CS4410/11: Operating Systems

CPU Scheduling (Recap)

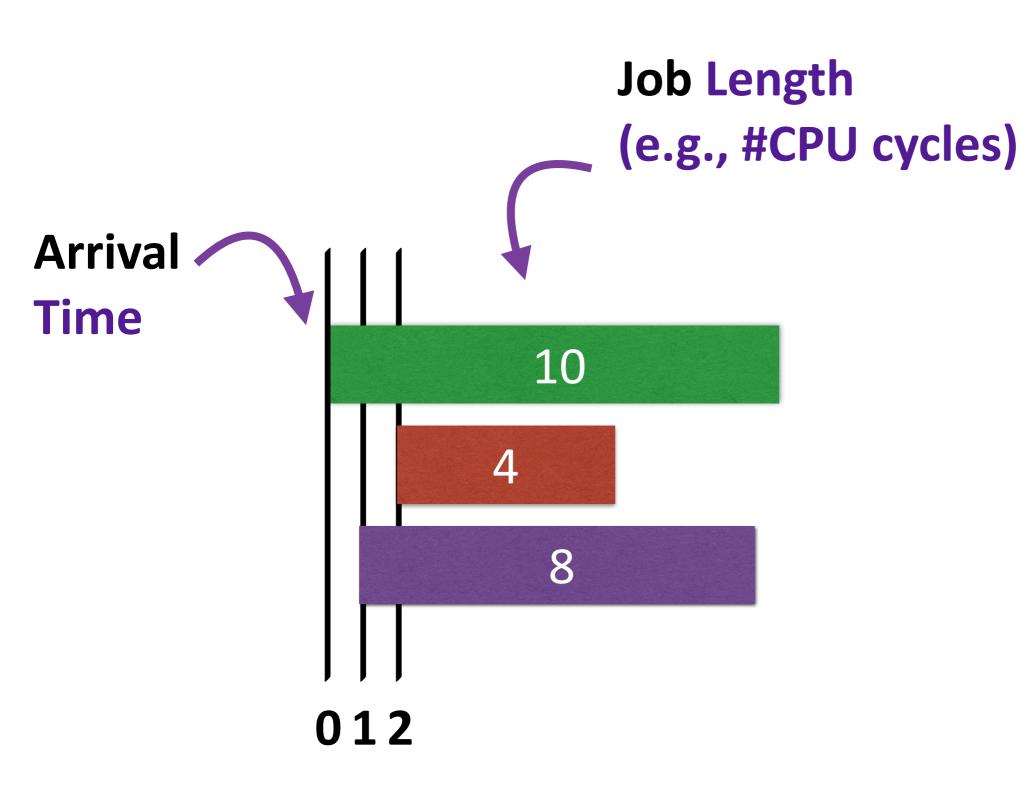
Networking

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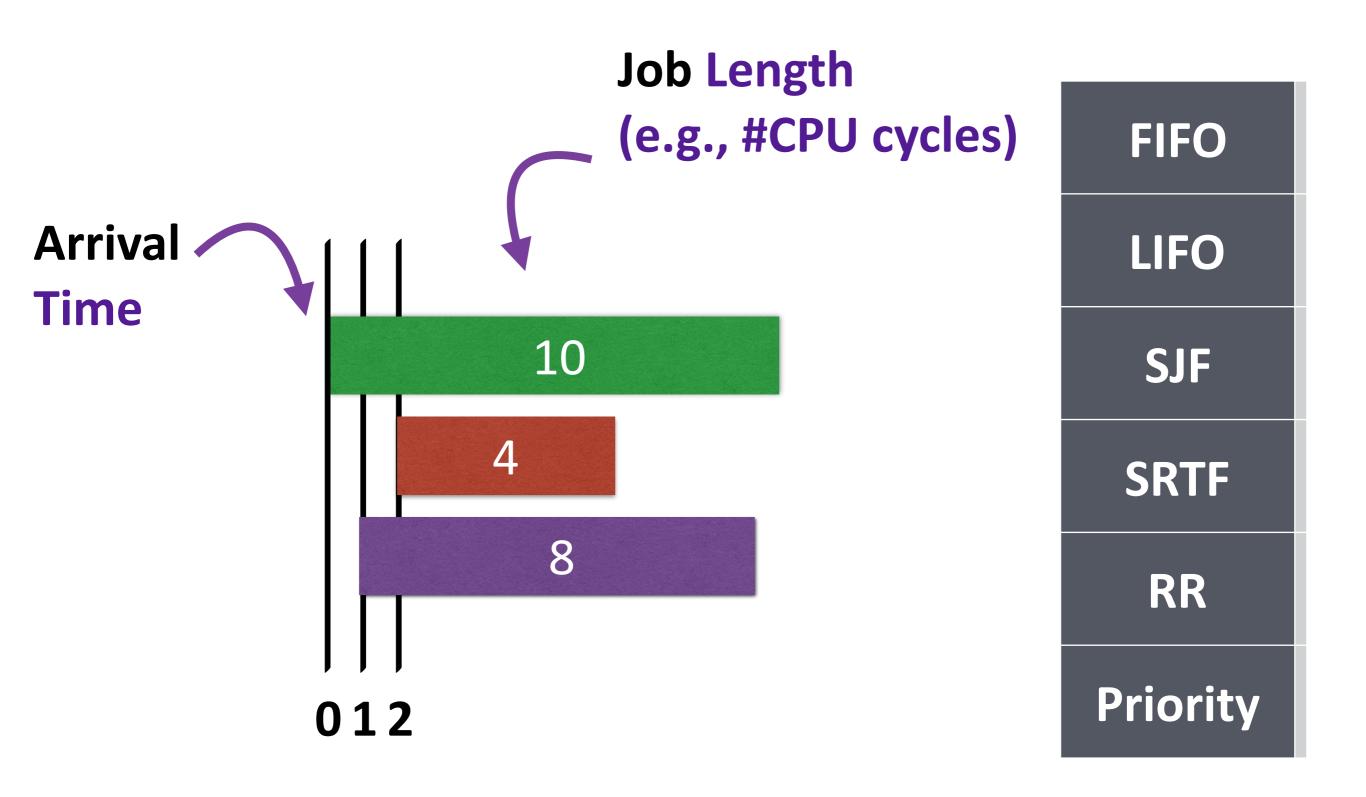
Slides based on material from Sirer, Rennesse, Rexford (Princeton)



CPU Scheduling — Example



CPU Scheduling — Example



Networking — What is it about?

- So far: focused on what happens on a "machine"!
- Networking
 - How do machines communicate?
- Lets start with a simple analogy
 - How to move stuff from München to Ithaca?

Networking — Key Concepts

Four "concepts"!

- Layering
 - Abstraction is the key to manage complexity
- Naming
 - A name for each computer, protocol, ..

• Protocols

• Computers, network devices speaking the same language

Resource Allocation

• Share resources (bandwidth, wireless spectrum, paths, ...)

