# CS4410/11: Operating Systems

# **CPU Scheduling (Recap)**

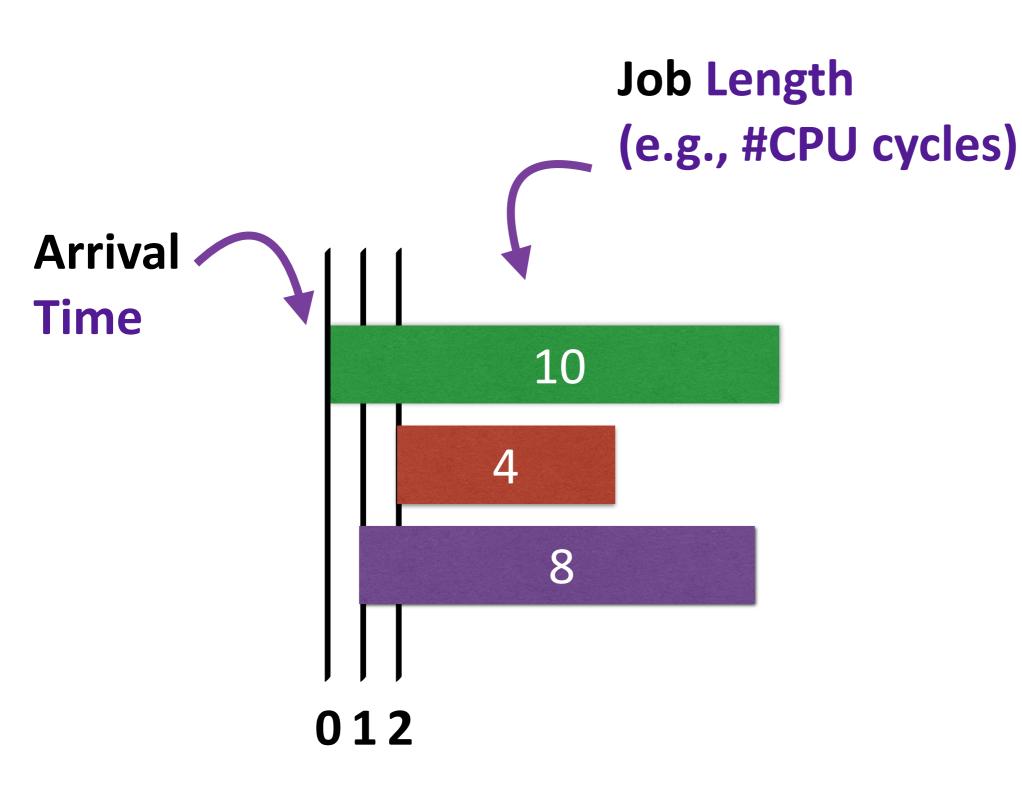
## Networking

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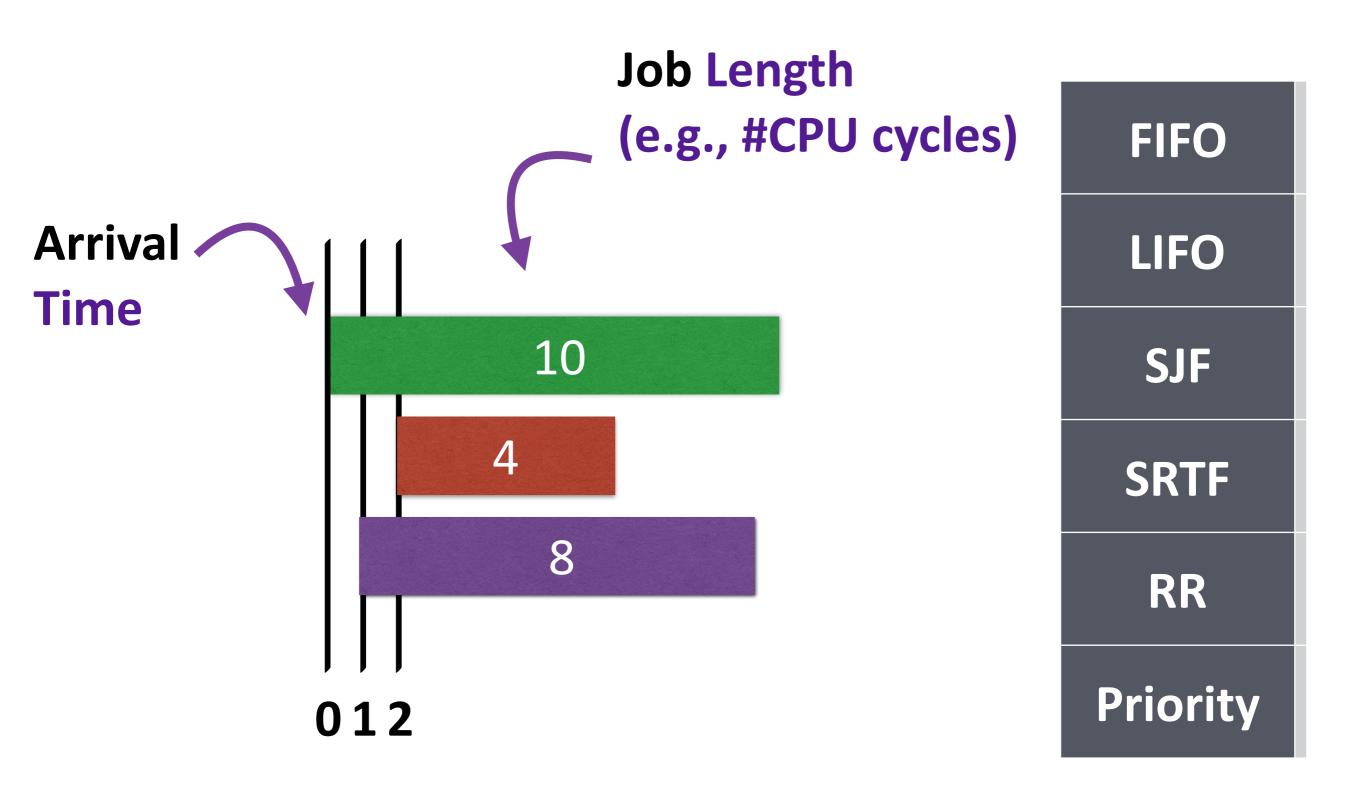
Slides based on material from Sirer, Rennesse, Rexford (Princeton)



### CPU Scheduling — Example



### CPU Scheduling — Example



### Networking — What is it about?

- So far: focused on what happens on a "machine"!
- Networking
  - How do machines communicate?
- Lets start with a simple analogy
  - How to move stuff from München to Ithaca?

### Networking — Key Concepts

### Four "concepts"!

- Layering
