## Networking — Network layer

### Three concepts

- Naming
  - A way to identify the source/destination
  - E.g., house address

#### Routing

- Finding "how to" move towards the destination
- E.g., which airplane should the stuff go on

#### Forwarding

- Actually "moving" towards the destination
- E.g., Using airplane/truck/rail

Lets come up with an approach? Generalize Ethernet ideas?

**Attempt 1: Broadcast** 

Send to everybody

- Send to everybody
- Goods

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  - Oh, well, everything else

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- Not-so-goods
  - Oh, well, everything else
  - Bandwidth overheads

### **Attempt 2: Time division Multiplexing**

Each source-destination pair assigned a time slot

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  - Can send data only during that slot

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- Goods
  - No collisions

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  - Can send data only during that slot
- Goods
  - No collisions
- Not-so-goods

- Each source-destination pair assigned a time slot
  - Can send data only during that slot
- Goods
  - No collisions
- Not-so-goods
  - Underutilization of resources

### **Attempt 3: Frequency division Multiplexing**

• Each source-destination pair assigned a subset of resources

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  - Can use only "assigned" resources (e.g., bandwidth)

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  - Predictable performance

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**Attempt 2 and 3: Circuit Switching** 

Source establishes connection

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  - Resources along the path are reserved

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  - Resources along the path are reserved
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- Source tears down connection
  - Free resources for others to use

#### **Circuit Switching**

• Goods:

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  - Predictable performance

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  - Reliable delivery

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  - Blocked connections

#### **Circuit Switching**

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- Reliable delivery
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#### Not-so-goods

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- Blocked connections
- Connection set up overheads

#### **Circuit Switching**

#### Goods:

- Predictable performance
- Reliable delivery
- Simple forwarding mechanism

#### Not-so-goods

- Resource underutilization
- Blocked connections
- Connection set up overheads
- Per-connection state in switches (scalability problem)

#### **Attempt 4: Packet Switching**

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- Each device stores a "look-up table"
  - Whats the next hop towards the destination?
- Destination receives the packet(s)
  - And reconstructs the message

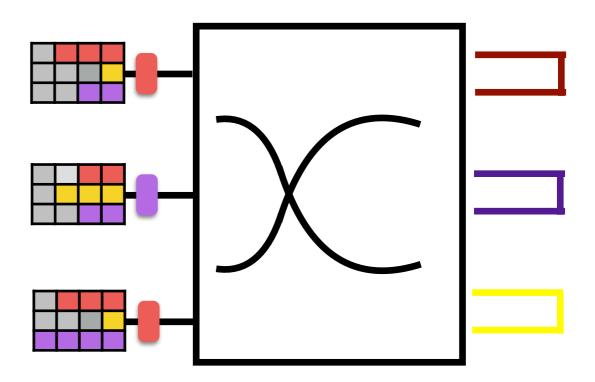
#### **Packet Switched forwarding**

Hop-by-hop forwarding

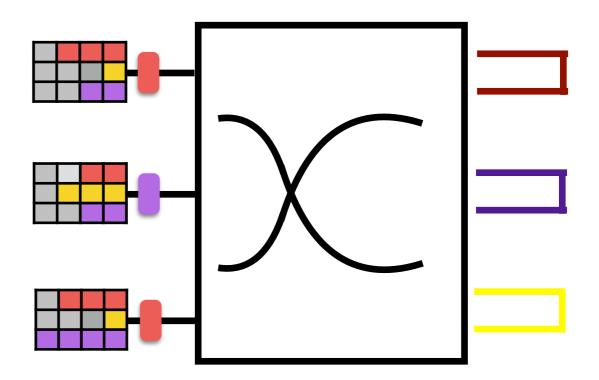
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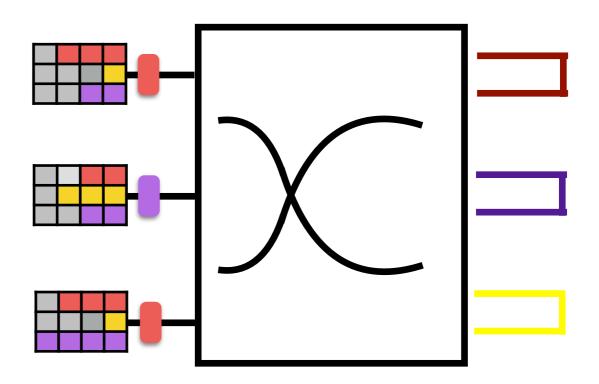
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- Each router has a "look-up table" (forwarding information base)
  - What should be stored in this table?
  - Prefix-based forwarding (longest-prefix matching)
    - Maps prefixes to the next-hop



#### **Packet Switching**

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### Network layer — Forwarding

#### **Packet Switching**

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- Not-so-goods:
  - Packet header overhead

### Network layer — Forwarding

#### **Packet Switching**

#### Goods:

- No resource underutilization
  - A source can send more if others don't use resources
- No blocked connection problem
- No per-connection state
- No set-up cost

#### Not-so-goods:

- Packet header overhead
- Network failures become a problem

### Networking — Network layer

#### Three concepts

#### Naming

- A way to identify the source/destination
- E.g., house address

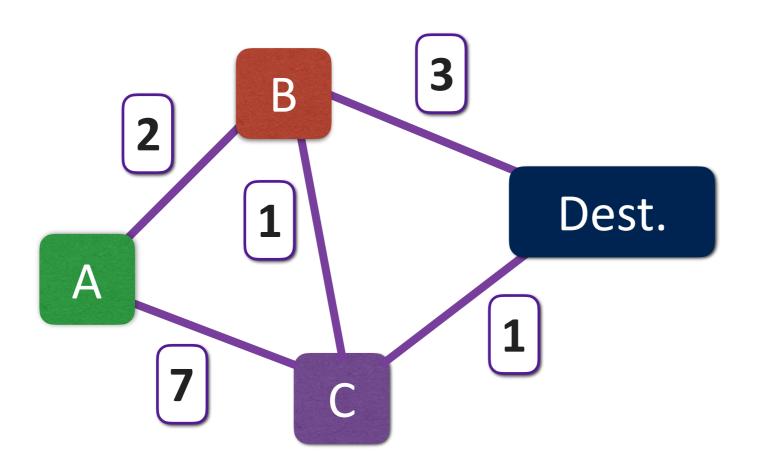
#### Routing

- Finding "how to" move towards the destination
- E.g., which airplane should the stuff go on

