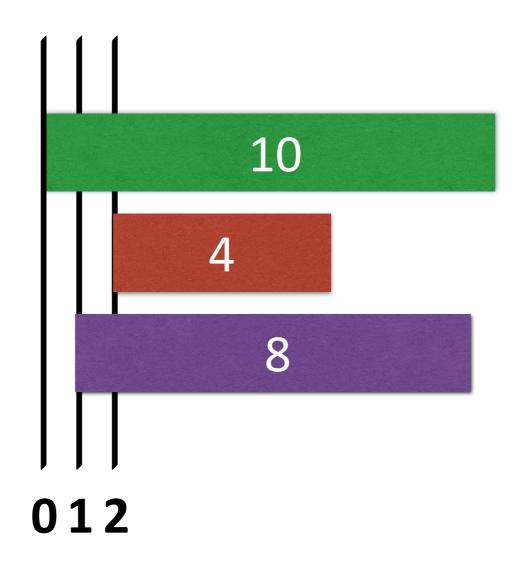
# CS4410/11: Operating Systems

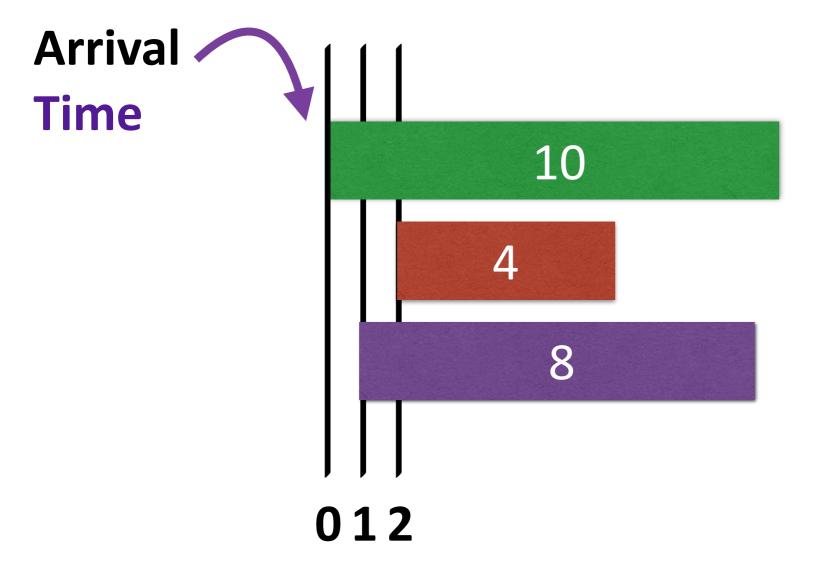
CPU Scheduling (Recap)

