## 15: Networking Basics

Last Modified:
7/3/2004 1:47:08 PMM

## Networking

$\square$ So far we have talked primarily about OS support for individualcomputer systems
$\square$ Today we are going to talk about networking computer systems together

## A Network.

$\square$ A network is simply a collection of nodes, connected by links, that communic ate and cooperate
○ $\mathcal{N o d e s}=$ End Hosts $(\mathcal{P C} s, \mathcal{P D} \mathcal{A}$, toasters?), Internal $\mathcal{N}$ odes (Routers, switches, fubs,..)


## Communication?

- If two entities are going to communicate, they must agree on the expected order and meaning of messages they exchange.
$\square$ Asking for the time protocol
O SUCCESS FUL PROTO CO LEXCHANVG
- Hi ...Hi..Got the time?...t wo oclock
- ABORTED PROTOCOL
- Hi.. Don't bother meXX
- PROTOCOL $\mathcal{M I S} \mathcal{M A} \mathcal{A} \mathcal{C H}$
- Allo..Hello..Q uelle feuere átil....XX\&6ankstare $>$


## Questions

$\square \mathcal{W h a t}$ will be the format of data exchanged? How do we agree on a language among all kinds of nodes?
$\square$ Transmission across links is faulty can corrupt/lose data. How can we reliably excfiange information?
$\square$ How do we find the right path between two nodes? If there are many how do we choose the best one?
$\square \mathcal{H o w}$ do nodes refer to one another or address one another?
$\square$ What is the operating systems role in all this?

## Networking protocols

$\square$ OKlets define the "Ianguage" for all interactions over the network??
O One single language that can support everything from we 6 browsing to email to ftp to distributed file systems?
$\square$ Human beings are able to fandle lots of complexity in the ir protocol processing. - Ambiguously defined protocols - Many protocols all at once
$\square \mathcal{H o w}$ do computers manage comple $\chi$ protocol processing?

## Layered Arcfitectures

$\square \mathcal{B r e a k - u p ~ d e s i g n ~ p r o b l e m ~ i n t o ~ s m a l l e r , ~ m o r e ~}$ manageable problems

- Layers
$\square$ Design protocols to support each well define d task
- Not one language for everything!!



## Packet S witcfing

$\square$ Packets indicate their destination
$\square \mathfrak{N}$ o predetermined path for a packet to take
$\square$ Each intermediate note routes the packet closer to its destination


## Passenger Forwarding ()



## Traceroute/tracert



## Internet Map

$\square$ Traceroute gives one slice through the Internet topology
$\square$ What does the Internet really look like?
O That is a actually a hard question to answer
O Internet Atlas Project

- hittp://www.caida.org/projects/interne tatlas/
- Tecfiniques, software, and protocols for mapping the Internet, focusing on Internet topology, performance, workload, and routing data


CAIDA: NS FNNET growth until $\underline{1995}$

Backbone nodes elevated


## $\mathcal{N S} \mathcal{F} \mathcal{N e}$ tworking Arcfitecture of Late 1990 s

$\square \mathcal{N S} \mathcal{F} \mathcal{N}(E T$ Backbone Project successfully transitioned to a ne wetworking architecture in 1995.

- vBNNS (very figh speed Backbone Network Services). N(SF funded, provided by $\mathcal{M C I}$
04 original $\mathfrak{N e}$ twork Access Points ( $\mathfrak{N} \mathcal{F} \mathcal{F}$ awarded)
- NSS funded Routing Arbiter project
- NetworkService Providers (not $\mathcal{N} S \mathcal{F}$ funde d)


## Network Access Point

$\square \mathcal{A l l o w s}$ Internet Service Providers (IS $\mathcal{P}_{S}$ ), government, research, and educational organizations to interconnect and exchange information
$\square$ IS $P$ s connect the ir networks to the $\mathcal{N} \mathcal{A} P$ for the purpose of exchanging traffic with other IS Ps
$\square S u c h e x c h a n g e$ of Internet traffic is often referred to as "peering"


```
DNS: Domain Name System
People:many identifiers: Domain Name System:
    SSNN, name, Passport #
Internet flosts, routers:
    O IPaddress (32 6it)
        used for addressing
        datagrams
    O "name", e.g.,
        gaia.cs.umass.edu - used
        by fumans
Q:map between IT
    addresses and name ?
\square distributed database
    implemented in fierarchy of
    many name servers
    \square application-layer protocol
    host, routers, name servers to
    communicate to resolve names
    (address/name trans(ation)
    O note:core Internet
        function implemented as
        application-layer protocol
    O complexity at network's
                        "edge"
```


## Names and addresses: why both?

$\square \mathcal{N}$ ame: www.google .com
$\square$ IP address (one of them): 216.239.39.147

- (Also Etfiernet or other link-layer addresses.)
$\square$ IP addresses are fixed-size numbers.
O 32 6its. $216.239 .39 .147=$ 101011000.11101111 .00100111 .10010011
$\square \mathcal{N}$ ames are memorizable, flexible:
- Variable-length
- Many names for a single IP address.
- Change address doesn't imply change name.
- iPv6 addresses are 128 6it-even farder to memorize!