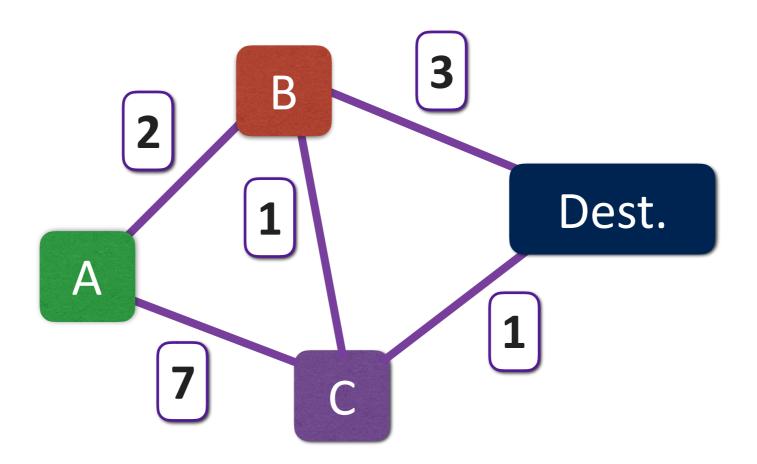
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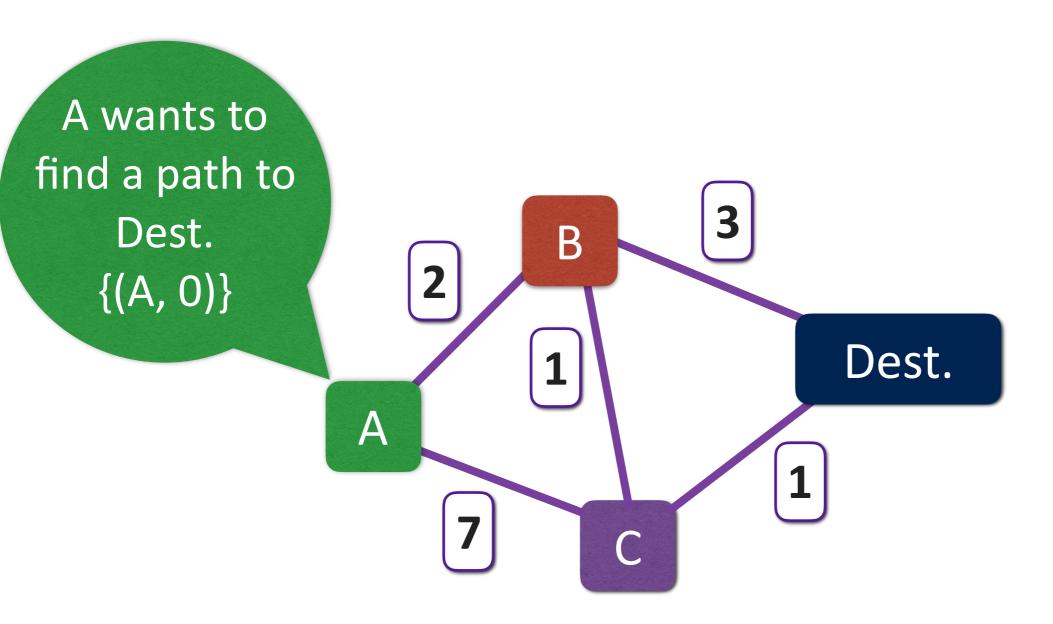
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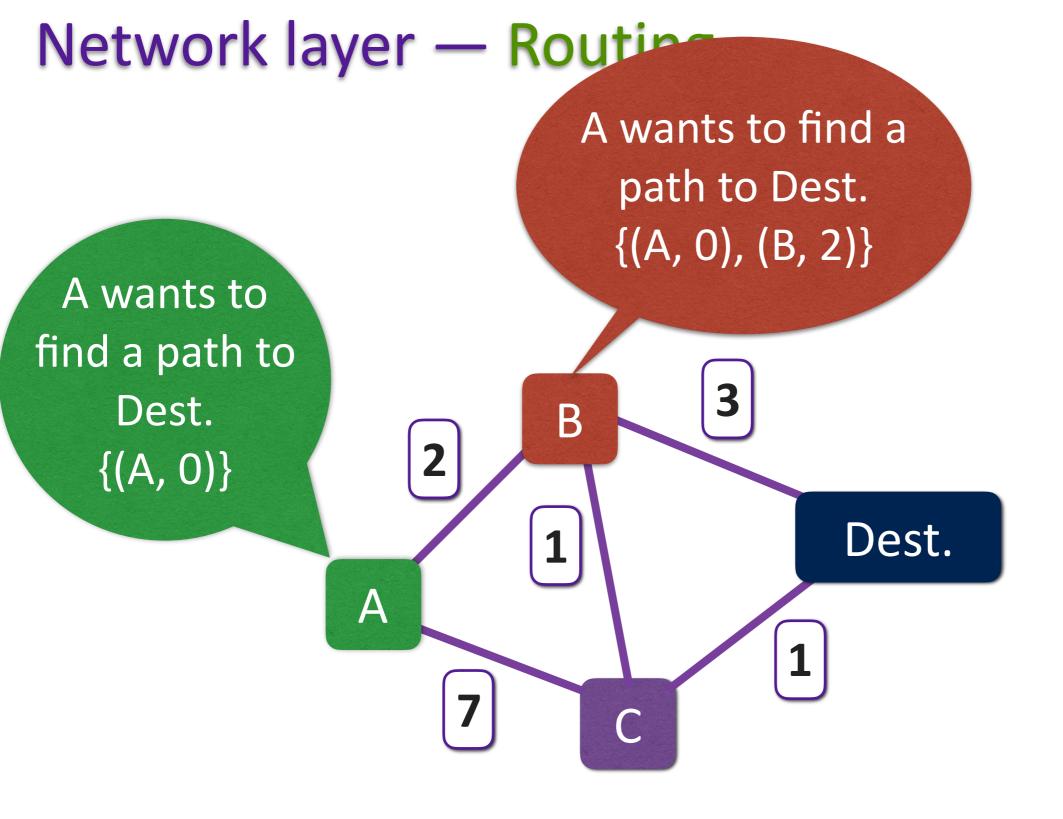
## Network layer — Example