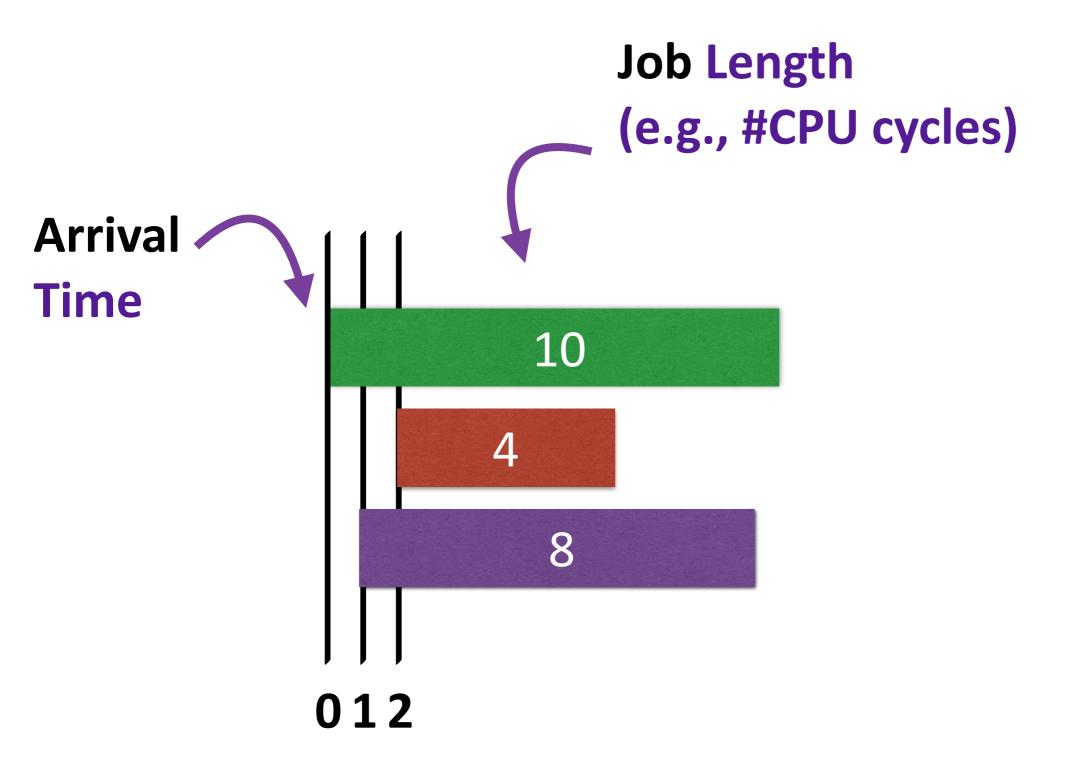
Networking

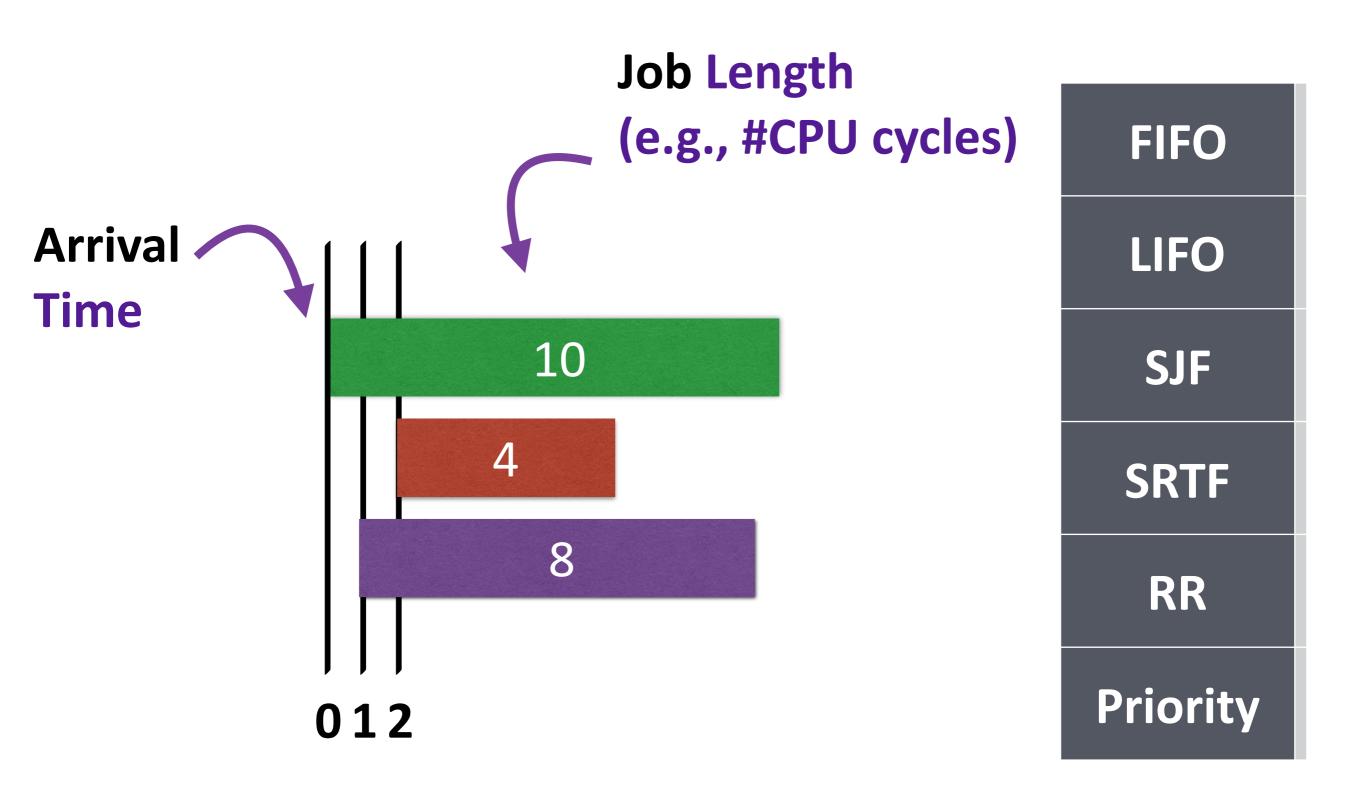
Rachit Agarwal Anne Bracy











So far: focused on what happens on a "machine"!

So far: focused on what happens on a "machine"!

- Networking
  - How do machines communicate?

So far: focused on what happens on a "machine"!

- Networking
  - How do machines communicate?
- Lets start with a simple analogy
  - How to move stuff from München to Ithaca?

Four "concepts"!

#### Four "concepts"!

- Layering
  - Abstraction is the key to manage complexity

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- Naming
  - A name for each computer, protocol, ...

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Abstraction is the key to manage complexity

#### Naming

A name for each computer, protocol, ...

#### Protocols

Computers, network devices speaking the same language

#### Four "concepts"!

#### Layering

Abstraction is the key to manage complexity

#### Naming

A name for each computer, protocol, ...

#### Protocols

Computers, network devices speaking the same language

#### Resource Allocation

Share resources (bandwidth, wireless spectrum, paths, ...)

Five "layers"!

Modularity

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  - Each layer relies on services from layer below

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### Five "layers"!

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#### Interfaces

Hide implementation details

### Five "layers"!

#### Modularity

- Each layer relies on services from layer below
- Each layer exports services to layer above

#### Interfaces

- Hide implementation details
- Layers can change without disturbing other layers