  - Abstraction is the key to manage complexity
- Naming
  - A name for each computer, protocol, ..

### • Protocols

• Computers, network devices speaking the same language

#### Resource Allocation

• Share resources (bandwidth, wireless spectrum, paths, ...)

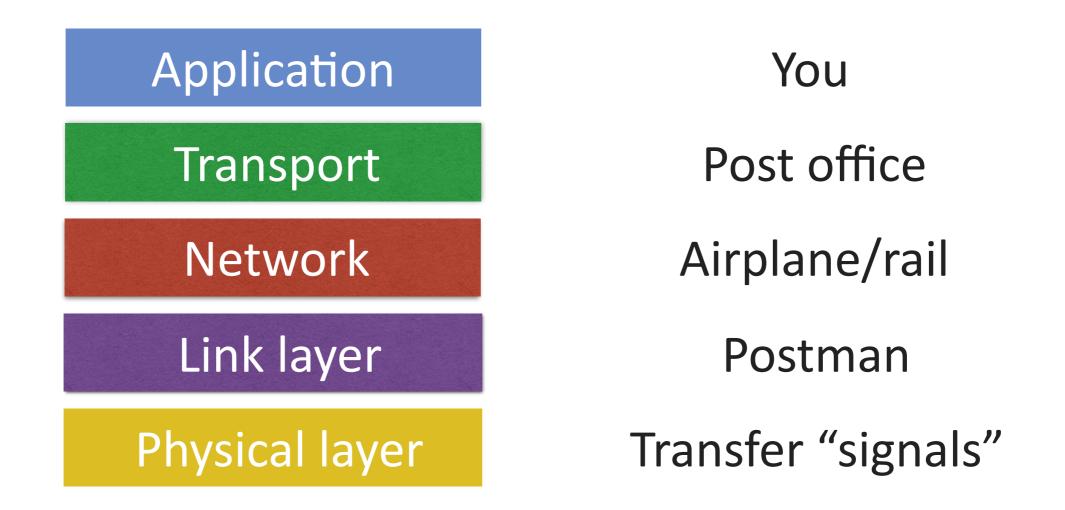
### Networking — A Stack of Protocol Layers

### Five "layers"!

- Modularity
  - Each layer relies on services from layer below
  - Each layer exports services to layer above
- Interfaces
  - Hide implementation details
  - Layers can change without disturbing other layers

### Networking — A Stack of Protocol Layers

Five "layers"!