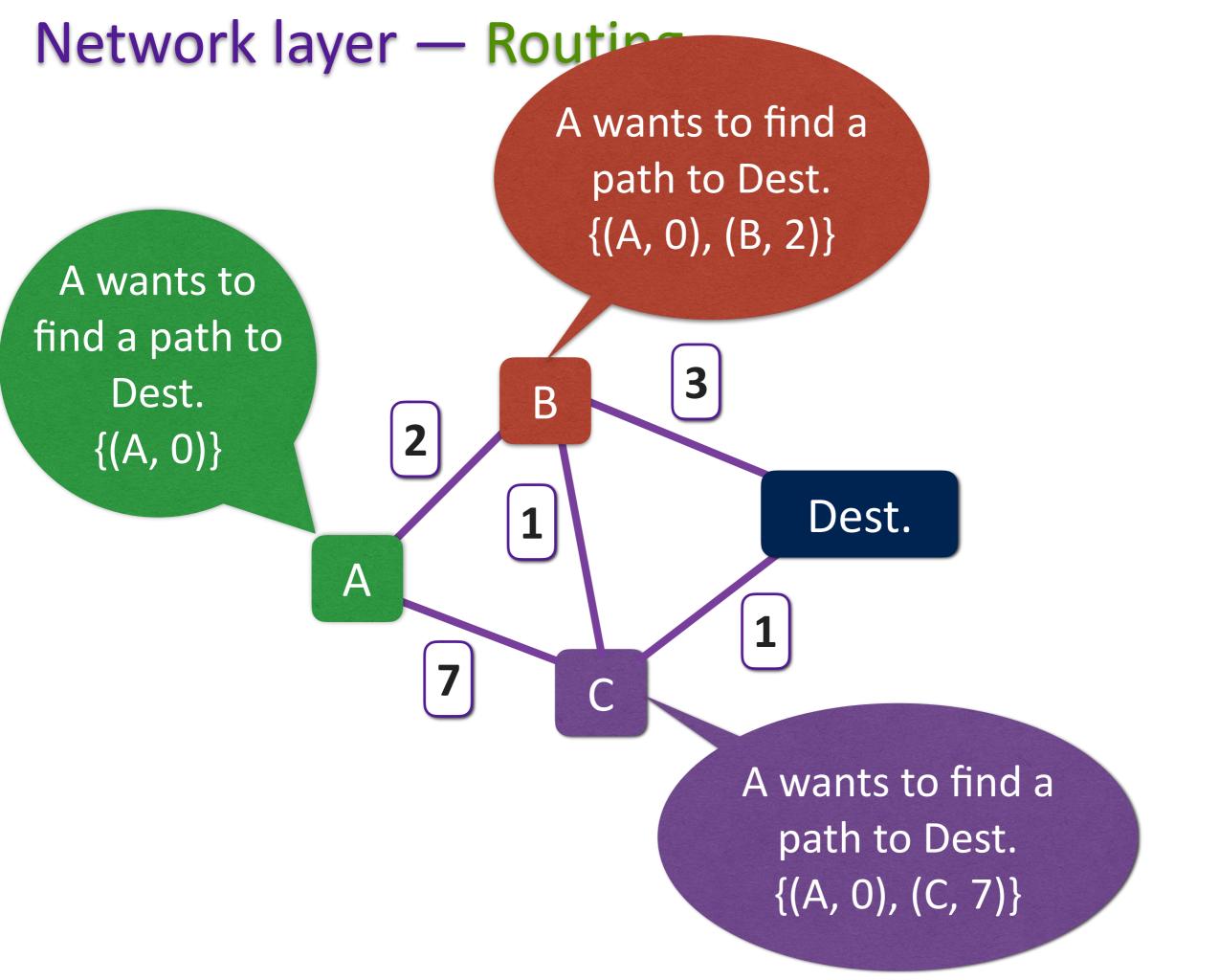


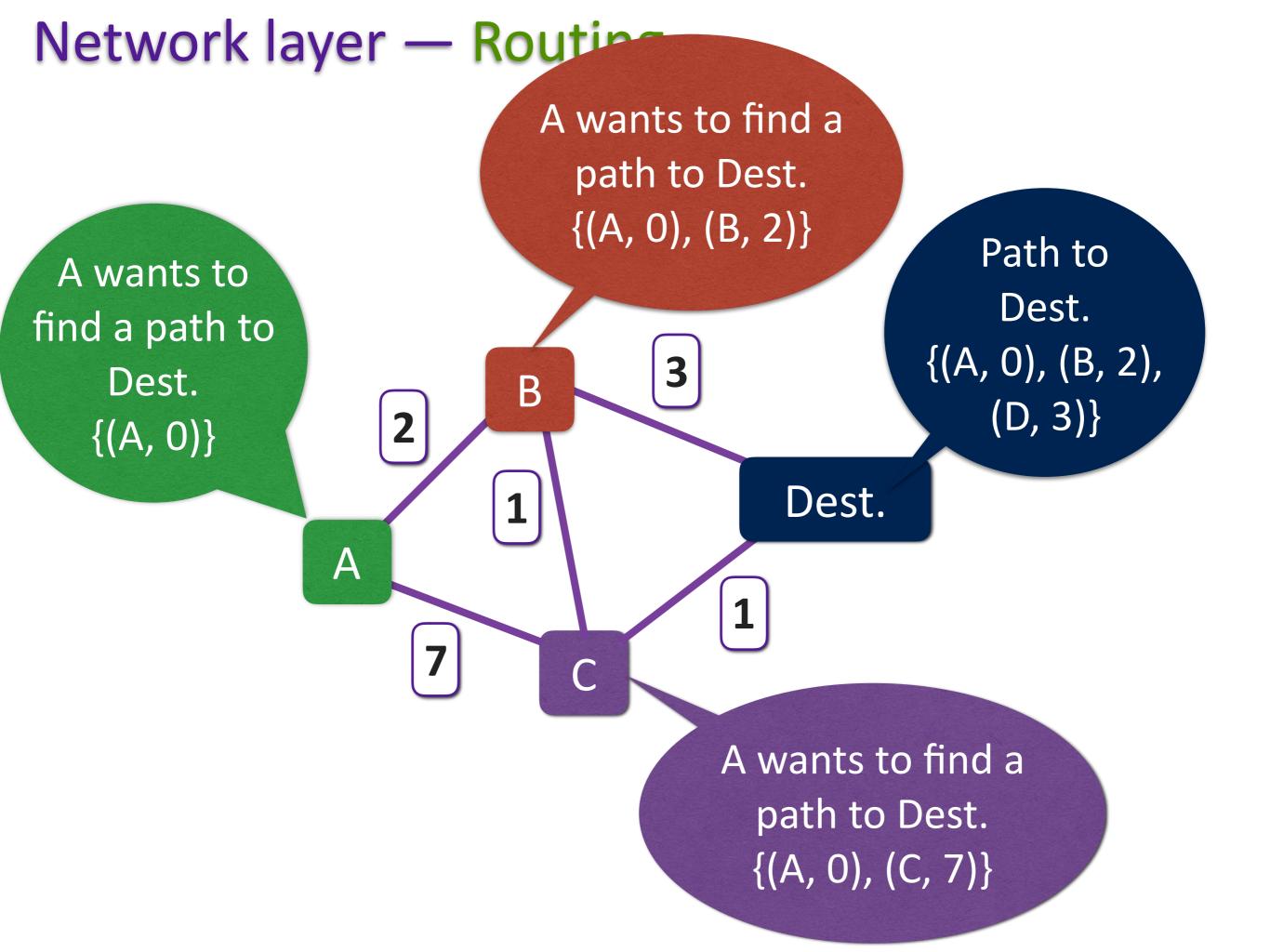
Lets come up with a routing scheme

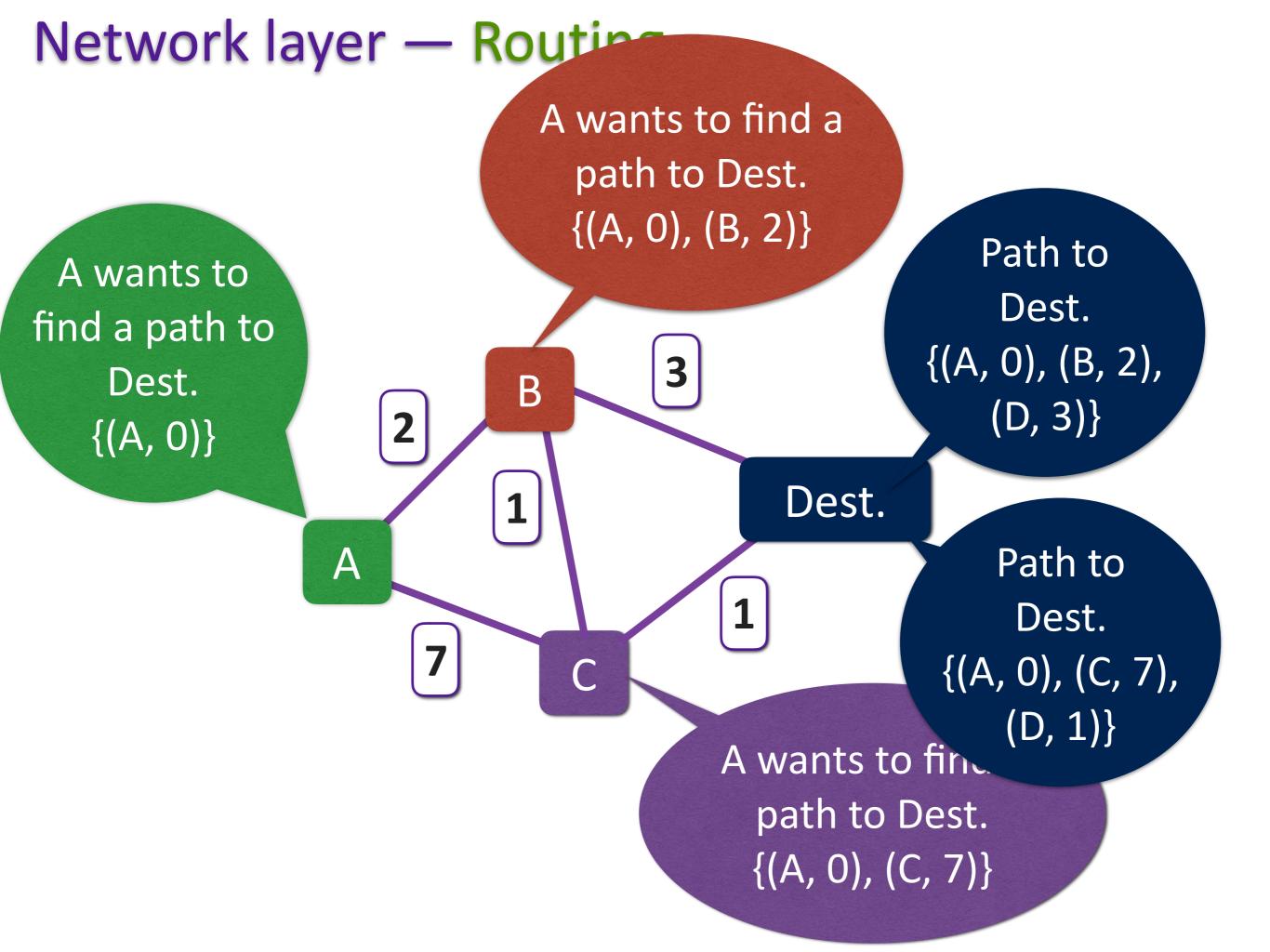


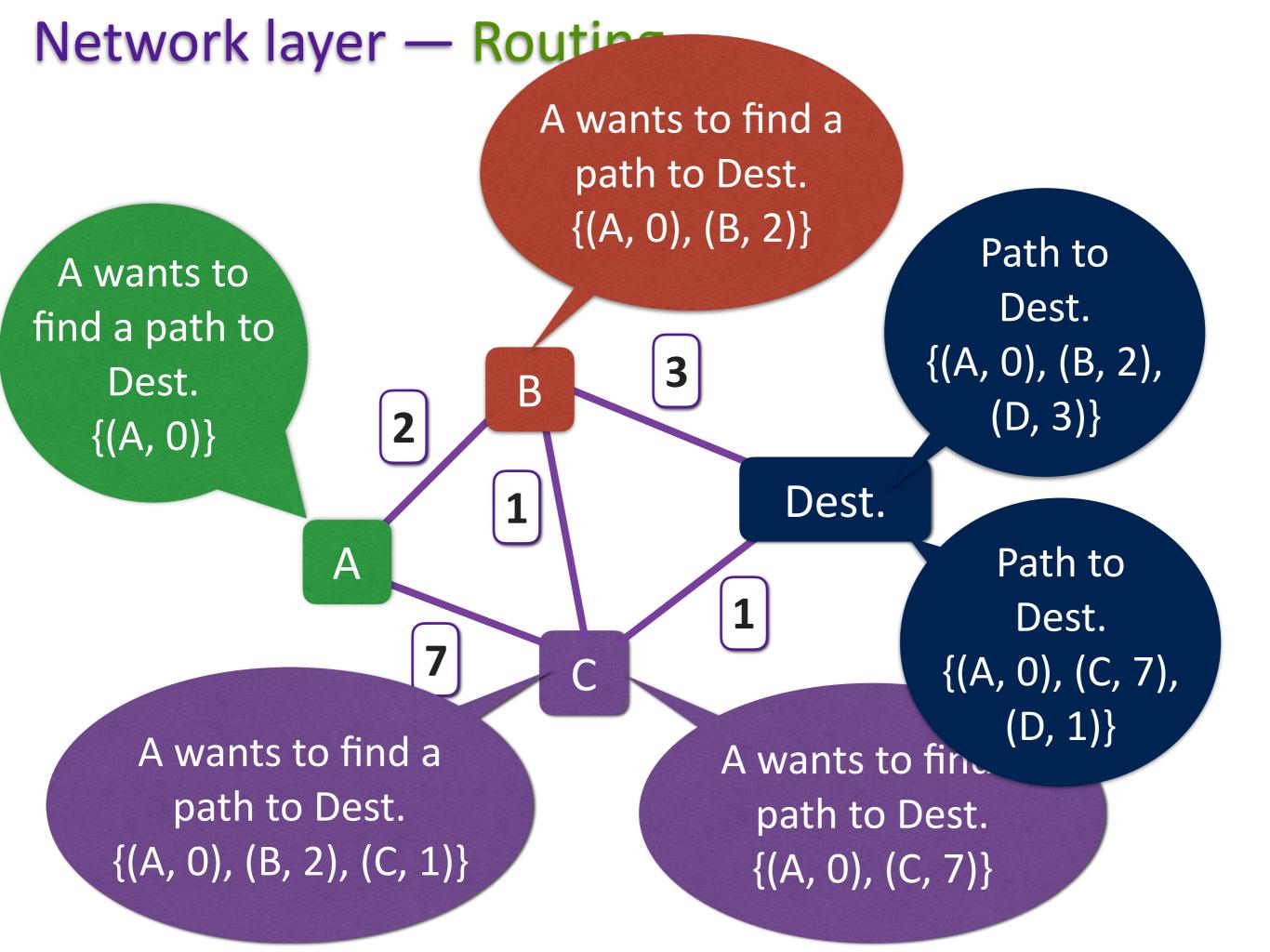


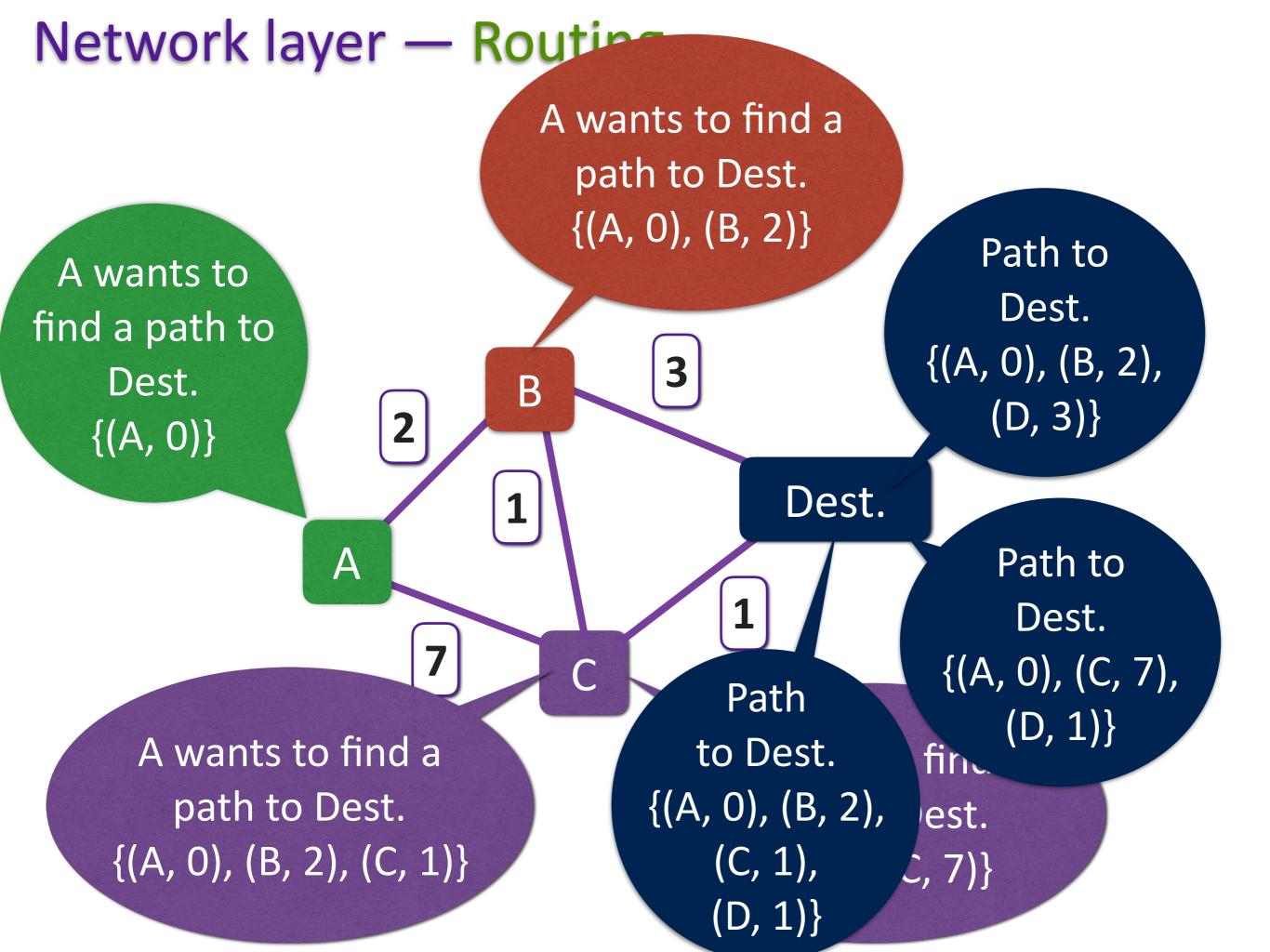












