

## Announcements

$\square$ For the next lecture, you MUST watch the tutorial on the shortest path algorithm beforehand:
http://www.cs.cornell.edu/courses/cs2110/2017f a/online/shortestPath/shortestPath.htm
$\square$ Thursday's lecture will assume that you understand it. Watch the tutorial once or twice and execute the algorithm on a small graph.

## Graphs



## Representing Graphs



Adjacency List


Adjacency Matrix


## Graph Algorithms

$\square$ Search

- Depth-first search
- Breadth-first search
$\square$ Shortest paths
- Dijkstra's algorithm
$\square$ Spanning trees
Algorithms based on properties
Minimum spanning trees
$\square$ Prim's algorithm
$\square$ Kruskal's algorithm


## Search on Graphs

$\square$ Given a graph ( $V, E$ ) and a vertex $u \in V$
$\square$ We want to "visit" each node that is reachable from $u$


There are many paths to some nodes.

How do we visit all nodes efficiently, without doing extra work?


## Depth-First Search

Intuition: Recursively visit all vertices that are reachable along unvisited paths.

```
/** Visit all nodes reachable
on unvisited paths from u.
Precondition: u is unvisited.
*/
public static void dfs(int u)
{
\[
\begin{gathered}
\text { visited[u] = true; } \\
\text { for all edges }(u, v): \\
\text { if(!visited[v]): } \\
\text { dfs(v); }
\end{gathered}
\]
```


$\mathrm{dfs}(1)$ visits the nodes in this order: $1,2,3,5,7,8$

## Depth-First Search

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```
/** Visit all nodes reachable
on unvisited paths from u.
Precondition: u is unvisited.
*/
public static void dfs(int u)
{
    visited[u] = true;
    for all edges (u,v):
        if(!visited[v]):
        dfs(v);
```

Suppose there are $n$ vertices that are reachable along unvisited paths and $e$ edges:

Worst-case running time? $O(n+e)$ Worst-case space? $O(n)$

## Depth-First Search in Java

```
public class Node {
    boolean visited;
    List<Node> neighbors;
```

Each vertex of the graph is an object of type Node
/** Visit all nodes reachable on unvisited paths from
this node.
Precondition: this node is unvisited.
public void dfs()
visited= true;
for (Node n : neighbors)
if (!n.visited) n.dfs();
\}
\}
\}

## Depth-First Search Iteratively

Intuition: Visit all vertices that are reachable along unvisited paths from the current node.

```
/** Visit all nodes reachable on
unvisited paths from u.
Precondition: u is unvisited.
public static void dfs(int u) {
    Stack s= (u);// Not Java!
    while (s is not empty) {
        u= s.pop();
        if (u not visited) {
        visit u;
        for each edge (u, v):
        s.push(v) ;
    }
    }
```


## Breadth-First Search

Intuition: Iteratively process the graph in "layers" moving further away from the source node.

```
/** Visit all nodes reachable on
unvisited paths from u.
Precondition: u is unvisited.
public static void bfs(int u) {
    Queue q= (u);// Not Java!
    while ( q is not empty ) {
        u= q.remove();
        if (u not visited) {
        visit u;
        for each (u, v):
        q.add (v);
        }
    }
                                    Queue:
\begin{tabular}{|l|l|l|l|l|l|l|}
\hline 2 & 5 & 7 & 3 & 5 & 8 & 5 \\
\hline
\end{tabular}
```

\}

## Analyzing BFS

Intuition: Iteratively process the graph in "layers" moving further away from the source node.

```
/** Visit all nodes reachable on
unvisited paths from u.
Precondition: u is unvisited. */
public static void bfs(int u) {
    Queue q= (u);// Not Java!
    while ( )
        u= q.remove();
        if (u not visited) {
        visit u;
        for each (u, v):
        q.add(v);
        }
    }
```

Suppose there are $n$ vertices that are reachable along unvisited paths and $e$ edges:

Worst-case running time? $O(n+e)$ Worst-case space? $O(e)$
\}

## Comparing Search Algorithms

## DFS

$\square$ Visits: 1, 2, 3, 5, 7, 8
$\square$ Time: $O(n+e)$
$\square$ Space: $O(n)$

## BFS

$\square$ Visits: 1, 2, 5, 7, 3, 8
$\square$ Time: $O(n+e)$
$\square$ Space: $O(e)$


