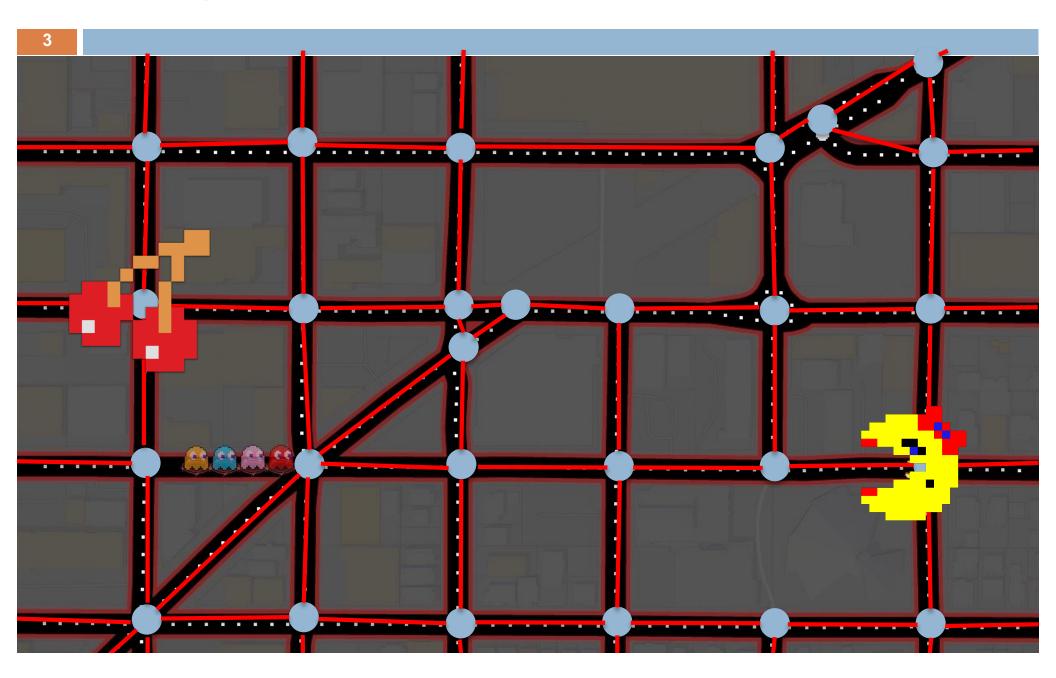


Lecture 17
CS 2110 Fall 2017

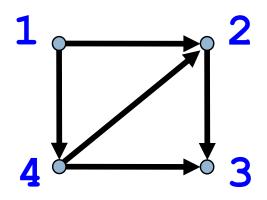
Announcements

- For the next lecture, you MUST watch the tutorial on the shortest path algorithm beforehand:
 - http://www.cs.cornell.edu/courses/cs2110/2017fa/online/shortestPath/shortestPath.html
- Thursday's lecture will assume that you understand it. Watch the tutorial once or twice and execute the algorithm on a small graph.

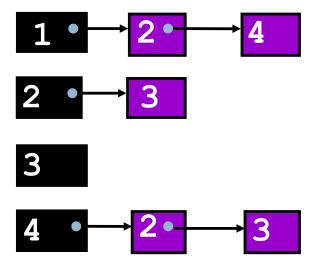
Graphs



Representing Graphs



Adjacency List



Adjacency Matrix

	1	2	3	4
1		1	0	1
2	0	0	1	0
3	0	0	0	0
4	0	1	1	0

Graph Algorithms

- Search
 - Depth-first search
 - Breadth-first search
- Shortest paths
 - Dijkstra's algorithm
- Spanning trees