#### **Attempt 1: Dynamic Source Routing**

Broadcast a Route Request Packet for destination d

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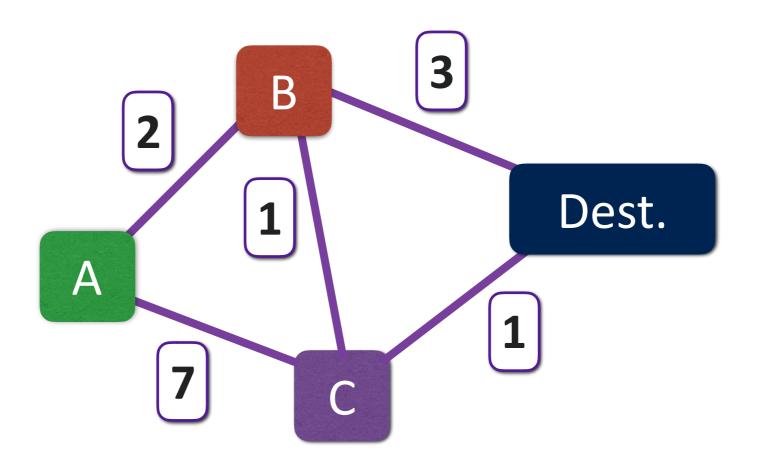
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    - Broadcast the Route Request Packet

