

Finish with Recursion

CS211
Fall 2000

Another Recursive Descent Example

- Goal: Determine if the brackets () [] {} on a line are balanced and properly nested
- Recursive definition for LegalExp
 - The empty string is a LegalExp
 - If E is a LegalExp then so are (E), [E], and {E}
 - One or more LegalExps on a line make a LegalExp
- Examples:
 - legal: () [()]
 - legal: () {} []
 - illegal: ([])
 - illegal: ()

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Step 1: Build a Tokenizer

- We want to divide an input line (a String) into tokens
 - The Java API on the Web includes a java.util package
 - This package contains java.util.StringTokenizer
- StringTokenizer has a nextToken() method that throws NoSuchElementException if it runs out of tokens
- The StringTokenizer constructor takes
 - an input string,
 - a list of delimiters (token separators), and
 - a boolean flag (true implies that delimiters are also tokens)
- StringTokenizer also has other methods
 - countTokens()
 - hasMoreTokens()
 - a couple others

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Using StringTokenizer

- Can StringTokenizer be used directly? Almost...
 - We want to skip all "uninteresting" tokens
 - Would like an eol token
- We can alter StringTokenizer to do these extra things by using either
 - inheritance or
 - aggregation
- Inheritance:
 - We make our own tokenizer by *extending* StringTokenizer
 - Existing methods are either inherited or overridden
- Using inheritance:
 - Either we override all methods of StringTokenizer
 - Or we accept some inherited methods with "surprising" behavior

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Tokenizer Code

- Aggregation
 - We make our own tokenizer by using a StringTokenizer *within* our tokenizer class
 - Our class has only the methods that we choose to write
 - ▲ nextToken()
 - ▲ pushBack()

```
class MyTokenizer {
    private StringTokenizer tokenizer;
    private String currentToken;
    private boolean pushed;
    public static String PARENS = "()[]{}";

    public MyTokenizer (String inputString) {
        tokenizer = new
        StringTokenizer(inputString,PARENS,true);
        currentToken = null;
        pushed = false;
    }

    public void pushBack () {
        pushed = true;
    }
}
```

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nextToken()

```
public String nextToken () {
    if (pushed) {
        pushed = false;
        return currentToken;
    }
    try {
        do {
            currentToken = tokenizer.nextToken();
        } while (PARENS.indexOf(currentToken) == -1);
    } catch (NoSuchElementException e) {
        currentToken = "eol";
    }
    return currentToken;
}
}
```

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Step 2: Build a Parser

- Follow the recursive definition
 - The empty string is a LegalExp
 - If E is a LegalExp then so are (E), [E], and {E}
 - One or more LegalExps on a line make a LegalExp

```
class LegalExpParser {
    MyTokenizer in;
    public void legalExp () {
        String matcher;
        String token = in.nextToken();
        while ("({[",indexOf(token) != -1) {
            if (token.equals("(")) matcher = ")";
            else if (token.equals("[") matcher = "]";
            else matcher = "}";
            legalExp();
            token = in.nextToken();
            if (!token.equals(matcher)) error;
            token = in.nextToken();
        }
        in.pushBack();
    }
}
```

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The Rest of the Code

```
public void eval (String inputString) {
    in = new MyTokenizer(inputString);
    legalExp();
    if (in.nextToken().equals("eof")) return;
    error;
} // end LegalExpParser

// Code using a LegalExpParser
// The String string is to be evaluated for brackets
LegalExpParser p = new LegalExpParser();
try {
    p.eval(string);
    System.out.println(string + " is OK");
} catch (IllegalArgumentException e) {
    System.out.println(e.getMessage());
}
```

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How Recursion Works

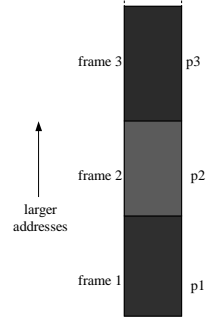
```
int fact (int n) {
    if (n==0) return 1;
    else return n*fact(n-1);
}
```

- How do we prevent the different 'n's from overwriting each other?
 - Give them different locations
- How do we know which 'n' to compute with at each moment?
 - Organize the 'n's into a stack; always use the topmost value

- How does a method-call know where to return to?
 - Save a *return address* before making the call
- How does the *callee* return the result to the *caller*?
 - Result is left on top of the stack
- How do we keep track of all this information?
 - Organize the information for each call into a *frame*

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Stack Frames



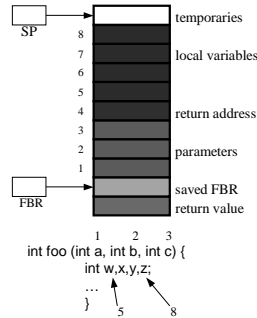
p1 calls p2 which calls p3

- FBR (Frame Base Register) points into current (topmost) frame
- Values within frames are addressed via offsets from FBR
- SP (Stack Pointer) moves up and down as temporary values are pushed and popped

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Typical Stack Frame

- One slot for each parameter (values computed and pushed by *caller*)
- One slot for return address (where to continue execution after method is done)
- One slot for each local variable (allocated by *callee*)
- Temporaries are pushed/popped on top of frame



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How Exceptions Work

- What a Java program does when an exception occurs


```
while not within a try that has a matching catch for this exception {
    Pop the top frame off the stack;
}
Execute the first matching catch-block;
Execute the finally block if there is one;
Continue with following code;
```
- When running a Java application
 - There is a "first" frame that catches all exceptions
 - It prints some semi-helpful information and then halts

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