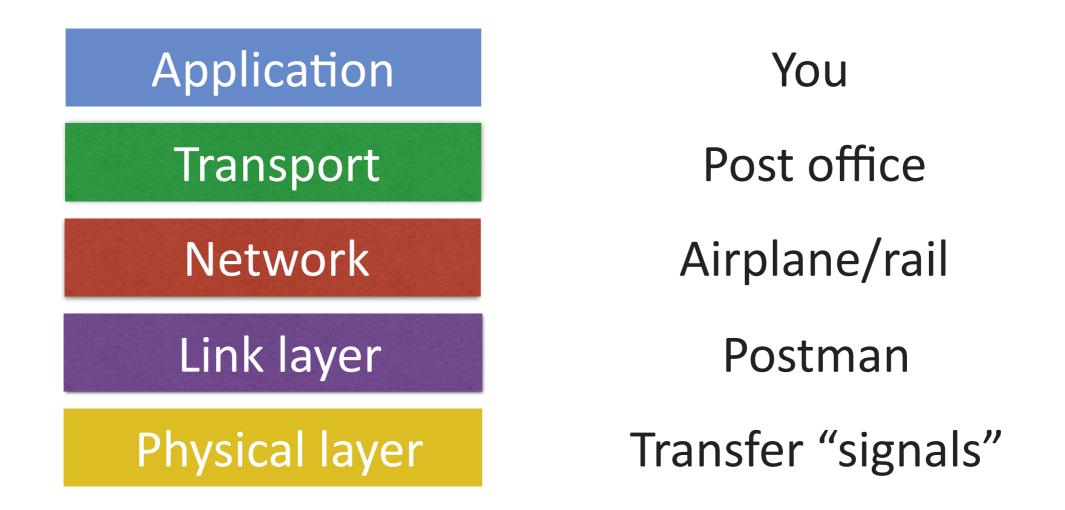
Networking — A Stack of Protocol Layers

Five "layers"!

- Modularity
 - Each layer relies on services from layer below
 - Each layer exports services to layer above
- Interfaces
 - Hide implementation details
 - Layers can change without disturbing other layers

Networking — A Stack of Protocol Layers

Five "layers"!



Networking — Physical layer

Transfer of bits

- 0s and 1s
- Not concerned with protocols

Application Transport Network Link Physical

Link = Medium + Adapters

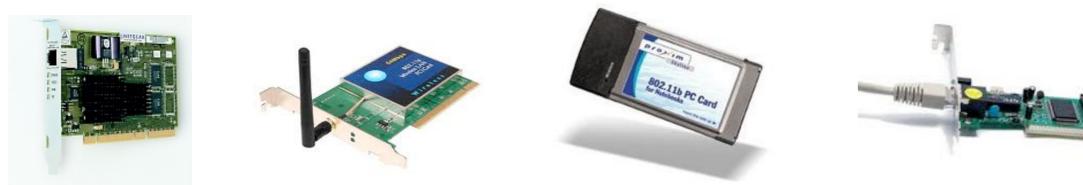
Communication Medium





Physical

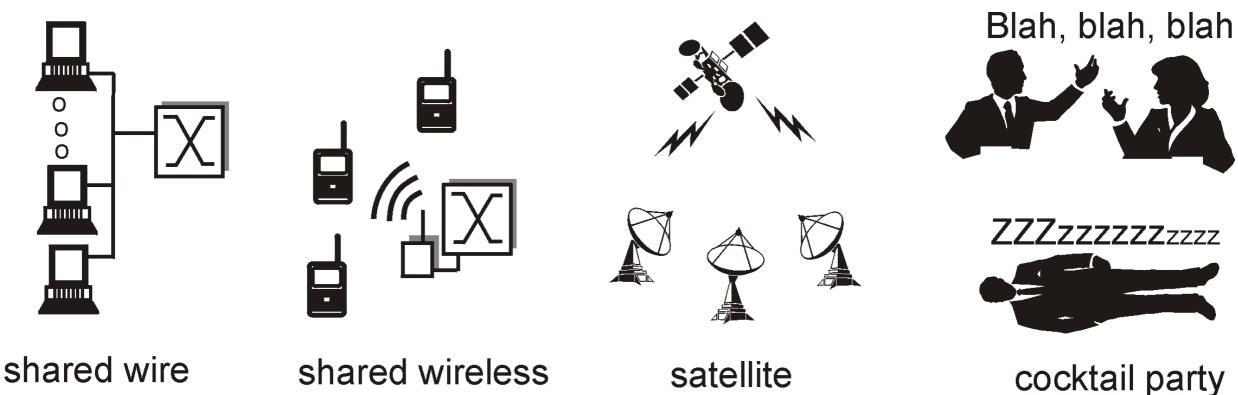
• Network Adapters (e.g., NIC — network interface card)