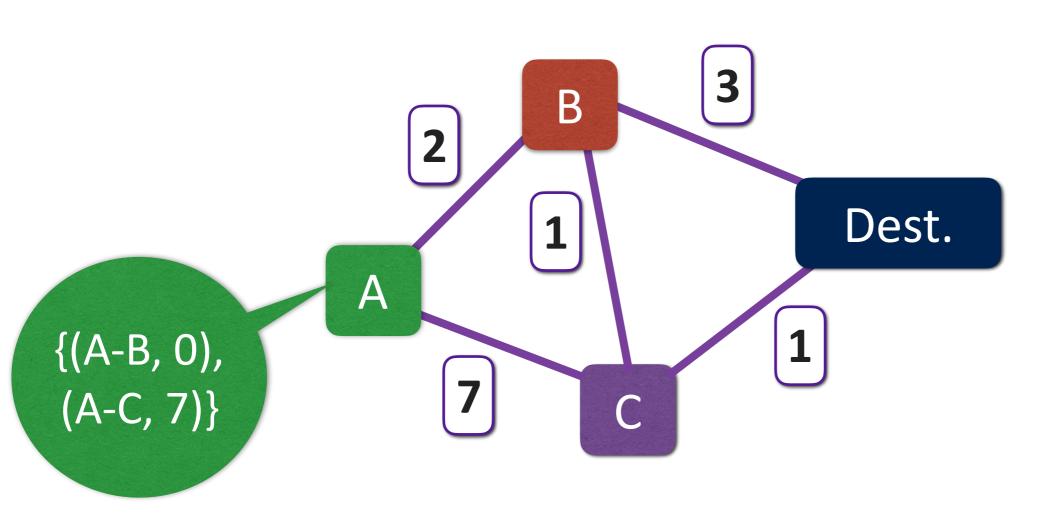
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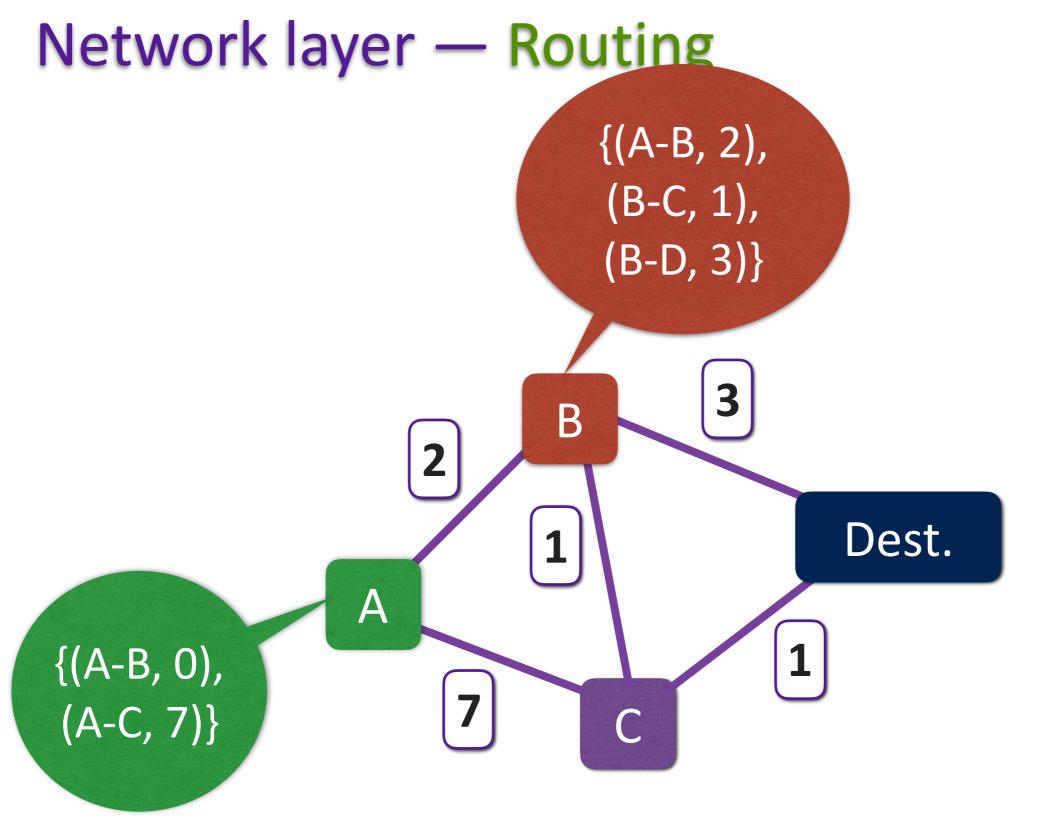
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    - Respond with a Route Reply packet