## Networking — Physical layer

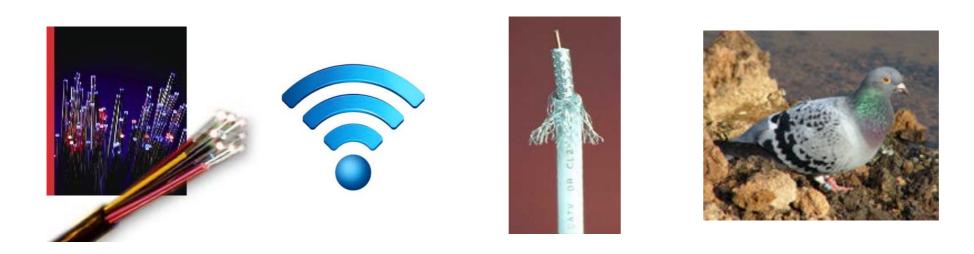
### Transfer of bits

- 0s and 1s
- Not concerned with protocols

Application Transport Network Link Physical

Link = Medium + Adapters

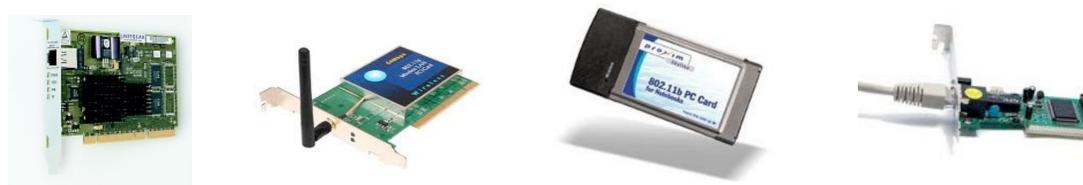
Communication Medium





Physical

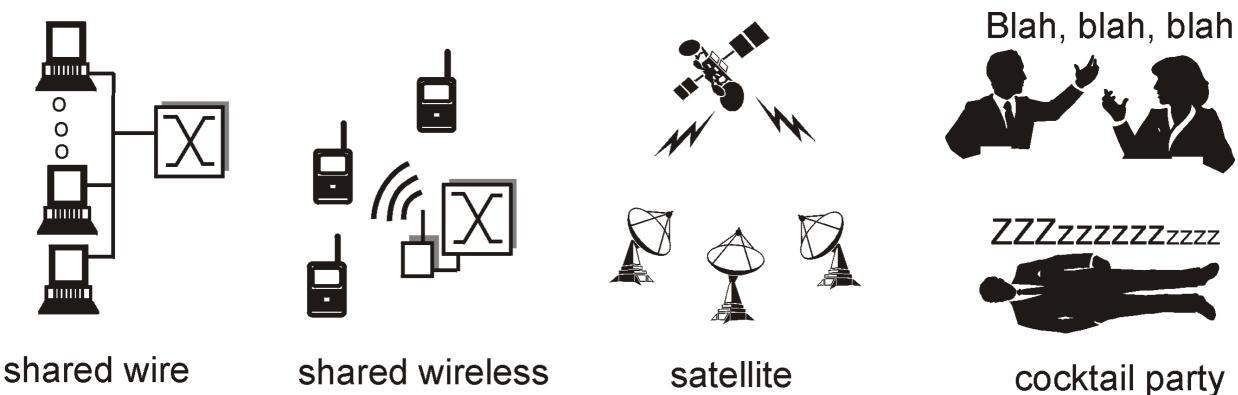
• Network Adapters (e.g., NIC — network interface card)