Five "layers"!

Physical layer

Five "layers"!

Physical layer

Five "layers"!

Link layer

Physical layer

Five "layers"!

Link layer

Physical layer

Postman

Five "layers"!

Network

Link layer

Physical layer

Postman

Five "layers"!

Network

Link layer

Physical layer

Airplane/rail

Postman

Five "layers"!

Transport

Network

Link layer

Physical layer

Airplane/rail

Postman

Five "layers"!

Transport

Network

Link layer

Physical layer

Post office

Airplane/rail

Postman

Five "layers"!

Application

Transport

Network

Link layer

Physical layer

Post office

Airplane/rail

Postman

Five "layers"!

Application

Transport

Network

Link layer

Physical layer

You

Post office

Airplane/rail

Postman

## Networking — Physical layer

Application

Transport

Network

Link

Physical

# Networking — Physical layer

Transfer of bits

Application

Transport

Network

Link

# Networking — Physical layer

- Transfer of bits
  - 0s and 1s

Application

Transport

Network

Link

# Networking — Physical layer

- Transfer of bits
  - 0s and 1s
  - Not concerned with protocols

Application

Transport

Network

\_ink

Application

Transport

Network

Link

**Link = Medium + Adapters** 

Application

Transport

Network

Link

**Link = Medium + Adapters** 

Communication Medium

Application

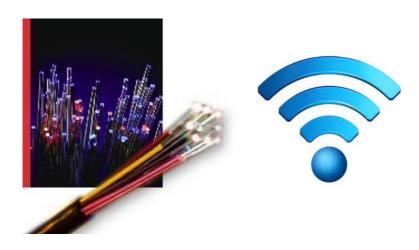
**Transport** 

Network

Link

## **Link = Medium + Adapters**

Communication Medium







Application

Transport

Vetwork

Link

Link = Medium + Adapters

Communication Medium







Application

Transport

Network

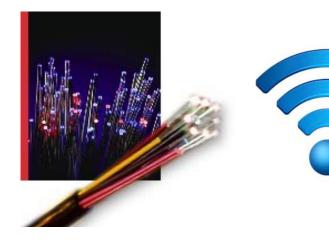
Link

Physical

Network Adapters (e.g., NIC — network interface card)