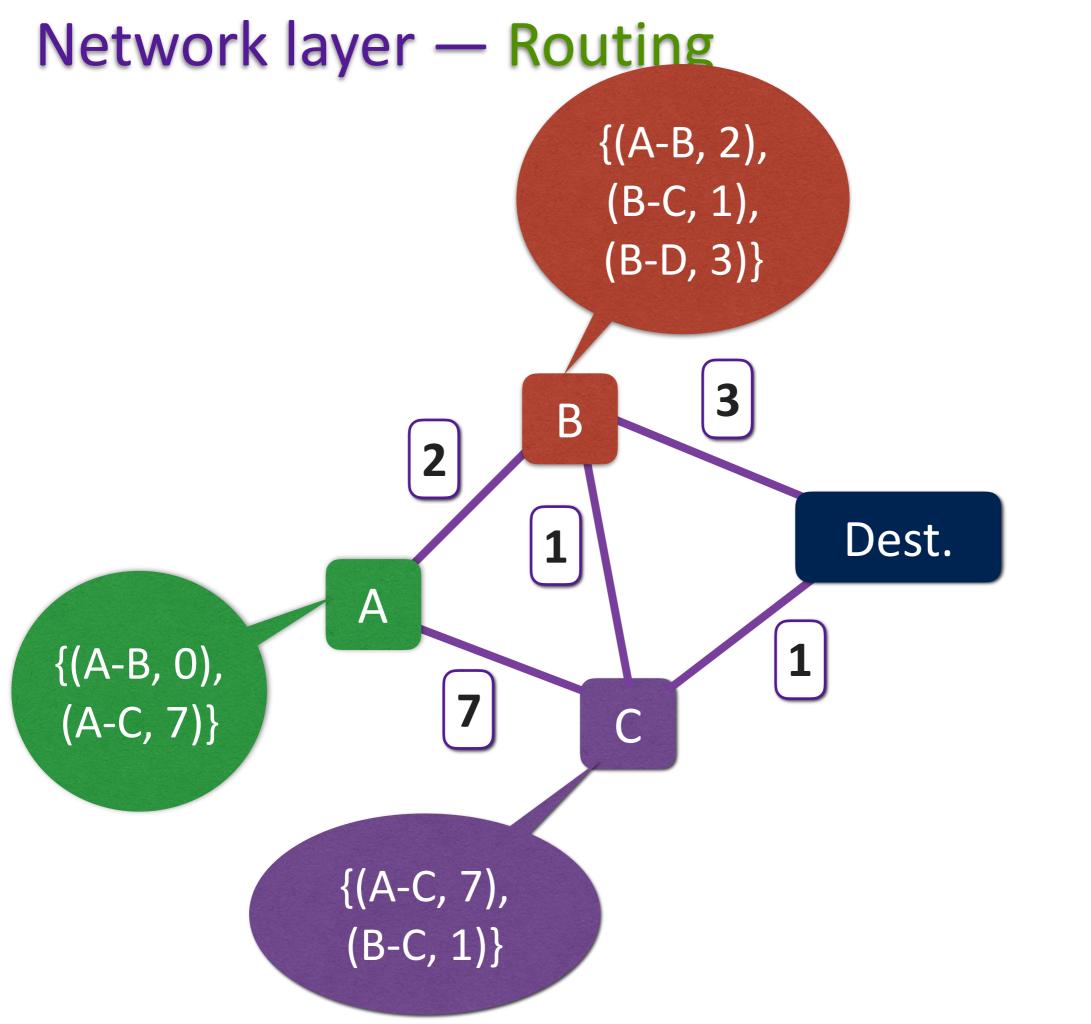
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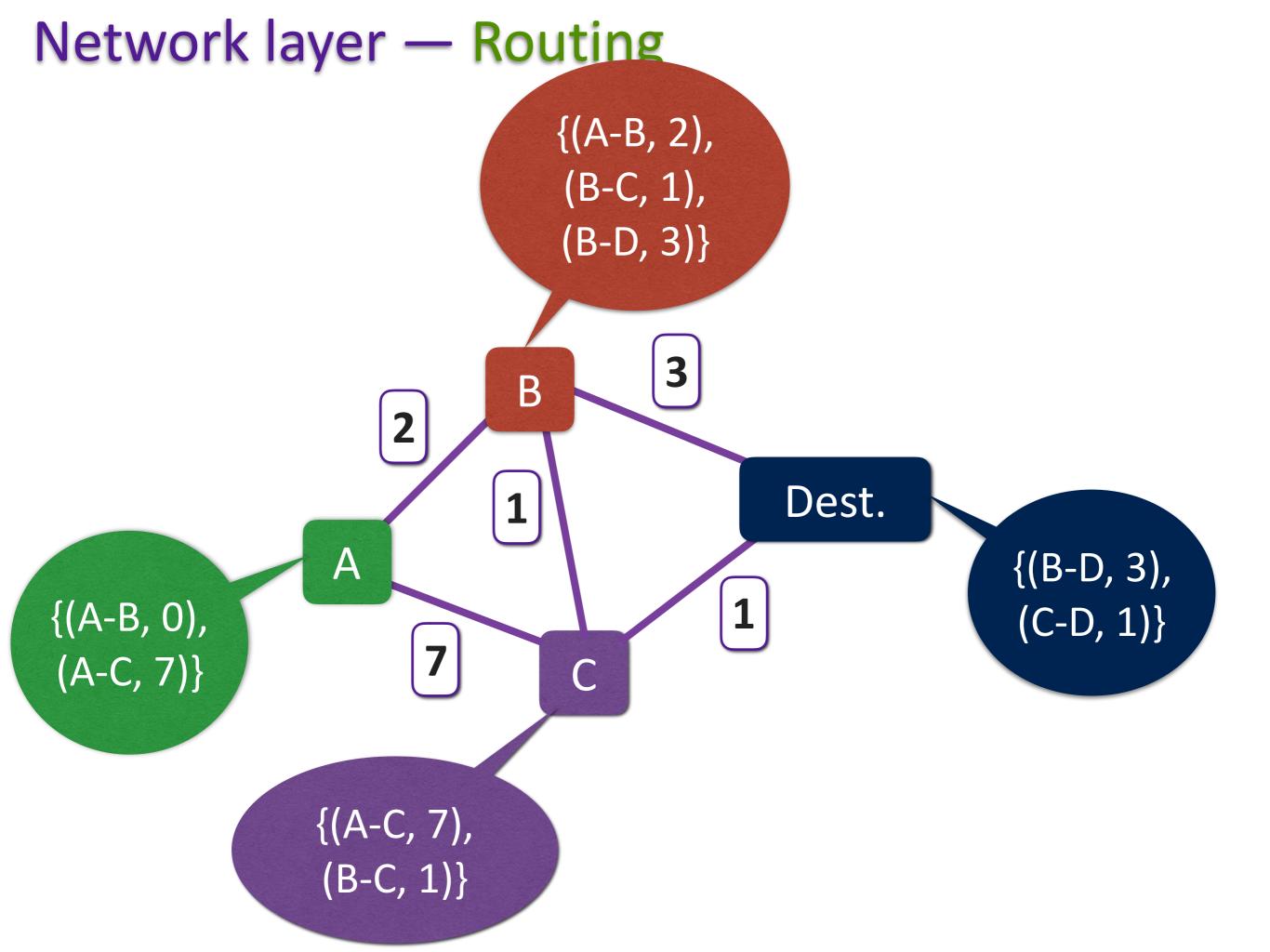
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- Challenge?