Algorithms based on properties

Minimum spanning trees

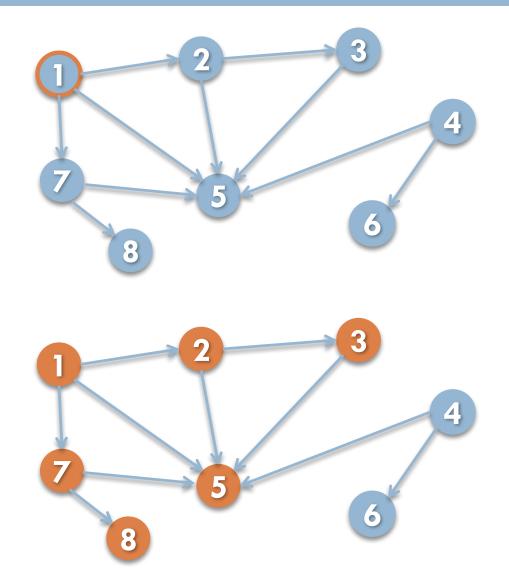
- Prim's algorithm
- Kruskal's algorithm

Search on Graphs

- □ Given a graph (V, E) and a vertex $u \in V$
- We want to "visit"
 each node that is
 reachable from u

There are many paths to some nodes.

How do we visit all nodes efficiently, without doing extra work?



Depth-First Search

Intuition: Recursively visit all vertices that are reachable along unvisited paths.

```
/** Visit all nodes reachable
on unvisited paths from u.
Precondition: u is unvisited.
* /
public static void dfs(int u)
    visited[u] = true;
    for all edges (u, v):
         if(!visited[v]):
             dfs(v);
                                 dfs(1) visits the nodes in this
                                 order: 1, 2, 3, 5, 7, 8
```

Depth-First Search

Intuition: Recursively visit all vertices that are reachable along unvisited paths.

```
/** Visit all nodes reachable
on unvisited paths from u.
Precondition: u is unvisited.
* /
public static void dfs(int u)
                            Suppose there are n vertices that
    visited[u] = true;
                            are reachable along unvisited
    for all edges (u, v):
                            paths and e edges:
         if(!visited[v]):
             dfs(v);
                            Worst-case running time? O(n + e)
                            Worst-case space? O(n)
```

Depth-First Search in Java

```
public class Node {
                                         Each vertex of the
      boolean visited;
                                         graph is an object
      List<Node> neighbors;
                                         of type Node
/** Visit all nodes reachable on unvisited paths from
this node.
Precondition: this node is unvisited.
                                         No need for a
      public void dfs() {
                                         parameter. The
             visited= true;
                                         object is the node.
             for (Node n: neighbors) {
                    if (!n.visited) n.dfs();
```

Depth-First Search Iteratively

Intuition: Visit all vertices that are reachable along unvisited paths from the current node.

```
/** Visit all nodes reachable on
unvisited paths from u.
Precondition: u is unvisited. */
public static void dfs(int u)
    Stack s=(u);// Not Java!
    while (s is not empty) {
        u = s.pop();
        if (u not visited) {
            visit u;
            for each edge (u, v):
                s.push(v);
                                       Stack:
```

Breadth-First Search

Intuition: Iteratively process the graph in "layers" moving further away from the source node.

```
/** Visit all nodes reachable on
unvisited paths from u.
Precondition: u is unvisited. */
public static void bfs(int u)
    Queue q=(u);// Not Java!
    while ( q is not empty ) {
        u= q.remove();
        if (u not visited) {
            visit u;
            for each (u, v):
                q.add(v);
```

Analyzing BFS

Intuition: Iteratively process the graph in "layers" moving further away from the source node.

```
/** Visit all nodes reachable on
unvisited paths from u.
Precondition: u is unvisited. */
public static void bfs(int u) {
    Queue q=(u);// Not Java!
    while (
        u= q.remove();
        if (u not visited)
            visit u;
            for each (u, v):
                q.add(v);
```

Suppose there are n vertices that are reachable along unvisited paths and e edges:

Worst-case running time? O(n + e)Worst-case space? O(e)

Comparing Search Algorithms

DFS

- \square Visits: 1, 2, 3, 5, 7, 8
- \square Time: O(n+e)
- \square Space: O(n)

BFS

- \square Visits: 1, 2, 5, 7, 3, 8
- \square Time: O(n+e)
- \square Space: O(e)