## **Link = Medium + Adapters**

Communication Medium







Application

**Transport** 

Vetwork

Link

Physical

Network Adapters (e.g., NIC — network interface card)









Application

Transport

Network

Link

**Broadcast links = Shared Medium** 

Application

Transport

Network

Link

## **Broadcast links = Shared Medium**

Everyone listens to everybody

Application

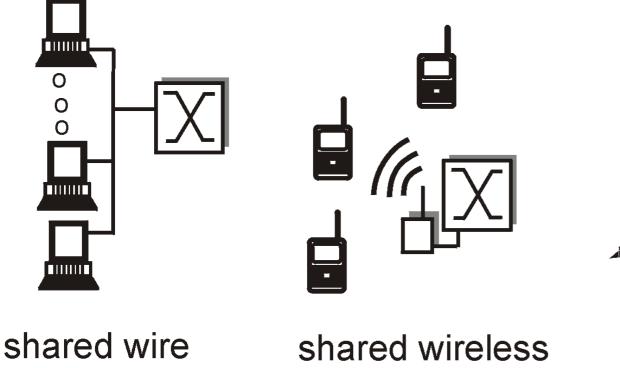
Transport

Vetwork

Link

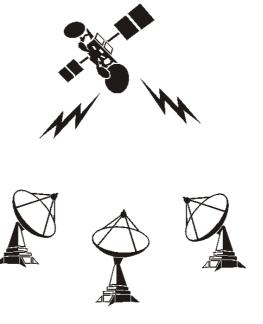
## **Broadcast links = Shared Medium**

Everyone listens to everybody



(e.g. Wavelan)

(e.g. Ethernet)



satellite

Link

Blah, blah, blah



**ZZZzzzzz**zzzzz

cocktail party

## **Broadcast links = Shared Medium**

Everyone listens to everybody

Application

Transport

Vetwork

Link

## **Broadcast links = Shared Medium**

Everyone listens to everybody

Application

Transport

Network

Link

## **Broadcast links = Shared Medium**

Everyone listens to everybody

Link

Physica

source

Adapter

## **Broadcast links = Shared Medium**

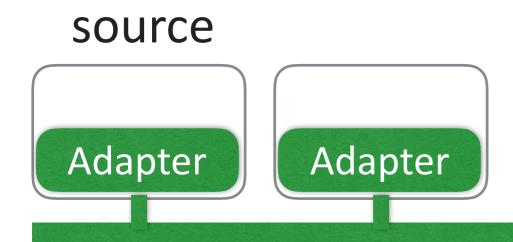
Everyone listens to everybody

Application

Transport

Network

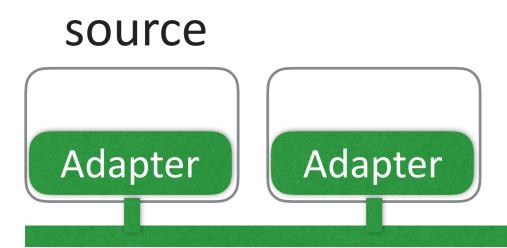
Link



## **Broadcast links = Shared Medium**

Everyone listens to everybody

Link destination





## **Broadcast links = Shared Medium**

Everyone listens to everybody

Application
Transport

Network

Link

Physical

destination







link-layer "protocol"

- Encoding data
  - Represented as a collection of 0s and 1s

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  - Detect and (optionally) correct errors

## Five "services"!

## Encoding data

Represented as a collection of 0s and 1s

## Framing

Put data packet into a frame; add receiver address

### Error detection and correction

Detect and (optionally) correct errors

### Flow control

- When to send/receive frames
- Depends on the protocol

## **Addresses**

Unique identifiers for sources and destinations

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  - Hierarchical allocation
    - Blocks: assigned to vendors (e.g., Dell) from IEEE
    - Adapters: assigned by the vendor from its block
- What if I want to send to everybody?
  - Special (broadcast) address: FF-FF-FF-FF-FF

### **Sharing a medium**

Ever been to a party?

- Ever been to a party?
  - Tried to have an interesting discussion?

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- Collisions

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Lets try to come up with a protocol to avoid collisions!