## Mapping $\mathfrak{N}$ ot 1 to 1

$\square$ One name may map to more than one IP address

- IP addresses are per network interface

O Multi-fomed machines have more than one ne tworkinterface - each with its own IP address
O Example: routers must be like this
$\square$ One IP address may map to more than one name
O One server machine may be the we 6 server (www.foo,com), mail server (mail.foo.com)etc.

## How to get names and numbers?

$\square$ Acquisition of $\mathcal{N}$ (ames and numbers are both regulated
O Why?

## How to get a name?

$\square$ First, get a domain name then you are free to assign sub names in that domain
O How to get a domain name coming up
$\square \mathcal{B e f o r e}$ you askfor a domain name though
O Should understand domain name structure...

- Know that you are responsible for providing authoritative $\mathcal{D N}(S$ server (actually a primary and one or more secondary $\mathcal{D N}(5$ servers) for that domain and registration information through "whois"


## Domain name structure


gTLDs= Generic Top Level Domains ccTLDs = Country Code Top Level Domains

## Top-level Domains (TLDs)

$\square$ Generic Top Leve (Domains (gTLDs)
O.com - commercial organizations

○.org-not-for-profit organizations

- .edu - educational organizations
- mil - military organizations
- gov-governmental organizations
O.net - network service providers

○ $\mathcal{N e}$ w: . 6 iz, .info, name, ...
$\square$ Country code Top Level Domains (cci $\mathcal{L D}$ ) O One for each country

## How to get a domain name?

믄1998, non-profit corporation, Internet Corporation for Assigned $\mathcal{N}$ ames and $\mathcal{N} u m b e r s$ (ICAN $\mathcal{N})$, was formed to assume responsibility from the USS Government
$\square I C A \mathcal{N} \mathcal{N}$ authorizes other companies to register domains in com, org and net and newg(TLDs

- Network Solutions is largest and in transitional period between US Govt and ICANN $\mathcal{N}$ had sole authority to register domains in com, org and net

How to get an $I \mathcal{P} \mathcal{A d d r e s s}$ ?
$\square \mathcal{A n s}$ wer 1: $\mathcal{N}$ ormally, answer is get an IP address from your upstream provider

- This is essential to maintain efficient routing!
$\square$ Answer 2: If you need lots of IP addresses then you can acquire your own block of them.
- IPaddress space is a scarce resource - must prove you have fully utilized a small block before can askfor a larger one and pay $\$ \$$ (gan 2002- $\$ 2250 /$ year for $/ 20$ and $\$ 18000 /$ year for a $/ 14$ )


## How to get lots of IT

Addresses? Internet Registries
RIPE NCC (Rise aux IP Europiens Ne twork Coordination Centre) for Europe, Middle.East, Africa
$\mathfrak{A P N} \operatorname{IC}(\mathfrak{A s i a}$ Pacific $\mathcal{N}$ etworkInformation Centre )for $\mathcal{A}$ sia and Pacific
$\mathfrak{A R I N}(\mathcal{A}$ merican Registry for Internet $\mathfrak{N}$ (umbers) for the Americas, the Caribbean, sub-safaran Africa
$\mathcal{N}$ ote: Once again regional distribution is important for efficient routing!
Can also get $\mathcal{A}$ utonomous System $\mathfrak{N u m n b e r s}$ ( $\mathcal{A S} \mathfrak{N} s$ from these registries


## Locating Resource

$\square$ www.cnn.com is the
name of a computer
(and, implicitly, of a
file in that computer)
$\square$ Ulse $\mathcal{D N}$ s to translate name to address


## Connection

$\square$ The protocol (fttp) sets up a connection (another protocol, tcp) between the host and cnn.com to transfer the page
$\square$ The connection transfers the page as a byte stream, without errors: flow control + error control


## Data flow

$\square$ The byte stream flows from end to end across many links and switches: routing (+ addressing)
$\square$ That stream is regulated and controlled by both ends: retransmission of erroneous or missing bytes; flow control


## Packets

$\square$ The network transports bytes grouped into packets
$\square$ The packets are "self. containe $d^{\prime \prime}$ and routers handle them one by one

- The end hosts worry about errors and
pacing
- Destination sends $\mathfrak{A C R s}$
- Source checks losses


## Port Numbers

$\square$ When a packet arrives at its destination, the operating system uses the destination port number to identify which application should receive it.
$\square \mathcal{T}$ fis is called demultiplexing.

## Bits

$\Delta$
$\square$ Equipment in each node sends the packets as a string of bits
$\square$ That equipment is not aware of the meaning of the bits