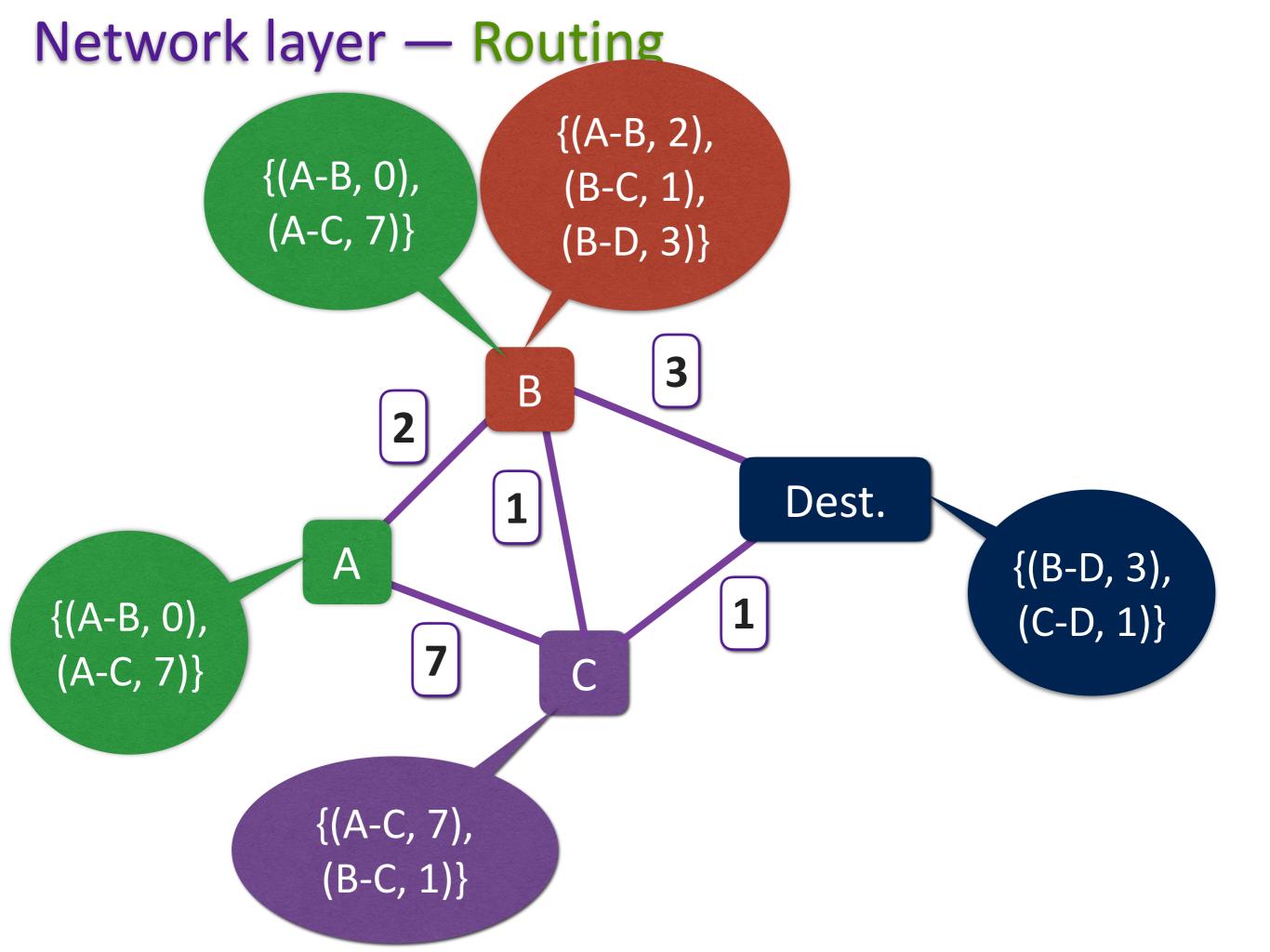


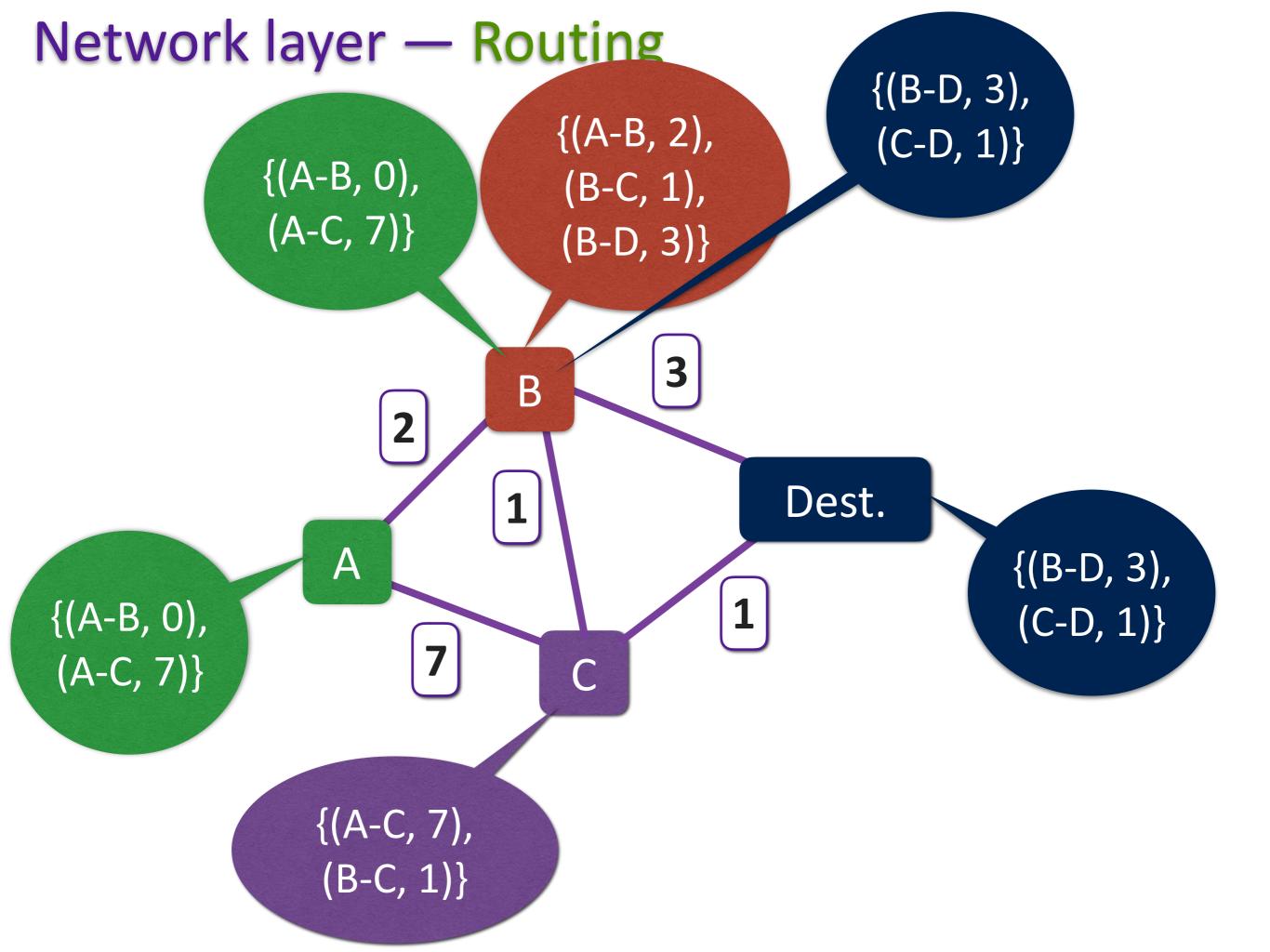


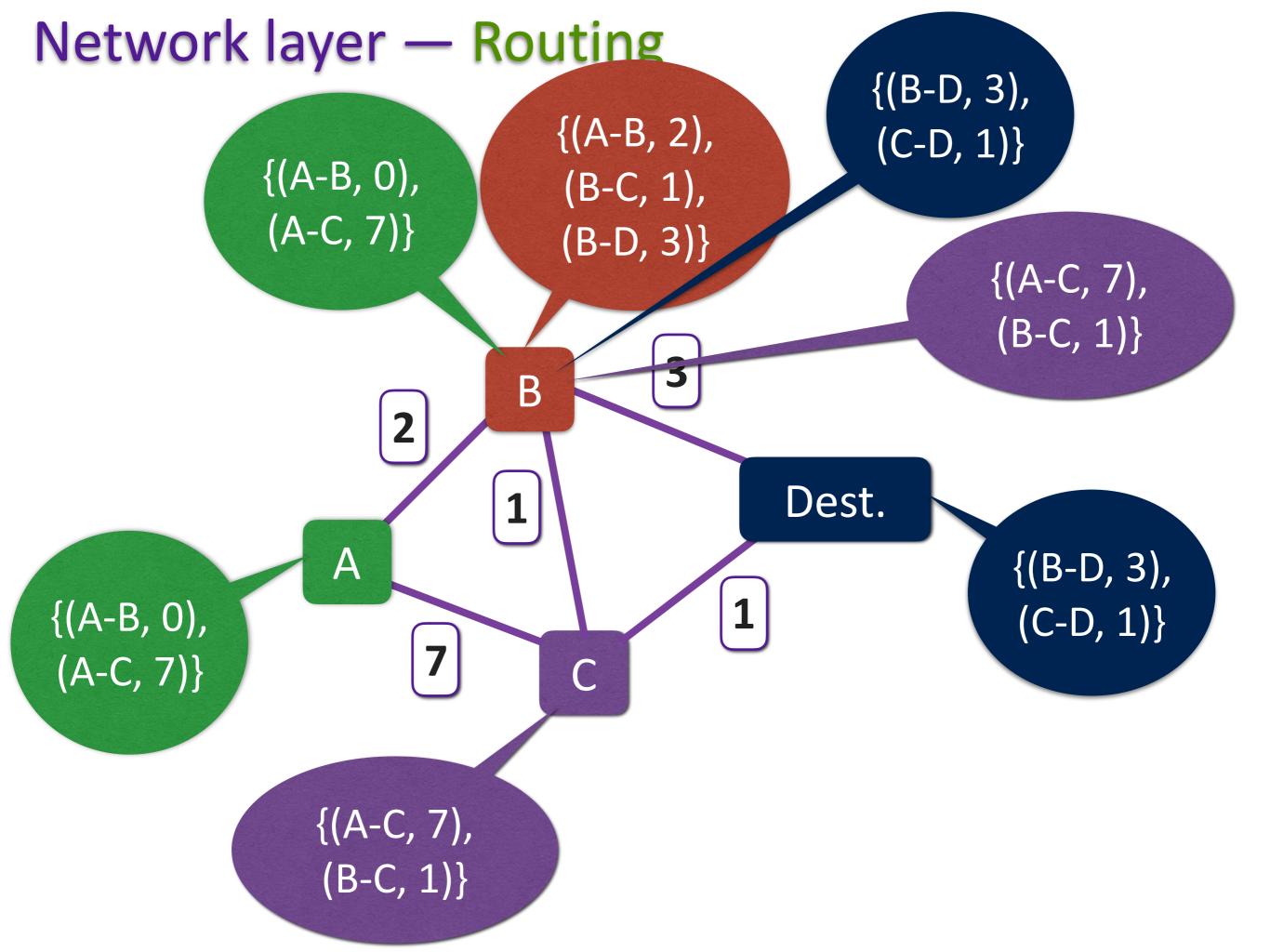












**Attempt 2: Link State Routing** 

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### **Attempt 3: Distance Vector Routing**

Each router

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    - d(A, dest) = min{d(A, neighbor) + d(neighbor, dest)}
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  - Broadcast to all its neighbors