Attempt 1: Time sharing

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Never have a collision

- Underutilization of resources
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  - When I have something to speak, I wait for my turn

Lets try another protocol to avoid collisions

Attempt 2: Frequency sharing

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#### Problem

Overheads ...

### Lets try another protocol to avoid collisions

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- For wireless and optical mediums
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- What if I want to talk to only a few people in the group?

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- E.g., Divide into groups; each group talks among themselves

- Overheads ...
- What if I want to talk to only a few people in the group?
- What if I want to talk to people in different groups?
- E.g., one person wants to announce something ...

**Attempt 3: Carrier sense, Collision detection, Random access** 

Carrier Sense

- Carrier Sense
  - Listen before speaking

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  - .... and don't interrupt

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Wait for a random period of time

#### **Attempt 3: Carrier sense, Collision detection, Random access**

#### Carrier Sense

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- .... and don't interrupt

#### Collision detection

- Detect simultaneous speaking
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#### Random access

- Wait for a random period of time
- .... before trying to talk again

**Comparing the three approaches** 

Time division

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  - No collisions

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  - Underutilization of resources!

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- Time division
  - No collisions
  - Underutilization of resources!
  - What if token is lost?
- Frequency division
  - Overheads
- Random access
  - Efficient at low load, inefficient at high load (collisions)

#### **Ethernet uses CSMA/CD**

Carrier Sense: continuously listen to the channel

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  - (Exponentially increasing waiting times)

#### **Interesting Properties**

Distributed

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  - Why is that good?

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  - Why is that good?
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  - No state in the network
  - Cheap physical links

Connection-less, unreliable service

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Connection less

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Destination adapter does not acknowledge

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  - I am going to ignore the protocol
- Untrusted data access
  - I want to listen to what others are talking

Five "layers"!

Application

Transport

Network

Link layer

Physical layer

Five "layers"!

Application

Transport

Network

Link layer

Physical layer

Five "layers"!

Application

Transport

Network

Link layer

Physical layer

Deliver locally

Five "layers"!

Application

Transport

Network

Link layer

Physical layer

Deliver globally

Deliver locally

Five "layers"!

Application

Transport

Network

Link layer

Physical layer

Deliver (un)reliably

Deliver globally

Deliver locally

Five "layers"!

Application

Transport

Network

Link layer

Physical layer

**Deliver** 

Deliver (un)reliably

Deliver globally

Deliver locally

Five "layers"!

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Transport

Network

Link layer

Physical layer

Deliver

Deliver (un)reliably

Deliver globally

Deliver locally

### Three concepts

- Naming
  - A way to identify the source/destination

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  - E.g., house address

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#### Forwarding

#### Three concepts

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- E.g., which airplane should the stuff go on

#### Forwarding

Actually "moving" towards the destination

#### Three concepts

#### Naming

- A way to identify the source/destination
- E.g., house address

#### Routing

- Finding "how to" move towards the destination
- E.g., which airplane should the stuff go on

#### Forwarding

- Actually "moving" towards the destination
- E.g., Using airplane/truck/rail

### **Naming**

• Give every computer a unique name

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  - Challenges?

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  - Scalability why?

- Give every computer a unique name
  - Challenges?
  - Scalability why?
  - Assignment why?

### Naming

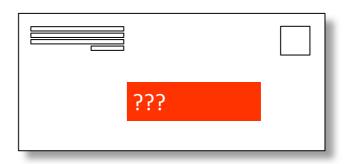
Hierarchical addressing

### **Naming**

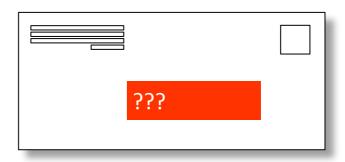
Hierarchical addressing



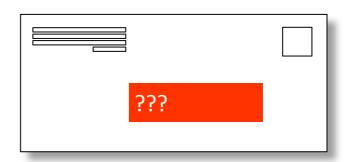
- Hierarchical addressing
  - E.g., addresses for houses



- Hierarchical addressing
  - E.g., addresses for houses
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- Hierarchical addressing
  - E.g., addresses for houses
  - Country: USA
  - City, State: Ithaca, NY
  - Number, Street: 306 State St.
  - Name: Rachit Agarwal