**Broadcast links = Shared Medium** 

Everyone listens to everybody

Link



(e.g. Ethernet)

shared wireless (e.g. Wavelan)

satellite

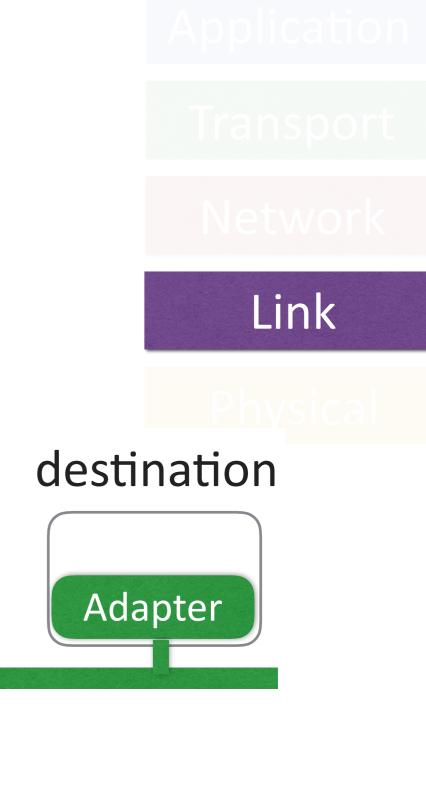
**Broadcast links = Shared Medium** 

Adapter

• Everyone listens to everybody

source

Adapter





link-layer "protocol"

Five "services"!

- Encoding data
  - Represented as a collection of 0s and 1s
- Framing
  - Put data packet into a frame; add receiver address

### Error detection and correction

- Detect and (optionally) correct errors
- Flow control
  - When to send/receive frames
  - Depends on the protocol

### Addresses

### Unique identifiers for sources and destinations

- "Hard-coded" in the adapter
- MAC address (e.g., 00-15-C5-49-04-A9)
- Hierarchical allocation
  - Blocks: assigned to vendors (e.g., Dell) from IEEE
  - Adapters: assigned by the vendor from its block
- What if I want to send to everybody?
  - Special (broadcast) address: FF-FF-FF-FF-FF-FF

### Sharing a medium

- Ever been to a party?
  - Tried to have an interesting discussion?
- Collisions



### Lets try to come up with a protocol to avoid collisions!

### Attempt 1: Time sharing

- Everybody gets a turn to speak
- Goods
  - Never have a collision
- Problem
  - Wasted resources
    - During my turn, I may have nothing to speak
    - When I have something to speak, I wait for my turn

### Lets try another protocol to avoid collisions

### • Attempt 2: Frequency sharing

- Each person is assigned a particular frequency
- E.g., Divide into groups; each group talks among themselves

### • Problem

- What if I want to talk to others?
- E.g., one person wants to announce something ...

#### **Attempt 3: Carrier sense, Collision detection, Random access**

#### Carrier Sense

- Listen before speaking
- .... and don't interrupt
- Collision detection
  - Detect simultaneous speaking
  - .... and shut up!
- Random access
  - Wait for a random period of time
  - .... before trying to talk again

### **Comparing the three approaches**

#### Time division

- No collisions
- Wasted resources!
- What if token is lost?

#### Frequency division

- Efficient and fair at high load
- Inefficient at low load!

#### Random access

• Efficient at low load, inefficient at high load (collisions)

### **Ethernet uses CSMA/CD**

- Carrier Sense: continuously listen to the channel
  - If idle: start transmitting
  - If busy: wait until idle
- Collision Detection: listen while transmitting
  - No collision: transmission complete
  - Collision: abort transmission; send jam signal
- Random access: exponential back off
  - After collision, transmit after "waiting time"
  - After k collisions, choose "waiting time" from {0, ..., 2<sup>k</sup>-1)
  - (Exponentially increasing waiting times)

**Interesting Properties** 

Distributed

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  - No Central arbitrer

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#### **Interesting Properties**

#### Distributed

- No Central arbitrer
- Why is that good?

#### Inexpensive

- No state in the network
- Cheap physical links

### **Connection-less, unreliable service**

#### Connection less

- E.g., I am going to talk to you without getting permission first
- Networking terminology: No "handshaking"
- Unreliable
  - Destination adapter does not acknowledge
    - Did you listen to what I said?
  - Adversarial behavior could bring the connections down
    - I am going to ignore the protocol
  - Untrusted data access
    - I want to listen to what others are talking