Broadcast links = Shared Medium

Everyone listens to everybody

Link



(e.g. Ethernet)

shared wireless (e.g. Wavelan)

satellite

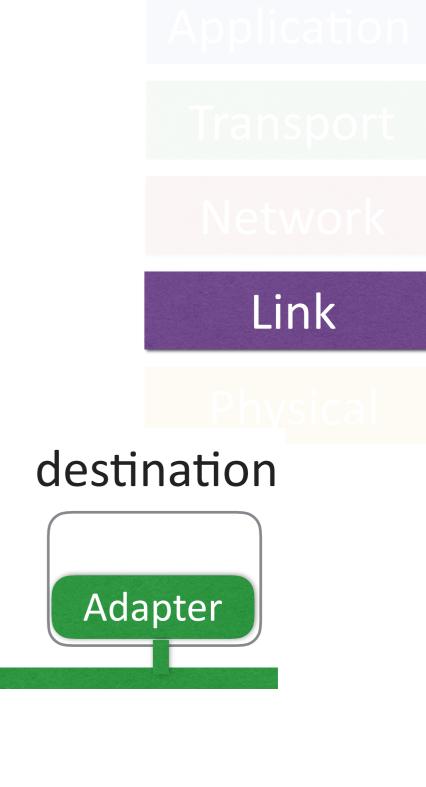
Broadcast links = Shared Medium

Adapter

• Everyone listens to everybody

source

Adapter





link-layer "protocol"

Five "services"!

- Encoding data
 - Represented as a collection of 0s and 1s
- Framing
 - Put data packet into a frame; add receiver address

Error detection and correction

- Detect and (optionally) correct errors
- Flow control
 - When to send/receive frames
 - Depends on the protocol

Addresses

Unique identifiers for sources and destinations

- "Hard-coded" in the adapter
- MAC address (e.g., 00-15-C5-49-04-A9)
- Hierarchical allocation
 - Blocks: assigned to vendors (e.g., Dell) from IEEE
 - Adapters: assigned by the vendor from its block
- What if I want to send to everybody?
 - Special (broadcast) address: FF-FF-FF-FF-FF-FF

Sharing a medium

- Ever been to a party?
 - Tried to have an interesting discussion?
- Collisions



Lets try to come up with a protocol to avoid collisions!

Attempt 1: Time sharing

- Everybody gets a turn to speak
- Goods
 - Never have a collision
- Problem
 - Wasted resources
 - During my turn, I may have nothing to speak
 - When I have something to speak, I wait for my turn

Lets try another protocol to avoid collisions

• Attempt 2: Frequency sharing

- Each person is assigned a particular frequency
- E.g., Divide into groups; each group talks among themselves

• Problem

- What if I want to talk to others?
- E.g., one person wants to announce something ...

Attempt 3: Carrier sense, Collision detection, Random access

Carrier Sense

- Listen before speaking
- and don't interrupt
- Collision detection
 - Detect simultaneous speaking
 - and shut up!
- Random access
 - Wait for a random period of time
 - before trying to talk again

Comparing the three approaches

Time division

- No collisions
- Wasted resources!
- What if token is lost?

Frequency division

- Efficient and fair at high load
- Inefficient at low load!

Random access

• Efficient at low load, inefficient at high load (collisions)

Ethernet uses CSMA/CD

- Carrier Sense: continuously listen to the channel
 - If idle: start transmitting
 - If busy: wait until idle
- Collision Detection: listen while transmitting
 - No collision: transmission complete
 - Collision: abort transmission; send jam signal
- Random access: exponential back off
 - After collision, transmit after "waiting time"
 - After k collisions, choose "waiting time" from {0, ..., 2^k-1)
 - (Exponentially increasing waiting times)

Interesting Properties

Distributed

- Distributed
 - No Central arbitrer

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 - Why is that good?

- Distributed
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 - Why is that good?
- Inexpensive

Interesting Properties

Distributed

- No Central arbitrer
- Why is that good?

Inexpensive

- No state in the network
- Cheap physical links

Connection-less, unreliable service

Connection less

- E.g., I am going to talk to you without getting permission first
- Networking terminology: No "handshaking"
- Unreliable
 - Destination adapter does not acknowledge
 - Did you listen to what I said?
 - Adversarial behavior could bring the connections down
 - I am going to ignore the protocol
 - Untrusted data access
 - I want to listen to what others are talking