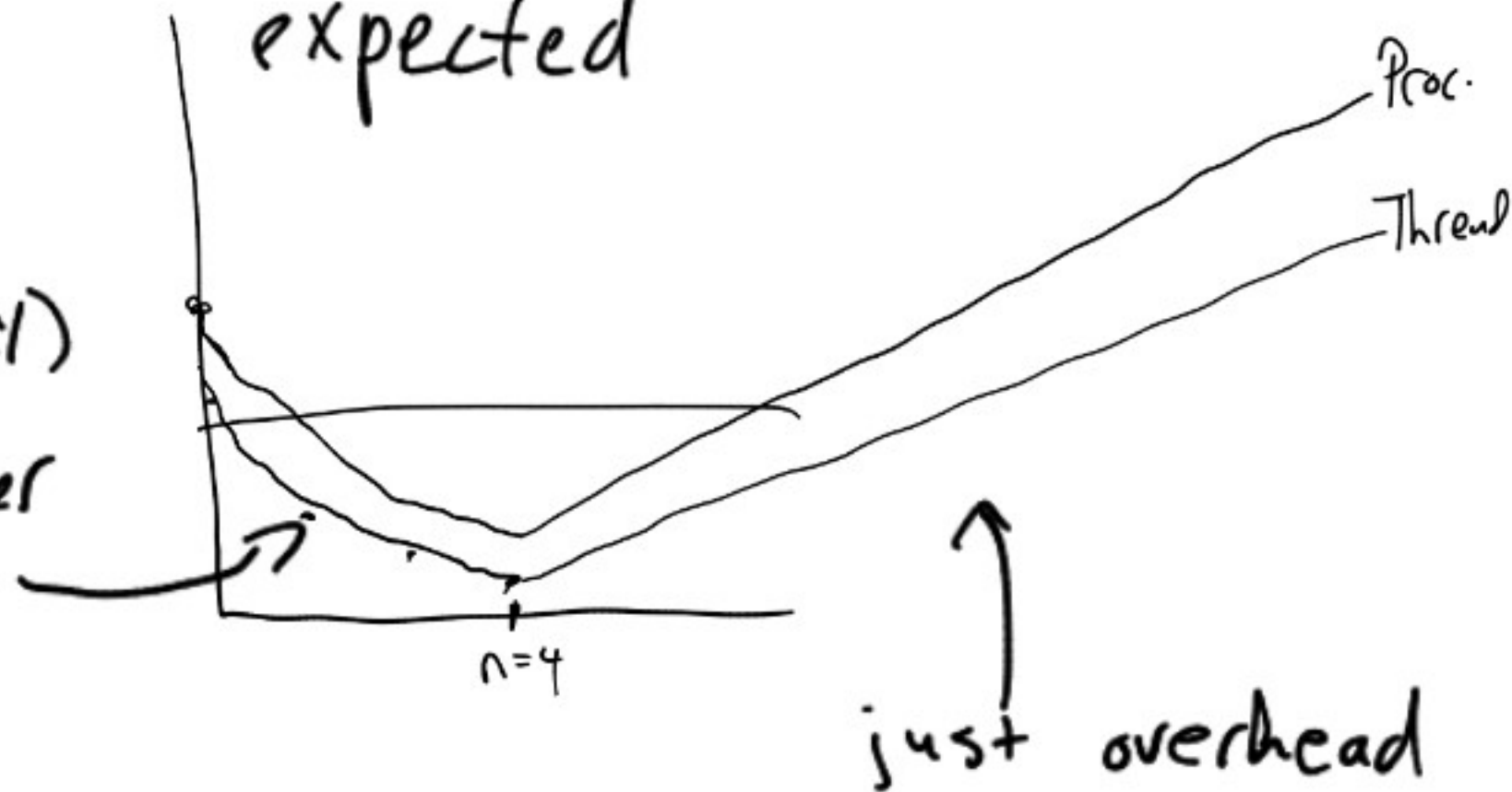
# CPU bound threads/procs

Only one CPU, so just overhead

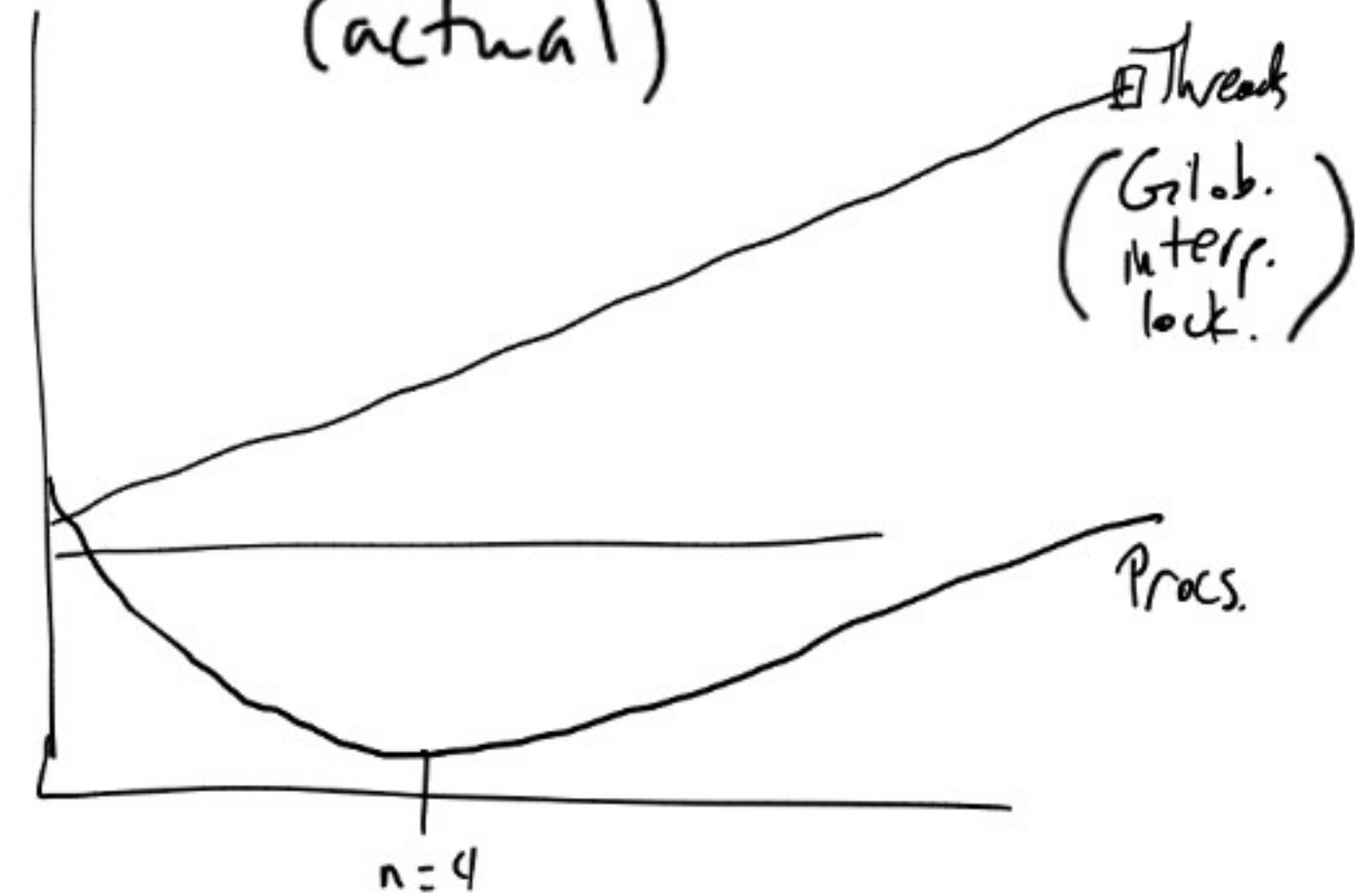


more threads ⇒ more CPUs (if avail) ⇒ faster

multicore (4) expected



multicore (actual)



class Semaphore (object):

def \_\_init\_\_(self, val):

self.val = val # protected by self.lock

self.lock = 0 # TAS lock

self.queue = [] # protected by self.lock  
(waiting)

# invariant:  $val \geq 0$ ,  $val = 0$  if  $q$  is nonempty

def P(self):

while TAS(self.lock):

yield() # or pass

if self.val > 0:

self.val --

self.lock = 0

return

else:

self.queue.enqueue(running-thread T<sub>0</sub>)

release lock: deschedule()  
current thread

return

def V(self):

while TAS(lock):

yield() # or pass

if queue.empty():

self.val ++

lock = 0

return

else:

dequeue a thread t,  
mark t as ready

lock = 0

return.

(scheduler)

## Monitor design pattern

- Invented by sir Tony Hoare. (semaphores invented by Dijkstra)
- closely related to Hoare logic
- Get support from PL.

Monitor object is like a normal OOP object,

• state (member vars) • methods

• at any time, only one thread running inside a given monitor.

• able to wait for certain conditions to hold  
(e.x. wait for count > 0)

# Semaphore impl. as a monitor (Hoare style)

```
class Semaphore (Monitor):  
    def __init__(value): ←  
        self.val = value #invariant: value ≥ 0  
  
    def P(self):  
        wait until self.value > 0 ←  
        self.val --  
  
    def V(self):  
        self.val ++
```

Hoare-style.  
In practice,  
need a little  
more effort  
from programmer.



## Bounded buffer monitor (Hoare style)

```
class Buffer (Monitor):  
    def __init__(self):  
        self.in = 0          # next empty slot  
        self.out = 0        # next full slot (if  $\neq$  in)  
        self.buf = Object[N]  
        # invariant:  $0 \leq \text{out} < N$   
        #  $\text{in} = \text{out}$  represents empty queue,  
        #  $\text{in} + 1 = \text{out}$  represents full queue
```

```
    def put(self, obj):  
        wait until  $\text{in} + 1 \neq \text{out} \pmod{N}$   
        buf [ $\text{in} + 1 \pmod{N}$ ] = obj
```

```
    def get(self, obj):  
        wait until  $\text{in} \neq \text{out}$   
        return buf [ $\text{out} \pmod{N}$ ]
```