#### Hierarchical addressing

Country City, State Street,
Number Occupant

Country	City, State	Street, Number	Occupant
(8 bits)	(8 bits)	(8 bits)	(8 bits)

Country	City, State	Street, Number	Occupant
(8 bits)	(8 bits)	(8 bits)	(8 bits)
1000000	0-1010100	10001011	00000-101

Country	City, State	Street, Number	Occupant
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128	84	139	5

#### Hierarchical addressing

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128	84	139	5

Network

### Hierarchical addressing

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1000000	0-1010100	10001011	00000-101
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Network Machine

#### Hierarchical addressing

Country	City, State	Street, Number	Occupant
(8 bits)	(8 bits)	(8 bits)	(8 bits)
1000000	0-1010100	10001011	00000-101
128	84	139	5

Network

Machine

IP address: 128.84.139.5



### Hierarchical addressing

Why is it more scalable?



#### Hierarchical addressing

Why is it more scalable?





- Why is it more scalable?
  - Need to keep track of next step only!





- Why is it more scalable?
  - Need to keep track of next step only!
  - Flight to: USA



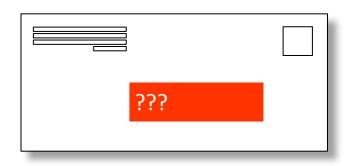


- Why is it more scalable?
  - Need to keep track of next step only!
  - Flight to: USA
  - Truck to: Ithaca, NY



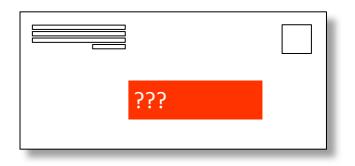


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  - Need to keep track of next step only!
  - Flight to: USA
  - Truck to: Ithaca, NY
  - Direction to: 306 State St.





- Why is it more scalable?
  - Need to keep track of next step only!
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  - Truck to: Ithaca, NY
  - Direction to: 306 State St.
  - Mailbox: Rachit Agarwal









### Hierarchical addressing

Why is it easier to assign?



- Why is it easier to assign?
  - Just assign a new machine a "local" address!



- Why is it easier to assign?
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  - E.g., adding a new machine to Cornell network



#### Hierarchical addressing

Why is it easier to assign?

- ???
- Just assign a new machine a "local" address!
- E.g., adding a new machine to Cornell network
- If last local address: 128.84.139.5

#### Hierarchical addressing

Why is it easier to assign?

- ???
- Just assign a new machine a "local" address!
- E.g., adding a new machine to Cornell network
- If last local address: 128.84.139.5
- New machine gets: 128.84.139.6

#### Three concepts

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  - A way to identify the source/destination
  - E.g., house address

#### Routing

- Finding "how to" move towards the destination
- E.g., which airplane should the stuff go on

#### Forwarding

- Actually "moving" towards the destination
- E.g., Using airplane/truck/rail

Lets come up with an approach? Generalize Ethernet ideas?

**Attempt 1: Broadcast** 

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Send to everybody

#### **Attempt 1: Broadcast**

- Send to everybody
- Goods

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  - Oh, well, simplicity

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  - Oh, well, everything else

- Send to everybody
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  - Oh, well, simplicity
- Not-so-goods
  - Oh, well, everything else
  - Bandwidth overheads

#### **Attempt 2: Time division Multiplexing**

Each source-destination pair assigned a time slot

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  - Can send data only during that slot

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  - No collisions

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  - No collisions
- Not-so-goods

- Each source-destination pair assigned a time slot
  - Can send data only during that slot
- Goods
  - No collisions
- Not-so-goods
  - Underutilization of resources

#### **Attempt 3: Frequency division Multiplexing**

• Each source-destination pair assigned a subset of resources

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  - Can use only "assigned" resources (e.g., bandwidth)

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  - Predictable performance

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#### **Attempt 2 and 3: Circuit Switching**

Source establishes connection

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  - Resources along the path are reserved

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- Source tears down connection

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  - Resources along the path are reserved
- Source sends data
  - Transmit data using the reserved resources
- Source tears down connection
  - Free resources for others to use

#### **Circuit Switching**

• Goods:

- Goods:
  - Predictable performance

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  - Reliable delivery

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  - Simple forwarding mechanism

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  - Predictable performance
  - Reliable delivery
  - Simple forwarding mechanism
- Not-so-goods
  - Resource underutilization

#### **Circuit Switching**

- Goods:
  - Predictable performance
  - Reliable delivery
  - Simple forwarding mechanism
- Not-so-goods
  - Resource underutilization
  - Blocked connections

#### **Circuit Switching**

#### Goods:

- Predictable performance
- Reliable delivery
- Simple forwarding mechanism

#### Not-so-goods

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- Blocked connections
- Connection set up overheads

#### **Circuit Switching**

#### Goods:

- Predictable performance
- Reliable delivery
- Simple forwarding mechanism

#### Not-so-goods

- Resource underutilization
- Blocked connections
- Connection set up overheads
- Per-connection state in switches (scalability problem)

#### **Attempt 4: Packet Switching**

Divide the message into packets

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- Put destination address in the header of each packet
  - Just like shipping stuff
- Each device stores a "look-up table"
  - Whats the next hop towards the destination?
- Destination receives the packet(s)
  - And reconstructs the message

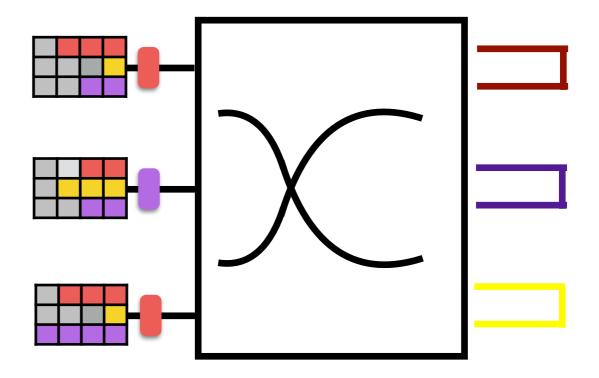
#### **Packet Switched forwarding**

Hop-by-hop forwarding

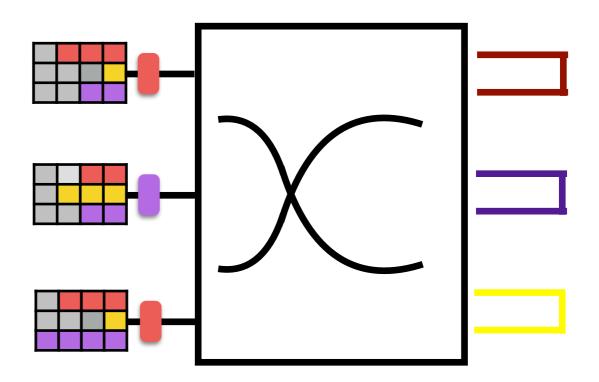
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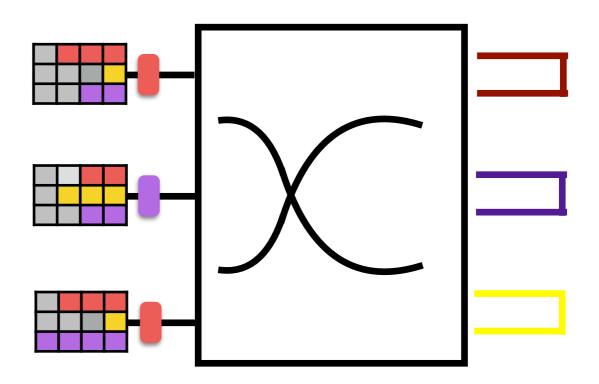
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- Hop-by-hop forwarding
- Each router has a "look-up table" (forwarding information base)
  - What should be stored in this table?
  - Prefix-based forwarding (longest-prefix matching)
    - Maps prefixes to the next-hop



### **Packet Switching**

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#### **Packet Switching**

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#### Not-so-goods:

Packet header overhead

#### **Packet Switching**

#### Goods:

- No resource underutilization
  - A source can send more if others don't use resources
- No blocked connection problem
- No per-connection state
- No set-up cost

#### Not-so-goods:

- Packet header overhead
- Network failures become a problem